

Contributions to Persistence Theory

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Abstract. Persistence theory discussed in this paper is an application of algebraic topology (Morse Theory [29]) to Data Analysis, precisely to qualitative understanding of point cloud data, or PCD for short. PCD can be geometrized as a filtration of simplicial complexes (Vietoris-Rips complex [25] [36]) and the homology changes of these complexes provide qualitative information about the data. Bar codes describe the changes in homology with coefficients in a fixed field. When the coefficient field is \mathbb{Z}_2 , the calculation of bar codes is done by ELZ algorithm (named after H. Edelsbrunner, D. Letscher, and A. Zomorodian [20]). When the coefficient field is \mathbb{R} , we propose an algorithm based on the Hodge decomposition [17]. With Dan Burghelea and Tamal K. Dey we developed a persistence theory which involves level sets discussed in Section 4. We introduce and discuss new computable invariants, the “relevant level persistence numbers” and the “positive and negative bar codes”, and explain how they are related to the bar codes for level persistence. We provide enhancements and modifications of ELZ algorithm to calculate such invariants and illustrate them by examples.

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1 Introduction

We view “Persistence Theory” as the Computer Science friendly application of Morse Theory to Data Analysis.

- *Data* considered in this thesis is called PCD (point cloud data) which is, mathematically, a *finite metric space* (Σ, d) of very large cardinality. A PCD often appears as a collection of points in some Euclidean space (\mathbb{R}^D, d) with d being the Euclidean distance. It is hard to visualize a PCD or to study its structure directly, since the cardinality is large and dimension of the Euclidean space in which it embeds can be more than three. However, we can understand qualitative features of a PCD by analyzing the family of simplicial complexes associated to it. In this thesis we use Vietoris-Rips complexes which were introduced first by Vietoris then by Rips in connection with the group theory (Vietoris [36], Hausmann [25]).
- Morse Theory [29] is part of geometric and algebraic topology, whose purpose is to describe the topology of reasonably nice spaces X , for example smooth manifolds, with the help of a generic real valued functions $f : X \rightarrow \mathbb{R}$, for example proper smooth functions with all critical points nondegenerate. It uses the local changes in the homology of the sub-levels $X_{-\infty, t} = f^{-1}((-\infty, t])$ to describe the global topology (homology) of the space X . The topology of the sub-levels for such generic functions changes for a discrete collection of t 's, called critical values. If X is compact then the topology changes for a finite collection of t 's, say $t_0 < t_1 < \dots < t_N$, and a finite filtration $X_{-\infty, t_0} \subseteq \dots \subseteq X_{-\infty, t_N} = X$ is obtained.

Persistence theory, as considered by H. Edelsbrunner, D. Letscher, A. Zomorodian [20], P. Frosini and C. Landi [21], also anticipated in R. Deheuvels [15], provides a slight change of perspective. The object to start with is a space X equipped with a finite filtration rather than a function $f : X \rightarrow \mathbb{R}$. This leads to the concept of *persistence vector space*, and then to *linear algebra of persistence vector spaces*, and to additional type of invariants, *bar codes* [5], [38]. (As recognized by Carlsson and Zamorodian [38] the concept of bar codes corresponds to “torsion and rank” of a finitely generated graded module over the ring of polynomials of one variable. This concept was not previously used in the works of computer scientists.)

When the underlying space X is a simplicial or more general polytopal complex, algorithms of reasonably low complexity [12], [26], [30] as proposed

in [20], [38] can be used to calculate efficiently the *bar codes* of the persistence vector space associated with a filtration of the complex. From the bar codes, one can derive the homology of the sub-levels and then of the underlying space.

Relationship with Data analysis

The Persistence Theory becomes relevant to Data Analysis due to remarkable new ideas:

One can understand the qualitative features of a PCD by analyzing geometric shapes, in this case simplicial complexes, associated to it. One associates to the finite metric space (Σ, d) and $\epsilon \geq 0$ a simplicial complex $R_\epsilon(\Sigma)$ (so called “Vietoris-Rips complex”). Since a single Vietoris-Rips complex $R_\epsilon(\Sigma)$ does not contain enough information of (Σ, d) , one considers all of these complexes for all $\epsilon \geq 0$. Fortunately only finitely many of them are different and they form a finite filtration of the standard simplex of dimension the cardinality of Σ less one. The result of this analysis is encoded in bar codes, which carry information about the qualitative features of the data and signal the missing parts and “accidental noises” in the data.

This theory had already nice applications in many areas of science cf. Carlsson [4] and one expects much more to follow.

Section 2, Mathematical Preliminaries, provides a concise presentation of the mathematical concepts behind the Persistence Theory. In this section, we recall the definitions of simplicial complexes, simplicial homology, polytopal complexes, cellular homology, singular homology and Betti numbers as well as the methods to calculate them when the coefficient field is \mathbb{Z}_2 or \mathbb{R} .

When the field is \mathbb{Z}_2 , the calculation is done by the *ELZ algorithm* (H. Edelsbrunner, D. Letscher, and A. Zomorodian [20]).

When the field is \mathbb{R} , the calculation can be done by the *Hodge decomposition*. This is a new method, at least as long as the calculation of bar codes is concerned, based on the elementary Hodge theory [17].

In Section 3, Persistent Homology of a PCD, one reviews the basic algebra of (tame) persistence vector spaces and introduces the *bar codes* of a tame persistence vector space. These are by now standard concepts but our presentation is slightly different from the existing literatures.

From a point cloud data (PCD), (Σ, d) , one derives a filtration of simplicial complexes (Vietoris-Rips complexes mentioned before)

$$R_{\epsilon_0}(\Sigma) \subseteq R_{\epsilon_1}(\Sigma) \subseteq \cdots \subseteq R_{\epsilon_N}(\Sigma).$$

Then r -dimensional homology groups with coefficients in a fixed field of the spaces of this filtration with the linear maps induced by inclusions define

a tame persistence vector space. The “*bar codes*” of this tame persistence vector space are referred to as the bar codes of the PCD. For the definition of bar code see Subsection 3.3.

When the coefficient field of homology groups is \mathbb{Z}_2 , the calculation of bar codes is done by the ELZ algorithm [20].

When the coefficient field of homology groups is \mathbb{R} , the calculation of bar codes is done by a slight generalization of the Hodge decomposition as described in Subsection 2.2. *This is a new contribution.*

The *simultaneous persistence* used in Section 4 in relation with *level persistence* is also a new contribution, and requires appropriate modifications / improvements of the ELZ algorithm.

In Section 4, Persistence Theory Refined, one considers continuous maps $f : X \rightarrow \mathbb{R}$. When f is weakly tame (Definition 3.11), a finite filtration $X_{-\infty, t_0} \subseteq X_{-\infty, t_1} \subseteq \cdots \subseteq X_{-\infty, t_N}$ is defined, where $X_{-\infty, t_i} = f^{-1}((-\infty, t_i])$ and $t_0 < t_1 < \cdots < t_N$ are critical values. The homology groups in each dimension of the above filtration provide a persistence vector space whose bar codes are referred to as the bar codes of the (weakly tame) map f in that dimension. In Section 4 the “persistence” for such filtration (which is the standard persistence considered in Section 3) is referred to as the *sub-level persistence* and its invariants as the *bar codes for the sub-level persistence* of f . The sub-level persistence analyzes the changes in the homology of sub-level sets.

In Section 4 we refine the sub-level persistence to *level persistence*. The level persistence considers the changes in the homology of the level sets $X_t = f^{-1}(t)$ rather than the changes in the homology of the sub-level sets $X_{-\infty, t} = f^{-1}((-\infty, t])$. In a more primitive form the level persistence was first considered in [16] under the name “interval persistence”. For tame maps (Definition 4.1) the level persistence is equivalent to Zigzag persistence previously introduced in [6].

In this work the maps considered will be tame, i.e. the topology of the levels changes at finitely many t . The theory requires the tameness hypotheses, however almost all continuous maps and in particular all maps of practical interest like simplicial maps or generic smooth maps on smooth manifolds are tame. Tame maps are weakly tame (Corollary 4.3).

The level persistence determines and is determined by the *relevant persistence numbers* introduced in Section 4. For a tame map the relevant persistence numbers are equivalent to the bar codes for the Zigzag persistence, a concept introduced by Carlsson and Silva [6]. We call these bar codes *bar codes for level persistence* and also provide their definition from a slightly different perspective.

For a tame map the relevant level persistence numbers consist of a finite collection of numbers while the bar codes for level persistence (equivalently for Zigzag persistence) consist of a finite collection of intervals of four types: closed, open, left open right closed and left closed right open. For a tame map $f : X \rightarrow \mathbb{R}$ they carry considerably more information than the bar codes for sub-level persistence of the map f . In fact the former implies the latter as explained and illustrated by examples in Section 4.

There are two fundamental concepts in the level persistence: *death* and *detectability* (or *observability*). These concepts should be compared with *birth* and *death* in sub-level persistence.

We also introduce in Section 4 the concept of positive and negative bar codes since they are related to relevant persistence numbers and can be calculated efficiently when the underlying space of the map is a simplicial complex and the map is linear on each simplex.

We reduce the calculation of bar codes for level persistence or of relevant persistence numbers to the calculation of bar codes for sub-level persistence. However the sub-level persistence is not for the map f but for other maps associated with f , which are derived via the construction of *cuts along levels*. We provide new algorithms to calculate “cuts along levels”, which are improvements on the existing ELZ algorithms [20].

Section 4 also contains a number of examples and pictures to describe the implementation of the algorithms.

The level persistence as presented in this thesis was also considered in papers [3] and [16].

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2 Mathematical Preliminaries

In this section we review the definitions of simplicial complex / simplicial homology[24] / polytopal complex / cellular homology[28] and singular homology and the methods to calculate these homologies when the coefficient field¹ is \mathbb{Z}_2 or \mathbb{R} .

2.1 Basic Knowledge

Affinely Independent [13]:

¹The homology considered here is always in a fixed field.

A set of points (vectors) $x_0, \dots, x_k \in \mathbb{R}^D$ is *affinely independent* iff the vectors

$$x_1 - x_0, x_2 - x_0, \dots, x_k - x_0$$

are linearly independent.

Convex Hull [2]:

Let $X = \{x_1, \dots, x_k\}$ be a finite set of points(vectors) in \mathbb{R}^D . The *convex hull* of X is the set of all convex combinations of the elements of X

$$H_{convex}(X) = \left\{ \sum_{i=1}^k \alpha_i x_i \mid \alpha_i \in \mathbb{R}, \alpha_i \geq 0, \sum_{i=1}^k \alpha_i = 1 \right\}$$

Affine Hull [14]:

Let $X = \{x_1, \dots, x_k\}$ be a finite set of points (vectors) in \mathbb{R}^D . The *affine hull* of X is the set of all affine combinations of the elements of X

$$\text{aff}(X) = \left\{ \sum_{i=1}^k \alpha_i x_i \mid \alpha_i \in \mathbb{R}, \sum_{i=1}^k \alpha_i = 1 \right\}$$

Geometric Simplices in a Real Vector Space [31]:

An n -geometric simplex σ in \mathbb{R}^D for $n \leq D$ is the convex hull of $n + 1$ affinely independent points $v_0, \dots, v_n \in \mathbb{R}^D$ and sometimes denoted by $[v_0, \dots, v_n]$. The points v_0, \dots, v_n are called *vertices* of σ .

The convex hull of a subset of vertices of σ containing $m + 1$ points is called an m -face of σ . The m -faces of σ are also m -geometric simplices.

Standard n -Simplex [24]:

The *standard n -simplex* $\Delta^n \in \mathbb{R}^{n+1}$ is the convex hull of vertices

$$v_0 = (1, 0, \dots, 0), \dots, v_n = (0, \dots, 0, 1).$$

Geometric Simplicial Complexes [31]:

A *geometric simplicial complex* $\mathcal{K} = \{ \sigma_\iota \mid \iota \in \mathcal{A} \}$ in \mathbb{R}^D is a collection of geometric simplices in \mathbb{R}^D that satisfies the following conditions:

1. Any face of a simplex of \mathcal{K} is also in \mathcal{K} .
2. The intersection of any two simplices $\sigma_1, \sigma_2 \in \mathcal{K}$ is a face of both σ_1 and σ_2 .

Underlining Space of a Geometric Simplicial Complex [31]:

Let's denote the union of simplices of a geometric simplicial complex \mathcal{K} as $|\mathcal{K}|$ and call it the *underlying space* of \mathcal{K} .

Simplicial Maps Between Geometric Simplicial Complexes [31]:

Let \mathcal{K} and \mathcal{L} be two geometric simplicial complexes. A map $f : |\mathcal{K}| \rightarrow |\mathcal{L}|$ is a *simplicial map* if it sends each simplex of \mathcal{K} to a simplex of \mathcal{L} by a linear map taking vertices to vertices.

Abstract Simplicial Complexes [31]:

An *abstract simplicial complex* $\mathcal{X} = (V, \mathcal{S})$ consists of a set V of vertices and a set \mathcal{S} of finite subsets of V called *abstract simplices*, which satisfy the following property: if $\sigma \in \mathcal{S}$ and $\tau \subseteq \sigma$, then $\tau \in \mathcal{S}$.

If $\sigma \in \mathcal{S}$ has $n + 1$ elements, the dimension of σ is n and σ is called an n -simplex.

If $\tau \subseteq \sigma$ has $m + 1$ elements, τ is called an m -face of σ .

Simplicial Maps Between Abstract Simplicial Complexes [31]:

Let $\mathcal{X}_1 = (V_1, \mathcal{S}_1)$ and $\mathcal{X}_2 = (V_2, \mathcal{S}_2)$ be two abstract simplicial complexes. A map $f : V_1 \rightarrow V_2$ is a *simplicial map* if whenever the vertices v_0, \dots, v_n of \mathcal{X}_1 span a simplex, the points $f(v_0), \dots, f(v_n)$ are vertices of a simplex of \mathcal{X}_2 .

Spatial Realization of an Abstract Simplicial Complex [31]:

A spatial realization $\mathcal{K} = \langle \mathcal{X} \rangle$ of an abstract simplicial complex \mathcal{X} is provided by an injective map $f : V \rightarrow \mathbb{R}^D$ that sends each simplex in \mathcal{X} to an affinely independent set in \mathbb{R}^D so that no interior of different geometric simplices intersect (in particular if $f(v_i)$ are all affinely independent). The union of the convex hulls of the images of all abstract simplices of \mathcal{X} is a geometric simplicial complex, which is a spatial realization of \mathcal{X} .

Spatial realizations always exist for D large enough and their underlying spaces are all homeomorphic to each other. If we have a geometric simplicial complex \mathcal{K} , we can also get an abstract simplicial complex \mathcal{X} with V the set of vertices of \mathcal{K} and

$$\mathcal{S} = \{ \sigma \subseteq V \mid \sigma \text{ are vertices of some simplex of } \mathcal{K} \}.$$

Then (V, \mathcal{S}) defines an abstract simplicial complex \mathcal{X} .

It is not hard to check the above assignments are essentially inverses to each other and they also induce a bijection between simplicial maps, so an abstract simplicial complex \mathcal{X} and its spatial realization \mathcal{K} can be regarded as equivalent objects.

Oriented Simplex (page 26, [31]):

Let σ be a simplex (geometric or abstract). One says that two orderings of its vertices are equivalent if they differ from one another by an even

permutation. If σ is a 0-simplex then, obviously, the ordering is unique. If $\dim \sigma > 0$ then the orderings of the vertices of σ fall into two equivalence classes. Each of these classes is called an *orientation* of σ . An *oriented simplex* (σ, ϵ) is a simplex σ together with an orientation ϵ .

If the points v_0, \dots, v_p are affinely independent, we shall use the symbol

$$\langle v_0, \dots, v_p \rangle$$

to denote the oriented simplex consisting of the simplex $\{v_0, \dots, v_p\}$ and the equivalence class of the ordering (v_0, \dots, v_p) .

Simplicial Homology with Coefficients in a Field [24]:

Let \mathcal{X} be an abstract simplicial complex as defined above. Let $C_n^\Delta(\mathcal{X}; \kappa)$ be the κ -vector space generated by the set of n -simplices in \mathcal{X} , where κ is a fixed field. Elements of $C_n^\Delta(\mathcal{X}; \kappa)$ are finite formal sums $\sum_i \lambda_i \sigma_i$ where $\lambda_i \in \kappa$, σ_i are n -simplices of \mathcal{X} . By choosing a total ordering of vertices of \mathcal{X} (which provides an orientation on each simplex) a boundary map $\partial_n^\Delta : C_n^\Delta(\mathcal{X}; \kappa) \rightarrow C_{n-1}^\Delta(\mathcal{X}; \kappa)$ is defined as the linear extension of the map below

$$\partial_n^\Delta(\langle v_0, \dots, v_n \rangle) = \sum_i (-1)^i \langle v_0, \dots, \hat{v}_i, \dots, v_n \rangle.$$

It's easy to verify that $\partial_n^\Delta \circ \partial_{n+1}^\Delta = 0$, so that we can define the *simplicial homology groups* as $\mathcal{H}_n^\Delta(\mathcal{X}; \kappa) = \ker \partial_n^\Delta / \text{im} \partial_{n+1}^\Delta$.

Choosing different orderings on the vertices of \mathcal{X} or different orientations on each simplex, the boundary maps might change (the sign of entries might change) but the homology groups remain the same.

Convex Polytope [23]:

An *extreme point* of a convex set S in a real vector space is a point in S which does not lie in any open line segment joining two points of S . The extreme points of a simplex in a real vector space are its vertices.

A *convex polytope* is the convex hull of a finite set of points. We call extreme points of a polytope its *vertices*. A simplex is a convex polytope such that its vertices are affinely independent.

Suppose P is a convex polytope with X the set of its vertices. The dimension of P is the dimension of the affine hull of X .

A point $x \in P$ is an *interior point* of P if it is interior in the sense of point set topology when P is regarded as a subspace of $\text{aff}(X)$.

Given $a_1, \dots, a_D, b \in \mathbb{R}$, consider the closed half-space defined by

$$\{ (r_1, \dots, r_D) \in \mathbb{R}^D \mid a_1 r_1 + \dots + a_D r_D \leq b \}$$

The boundary of the above half-space is

$$\{ (r_1, \dots, r_D) \in \mathbb{R}^D \mid a_1 r_1 + \dots + a_D r_D = b \}$$

A *face* of a convex polytope is any intersection of the polytope with a closed half-space such that no interior points of the polytope lie on the boundary of the half-space.

The dimension of a face is the dimension of its affine hull. Faces of a convex polytope are also convex polytopes. Vertices of a polytope are its 0-dimensional faces.

Polytopal Complex [23]:

Polytopal complex, often named as polyhedra complex or cellular complex, consists of a collection of convex polytopes in some Euclidean space \mathbb{R}^D , satisfying two conditions:

- (i) Every face of a polytope in Π is also in Π .
- (ii) The intersection of any two polytopes in Π is a face of both.

The polytopes in Π are called *cells*. An n -dimensional polytope in Π is called an n -cell.

For example, the set of all faces of a convex polytope defines a polytopal complex.

A geometric simplicial complex is a polytopal complex in which every cell is a geometric simplex.

Oriented Convex Polytope [28]:

Let σ be a d -dimensional polytope in \mathbb{R}^D and $V(\sigma)$ be the set of vertices of σ . The map $\epsilon : \binom{V(\sigma)}{d+1} \rightarrow \{1, -1\}$ is an *orientation* of σ if the following conditions are satisfied:

(1) $\epsilon(v_0, \dots, v'_i, \dots, v_d) = \epsilon(v_0, \dots, v_i, \dots, v_d)$ if v_i and v'_i are in the same open half-space of \mathbb{R}^D delimited by one of the supporting hyperplanes of $v_0, \dots, v_i, \dots, v_d$ (hyperplanes that pass through $v_0, \dots, v_i, \dots, v_d$) and $\epsilon(v_0, \dots, v'_i, \dots, v_d) = -\epsilon(v_0, \dots, v_i, \dots, v_d)$ if not.

(2) $\epsilon(v_0, v_1, \dots, v_d) = \text{sign}(\pi) \epsilon(v_{\pi(0)}, v_{\pi(1)}, \dots, v_{\pi(d)})$ for every permutation π .

0-dimensional polytope admits a unique orientation. Every d -dimensional ($d \geq 1$) polytope admits exactly two distinct orientations. If ϵ is one of them, $-\epsilon$ is the other one.

An orientation ϵ on σ induces an orientation $\epsilon|_\tau$ on its codimension-1 face τ . The convention is that a base in the supporting hyperplane of τ followed by a vector pointing towards the interior of σ provides a base of the supporting affine space of σ which defines the orientation ϵ .

Let σ be an n -cell, when τ is a $(n-1)$ -face of σ , one defines the orientation ϵ_τ by

$$\epsilon|_\tau(v_0, \dots, v_{n-1}) := \epsilon(v_0, \dots, v_{n-1}, v_n)$$

for any $v_0, \dots, v_{n-1} \in V(\tau)$ and $v_n \in V(\sigma) \setminus V(\tau)$.

Cellular Homology with Coefficients in a Field [28]:

Let Π be a polytopal complex as defined above. Let $C_n(\Pi; \kappa)$ be the vector space generated by the set of n -cells in Π with coefficients in a fixed field κ . Elements of $C_n(\Pi; \kappa)$ are finite formal sums $\sum_i \lambda_i \sigma_i$ where $\lambda_i \in \kappa$ and σ_i are n -cells of Π .

For any cell σ choose an orientation $\epsilon(\sigma)$.

The boundary map $\partial_n : C_n(\Pi; \kappa) \rightarrow C_{n-1}(\Pi; \kappa)$ is defined as the linear extension of the map below

$$\partial_n(\sigma) = \sum_{\tau} I_{\sigma, \tau} \tau.$$

$I_{\sigma, \tau} = 0$ if τ is not a $(n-1)$ -face of σ and $I_{\sigma, \tau} = \pm 1$ if τ is a $(n-1)$ -face of σ and $\epsilon(\sigma)|_\tau = \pm \epsilon(\tau)$, where $\epsilon(\sigma)$ and $\epsilon(\tau)$ are orientations of σ and τ respectively.

We can verify that $\partial_n \circ \partial_{n+1} = 0$, so the *polytopal homology groups* are defined as $\mathcal{H}_n(\Pi; \kappa) = \ker \partial_n / \text{im} \partial_{n+1}$. $\mathcal{H}_n(\Pi; \kappa)$ is independent of the chosen orientation although the boundary maps ∂_n do depend.

When $\kappa = \mathbb{Z}_2$, orientation becomes irrelevant and a boundary map is defined by

$$\partial_n(\sigma) = \sum_i \lambda_i \sigma_i$$

where σ is an n -cell and σ_i 's are $(n-1)$ -cells, $\lambda_i = 1$ iff σ_i is a face of σ .

Note. There is a canonical subdivision of a polytopal complex into a simplicial complex. Each cell σ can be decomposed into a simplicial complex. Choose an interior point x_σ of the cell σ . The cell σ can be decomposed as the union of cones with base codimension-1 faces of σ and apex x_σ . It's obvious that 2-dimension cell can be decomposed as a simplicial complex. By induction, the cell σ can be decomposed into a simplicial complex. This decomposition is possible because all cells are convex polytopes. Once each cell is decomposed into a simplicial complex, we can regard the polytopal complex as the union of those simplicial complexes, hence as a simplicial complex.

Once we decompose a polytopal complex into a simplicial complex, we can calculate the homology groups of this simplicial complex. In fact we can use this simplicial homology to replace the homology of the polytopal

complex, since the underlying spaces of the simplicial complex and polytopal complex are the same and two complexes with the same underlying spaces have the same homology groups (as pointed below).

It is easier to write programs to get the boundary maps and calculate the homology groups of a simplicial complex than a polytopal complex. However, it costs much more time to calculate the homology groups by first decomposing a polytopal complex into a simplicial complex than to calculate directly without decomposition, since the decomposition increases number of cells greatly.

It turns out that the homology of a simplicial or polytopal complex depends only on the underlying topological space. The easiest way to see this is to define the homology for a topological space, independent of simplices or cells, and show that it leads to the same results as the one described above. The homology defined in this way is known as *singular homology* and is defined below.

Definition 2.1. Singular Homology with Coefficients in a Field(see page 108, 153 [24])

A singular n -simplex in a topological space X is defined as a continuous map $\sigma : \Delta^n \rightarrow X$ on the standard n -simplex. Let $C_n(X; \kappa)$ be the vector space with basis the set of singular n -simplices in X and coefficients in a fixed field κ . Elements of $C_n(X; \kappa)$ are finite formal sums $\sum_i \lambda_i \sigma_i$ for $\lambda_i \in \kappa$ and $\sigma_i : \Delta^n \rightarrow X$. A boundary map $\partial_n : C_n(X; \kappa) \rightarrow C_{n-1}(X; \kappa)$ is defined by the formula:

$$\partial_n(\sigma) = \sum_i (-1)^i \sigma|[v_0, \dots, \hat{v}_i, \dots, v_n]$$

where $\sigma : \Delta^n \rightarrow X$ is a singular n -simplex, v_0, \dots, v_n are vertices of Δ^n , \hat{v}_i means v_i is deleted. Implicit in this formula is the canonical identification of $[v_0, \dots, \hat{v}_i, \dots, v_n]$ with Δ^{n-1} , preserving the ordering of vertices, so that $\sigma|[v_0, \dots, \hat{v}_i, \dots, v_n]$ is regarded as a map $\Delta^{n-1} \rightarrow X$, that is, a singular $(n-1)$ -simplex.

It's easy to verify that $\partial_n \circ \partial_{n+1} = 0$, so that we can define the singular homology group with coefficients in κ as $\mathcal{H}_n(X; \kappa) = \ker \partial_n / \text{im} \partial_{n+1}$.

$\mathcal{H}_n(X; \kappa)$ is a vector space since κ is a field.

Definition 2.2. Betti Numbers

Given a field κ one can define $b_n(X; \kappa)$, the n -th Betti number with coefficients in κ , as the dimension of the vector space $\mathcal{H}_n(X; \kappa)$.

We will show how to calculate Betti numbers of a polytopal complex Π with \mathbb{Z}_2 or \mathbb{R} coefficients.

When the coefficients are in \mathbb{Z}_2 , we have the following method referred to below as the *ELZ algorithm* ([19]).

First order all cells, so that $i < j$ if σ_i is a face of σ_j . Then form a boundary matrix ∂ , so that the entry on i -th row and j -th column is

$$\partial[i, j] = \begin{cases} 1 & \text{if } \sigma_i \text{ is a codimension-1 face of } \sigma_j; \\ 0 & \text{otherwise.} \end{cases}$$

The matrix ∂ is an upper triangular matrix with zeros on diagonal. Let $low(j)$ be the row index of the lowest 1 in column j . If the entire column is zero, then $low(j)$ is undefined. We call a matrix *reduced* if $low(j) \neq low(j_0)$ whenever j and j_0 , with $j \neq j_0$, specify two non-zero columns. The algorithm below starts with the boundary matrix ∂ and changes it by adding columns from left to right. All matrices in this process are upper triangular with zeros on diagonal. Finally we get a reduced matrix R .

Algorithm 2.1. The ELZ Algorithm

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R = ∂
for j = 1 to m do
  while there exists j0 < j with low(j0) = low(j) do
    add column j0 to column j
  endwhile
endfor

```

From the reduced matrix R one can read off the Betti numbers. It can be proven [19] that the zero columns of R correspond to generators of cycles and non-zero columns of R correspond to generators of boundaries. Since the homology group is quotient of cycles over boundaries, the Betti number is the number of generators of cycle minus the number of generators of boundary.

Therefore the k -th Betti number

$$b_k(\Pi; \mathbb{Z}_2) = \begin{aligned} & \text{the number of zero columns which correspond to a } k\text{-cycle} \\ & - \text{the number of non-zero columns which correspond to a} \\ & (k+1)\text{-boundary} \end{aligned}$$

Algorithms for the calculation of Betti numbers with coefficient in any finite field also exist, however they are slightly more complex.

When the coefficients are in \mathbb{R} , we develop a different method using the Hodge decomposition introduced below.

According to Observation 2.7 at the end of Subsection 2.2, the k -th Betti number can be calculated by the following formula

$$b_k(\Pi; \mathbb{R}) = \text{number of } k\text{-cells in } \Pi - \text{rank}(\partial_{k+1}) - \text{rank}(\partial_k).$$

2.2 Hodge Decomposition of a Chain Complex with \mathbb{R} Coefficients

The constructions below follow the standard “Hodge decomposition” familiar in Riemannian geometry. This finite dimensional elementary formulation was first considered by B. Eckmann [17]. One starts with a complex of finite dimensional \mathbb{R} -vector spaces

$$\cdots \xrightarrow{\partial_{r+2}} C_{r+1} \xrightarrow{\partial_{r+1}} C_r \xrightarrow{\partial_r} C_{r-1} \xrightarrow{\partial_{r-1}} \cdots \xrightarrow{\partial_2} C_1 \xrightarrow{\partial_1} C_0 \xrightarrow{\partial_0} C_{-1} = 0.$$

The vector spaces in this complex are equipped with positive definite inner products. If C_r 's are equipped with a base, one can take the unique inner product which makes this base orthonormal. In the cases under consideration, each complex comes from a finite simplicial / polytopal complex L , with C_r the \mathbb{R} -vector space generated by the r -simplices / r -cells of L and ∂_r the boundary maps. $\partial_{r+1} : C_{r+1} \rightarrow C_r$ are linear operators between inner product spaces. Let $\delta_r = \partial_{r+1}^*$ be the adjoint operator of ∂_{r+1} .

Lemma 2.1. *If A and B are two finite dimensional inner product spaces, $f : A \rightarrow B$ is a linear map and $f^* : B \rightarrow A$ is its adjoint, i.e. $\langle b, f(a) \rangle = \langle f^*(b), a \rangle$ for any $a \in A$ and $b \in B$, then*

- i) $\ker(f) = (\text{im}(f^*))^\perp$;*
- ii) For any $a \in A$, if $f \circ f^* \circ f(a) = 0$ then $f(a) = 0$.*

Define $\Delta_r = \partial_{r+1} \circ \delta_r + \delta_{r-1} \circ \partial_r : C_r \rightarrow C_r$ for $r \in \mathbb{Z}_{\geq 0}$, $(C_r)_+ = \text{im}(\partial_{r+1})$, $(C_r)_- = \text{im}(\delta_{r-1})$ and $H_r = \ker(\Delta_r)$.

Proposition 2.2. *i) $H_r = \ker(\delta_r) \cap \ker(\partial_r)$;*

*ii) (**Hodge Decomposition**) $C_r = (C_r)_+ \oplus H_r \oplus (C_r)_-$, where $(C_r)_+$, H_r and $(C_r)_-$ are pairwise orthogonal;*

iii) There is a canonical isomorphism

$$j_r : H_r = \ker(\delta_r) \cap \ker(\partial_r) \rightarrow \mathcal{H}_r = \ker(\partial_r) / \text{im}(\partial_{r+1}).$$

We will give an algorithm for calculating the Hodge decomposition, i.e., given a chain complex $\mathcal{C}(L)$, we will calculate the three orthogonal projections:

$$(p_r)_+ : C_r \rightarrow C_r \quad \text{with } (p_r)_+(C_r) = (C_r)_+$$

$$(p_r)_- : C_r \rightarrow C_r \quad \text{with } (p_r)_-(C_r) = (C_r)_-$$

and

$$(p_r)_H : C_r \rightarrow C_r \quad \text{with } (p_r)_H(C_r) = H_r$$

for each $r \geq 0$ respectively.

Lemma 2.3. *Given any $m \times n$ matrix A over \mathbb{R} , if the rank of A is k then there exists an $m \times k$ matrix $[A]$ of the form*

$$[A] = (v_1, v_2, \dots, v_k)_{m \times k} \quad (2.1)$$

where $\{v_1, v_2, \dots, v_k\}$ is a collection of orthonormal column vectors which is equivalent to the collection of column vectors of A .

Two collections of vectors are equivalent if they generate the same subspace. Moreover, there is a canonical construction of such orthonormal column vectors known as Gram-Schmidt Orthonormalization.

Note. 1. Given a linear map $A : \mathbb{R}^n \rightarrow \mathbb{R}^m$, one view A as a $m \times n$ matrix with respect to the standard basis of \mathbb{R}^n and \mathbb{R}^m . The collection of column vectors of $[A]$ represents an orthonormal basis of $\text{im}(A)$.

2. **Matlab** contains a function **orth** with input an $m \times n$ matrix A and output a matrix $[A]$.

3. $[A][A]^T$ is unique although $[A]$ is not (See **Lemma 2.4**).

Lemma 2.4. *Given a linear map $A : \mathbb{R}^n \rightarrow \mathbb{R}^m$, the linear map*

$$\begin{aligned} p_A : \mathbb{R}^m &\rightarrow \mathbb{R}^m \\ y &\mapsto [A][A]^T y \end{aligned}$$

is the orthogonal projection on $\text{im}(A)$.

Proposition 2.5. *Given a chain complex $\mathcal{C}(L)$ of a finite simplicial/polytopal complex L over \mathbb{R}*

$$\cdots \xrightarrow{\partial_{r+2}} C_{r+1} \xrightarrow{\partial_{r+1}} C_r \xrightarrow{\partial_r} C_{r-1} \xrightarrow{\partial_{r-1}} \cdots \xrightarrow{\partial_2} C_1 \xrightarrow{\partial_1} C_0 \xrightarrow{\partial_0} C_{-1} = 0$$

Each ∂_r can be regarded as an $n_{r-1} \times n_r$ matrix with respect to the standard basis of C_r and C_{r-1} . The following linear maps are orthogonal projections onto $(C_r)_+, (C_r)_-$ and H_r , respectively.

$$\begin{aligned} (p_r)_+ : C_r &\rightarrow C_r \\ y &\mapsto [\partial_{r+1}][\partial_{r+1}]^T y \\ (p_r)_- : C_r &\rightarrow C_r \\ y &\mapsto [(\partial_r)^T][(\partial_r)^T]^T y \\ (p_r)_H : C_r &\rightarrow C_r \\ y &\mapsto (I_{n_r} - [\partial_{r+1}][\partial_{r+1}]^T - [(\partial_r)^T][(\partial_r)^T]^T) y \end{aligned}$$

The linear map

$$\begin{aligned} k_r : \mathcal{H}_r = \ker(\partial_r)/\text{im}(\partial_{r+1}) &\rightarrow H_r = \ker(\delta_r) \cap \ker(\partial_r) \\ y + \text{im}(\partial_{r+1}) &\mapsto (p_r)_H(y) \end{aligned}$$

is the inverse of j_r , which verifies Proposition 2.2 iii).

Observation 2.6. *The rank of a real matrix A equals to the rank of AA^T or $A^T A$.*

By the Hodge decomposition, $\dim(H_r) = \dim(C_r) - \dim((C_r)^+) - \dim((C_r)^-)$. By the above observation, $\dim((C_r)^+) = \text{rank}([\partial_{r+1}][\partial_{r+1}]^T) = \text{rank}([\partial_{r+1}]) = \text{rank}(\partial_{r+1})$. We get $\dim((C_r)^-) = \text{rank}(\partial_r)$ similarly.

Observation 2.7. *The r -th Betti number*

$$b_r(L; \mathbb{R}) = \dim(\mathcal{H}_r(L; \mathbb{R})) = \dim(C_r) - \text{rank}(\partial_{r+1}) - \text{rank}(\partial_r).$$

3 Persistent Homology of a PCD

3.1 Introduction

Informally, we call a finite set of points $X \subseteq \mathbb{R}^m$ a *point cloud data* (PCD for short). A PCD can be regarded as a finite metric space (see Definition 3.1).

A first new idea in Data Analysis is to regard a PCD X as a filtration of the standard n -simplex Δ^n ($n = \sharp X - 1$)² via a construction known as “Vietoris-Rips complex” (see Subsection 3.2). The first term of the filtration is the 0-skeleton of Δ^n or X itself, the last term is Δ^n with the remaining components giving an idea of the qualitative features of the set $X \subseteq \mathbb{R}^m$.

Another new idea is to use homology to describe the topological changes of the components of this filtration (topological features “are born” in some component, and “die” in some other component). This leads to the persistent homology of this filtered simplicial complex. The persistent homology provides the tools to measure and explain the qualitative patterns of a PCD. The role of the numerical invariants Betti numbers, when the homology of a space (simplicial / polytopal complex) is considered, is taken by the invariants **bar codes**, when the persistent homology of a filtered simplicial / polytopal complex is under consideration. The homology considered in this work is always with coefficients in a field κ ($\kappa = \mathbb{Z}_2$ or \mathbb{R}), so the homology groups are actually vector spaces. The linear algebra for persistent homology is “persistent linear algebra” discussed in Subsection 3.3. The invariant “dimension” for a vector space is replaced by the invariant “bar codes” for a persistence vector space. The *bar codes* provides a complete invariant for a persistence vector space (cf. Theorem 3.5) as the *dimension* provides a complete invariant for a vector space. For a filtered simplicial / cell complex as opposed to simplicial / cell complex, we will have “bar codes” derived from persistent

²Here and in the remaining part of this text “ \sharp ” will denote “cardinality”.

homology as opposed to Betti numbers derived from homology. The bar codes can be explicitly calculated by algorithms of the same complexity as the algorithms used for the calculation of the Betti numbers. For $\kappa = \mathbb{Z}_2$, the matrix reduction and pairing algorithm described in [19] and referred below as the ELZ algorithm is the one we will use. For $\kappa = \mathbb{R}$, we will use elementary Hodge theory to produce a new algorithm for the calculation of the bar codes.

Subsection 3.2 discusses PCD and Vietoris-Rips filtration. Subsection 3.3 discusses persistent linear algebra, including persistence vector spaces and their bar codes. The essential features of a persistence vector space are the concept of “birth and death time”³ of its elements and the fact that the “bar codes” provides a complete invariant which carries significant information about “birth and death time” of its elements. Subsection 3.4 defines persistence vector space $\mathcal{H}_r(\mathcal{K})$ and the bar codes $\mathcal{B}(\mathcal{K})$ of a filtered simplicial/polytopal complex \mathcal{K} and uses the Hodge decomposition to calculate $\mathcal{B}(\mathcal{K})$. Subsection 3.5 gives algorithms to calculate the bar codes of a PCD with coefficients $\kappa = \mathbb{Z}_2$ or \mathbb{R} . In case of the field \mathbb{Z}_2 , the algorithm was introduced in [20]. Subsection 3.6 shows a numerical experiment of the bar codes of a PCD with coefficients $\kappa = \mathbb{R}$.

3.2 PCD, Vietoris-Rips Complex and Filtration

Informally, we call a finite set of points $X \subseteq \mathbb{R}^m$ a point cloud data (PCD for short).

Example 3.1. A PCD in \mathbb{R}^2 with 5 points (see Figure 1).

The inclusion $X \subseteq \mathbb{R}^m$ defines a finite metric space (X, d) with d being the Euclidean distance in \mathbb{R}^m . From mathematical point of view we will use the following definition

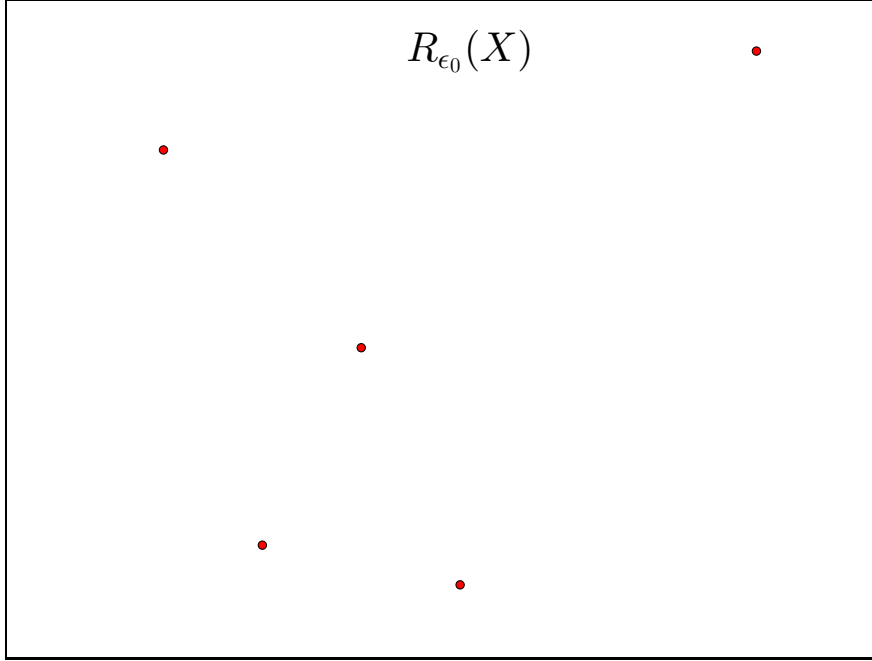
Definition 3.1. A point cloud data is a finite metric space $X = (X, d)$.

Let X be a finite metric space. Given $\epsilon \geq 0$, the Vietoris-Rips complex $R_\epsilon(X)$ of PCD X has X as the set of vertices⁴. A k -simplex is any subset of vertices $\sigma = [x_0, x_1, \dots, x_k]$ with the property that $d(x_i, x_j) \leq \epsilon$ for all pairs $x_i, x_j \in \sigma$. One obtains in this way an *abstract simplicial complex*.

Notice that Vietoris-Rips complex will be determined by its one skeleton. If $\epsilon < \epsilon'$ then there is an inclusion $R_\epsilon(X) \hookrightarrow R_{\epsilon'}(X)$.

³Time here means the index of the component of the filtration.

⁴Other types of simplicial complexes can be associated to the metric space and ϵ but they calculate essentially the same invariants and are most often less economical. They will not be discussed here.

Figure 1: A PCD in \mathbb{R}^2 with 5 points.

Since a PCD X is a finite set, Vietoris-Rips complex $R_\epsilon(X)$ will change at only finitely many epsilons

$$0 = \epsilon_0 < \epsilon_1 < \cdots < \epsilon_N = \sup_{x, x' \in X} (d(x, x')). \quad (3.0)$$

We define the filtration of Vietoris-Rips complexes of PCD X as

$$R_{\epsilon_0}(X) \subseteq R_{\epsilon_1}(X) \subseteq \cdots \subseteq R_{\epsilon_N}(X). \quad (3.1)$$

In this filtration, $R_{\epsilon_0}(X)$ is a zero dimensional simplicial complex with $\sharp X$ vertices and $R_{\epsilon_N}(X)$ is a $\sharp X - 1$ standard simplex.

Example 3.2. For the example of PCD described in Figure 1, one obtain eleven ϵ' s : $\epsilon_0, \epsilon_1, \dots, \epsilon_{10}$ and a filtration with eleven components as indicated below.

The invariant we will calculate will be a collection of intervals for any r , $0 \leq r \leq \dim R_{\epsilon_N}(X)$. Each interval will have as *ends* the numbers $\epsilon_0, \dots, \epsilon_N, +\infty$ or equivalently $0, 1, \dots, N$, or $+\infty$ with the convention that i specifies the number ϵ_i .

If we are interested in the bar codes for $r \leq m - 1$, one can work with $R_\epsilon(X, m)$, the m -skeleton of $R_\epsilon(X)$. The m -restricted filtered simplicial complex

$$R_{\epsilon_0}(X, m) \subseteq R_{\epsilon_1}(X, m) \subseteq \cdots \subseteq R_{\epsilon_N}(X, m). \quad (3.2)$$

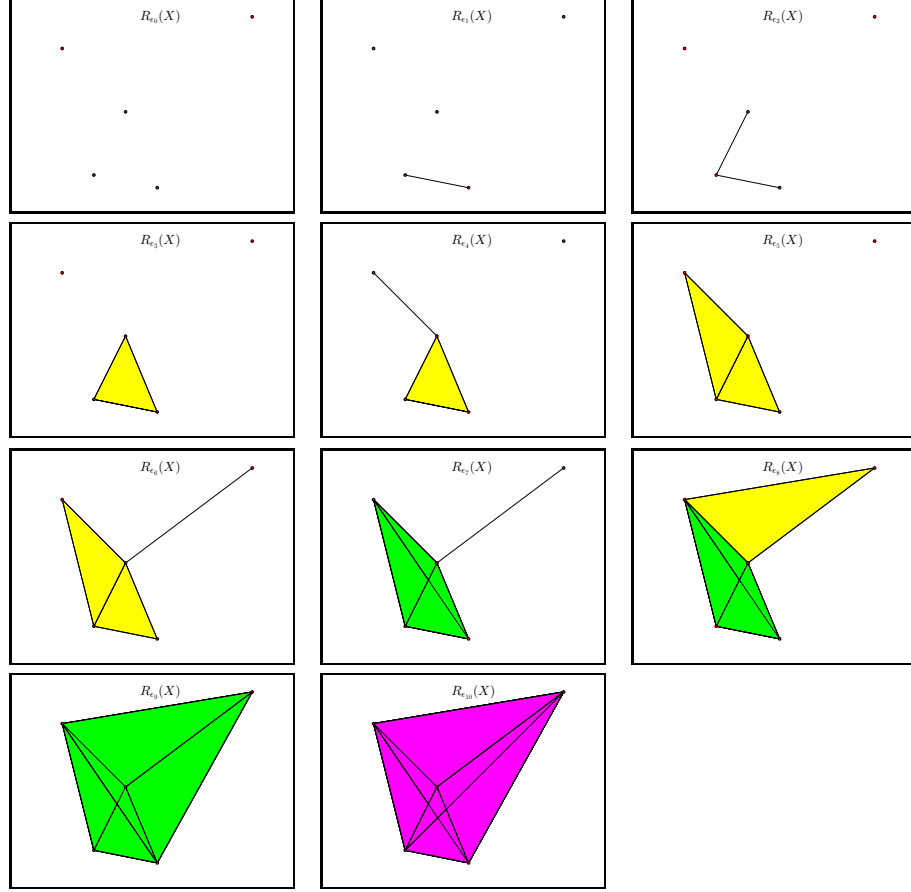


Figure 2: Filtration of Vietoris-Rips complexes of the PCD in Example 3.1.

has the same persistent homology and bar codes as the original filtration (3.1) up to dimension $m-1$ (see Subsection 3.4 for the definition of persistent homology and bar codes). This supposes a much smaller amount of data to be stored in a computer and permits to calculate all bar codes for $r \leq m-1$.

Given a positive integer number $P (P \leq N)$, we can make a further restriction of the filtration (3.2) by stopping at level P

$$R_{\epsilon_0}(X, m) \subseteq R_{\epsilon_1}(X, m) \subseteq \cdots \subseteq R_{\epsilon_P}(X, m). \quad (3.3)$$

This will permit the calculation of all bar codes for $r \leq m-1$ with both ends less than P or the bar codes for $r \leq m-1$ with left end less than P and right end $\geq P$.

3.3 Persistent Linear Algebra (a Gentle Introduction)

Definition 3.2. 1) A persistence vector space \mathcal{V} is a sequence

$$\{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$$

with V_n vector spaces over a field κ and φ_n linear maps.

A persistence vector space is **tame** iff each V_n has finite dimension and φ_n is an isomorphism for n large enough.

The main features of this concept are “birth and death time”⁵ for elements and information about these is provided by the bar codes.

2) A linear map of persistence vector spaces $\omega : \mathcal{V} \rightarrow \mathcal{W}$, where $\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$ and $\mathcal{W} = \{W_n, \phi_n : W_n \rightarrow W_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$ is a commutative diagram

$$\begin{array}{ccccccc} V_0 & \xrightarrow{\varphi_0} & V_1 & \xrightarrow{\varphi_1} & \dots & \xrightarrow{\varphi_{n-1}} & V_n & \xrightarrow{\varphi_n} & V_{n+1} & \xrightarrow{\varphi_{n+1}} & \dots \\ \downarrow \omega_0 & & \downarrow \omega_1 & & & & \downarrow \omega_n & & \downarrow \omega_{n+1} & & \\ W_0 & \xrightarrow{\phi_0} & W_1 & \xrightarrow{\phi_1} & \dots & \xrightarrow{\phi_{n-1}} & W_n & \xrightarrow{\phi_n} & W_{n+1} & \xrightarrow{\phi_{n+1}} & \dots \end{array}$$

where ω_n is a linear map from V_n to W_n for each $n \geq 0$.

A linear map of persistence vector spaces $\omega : \mathcal{V} \rightarrow \mathcal{W}$ is an isomorphism if there exists another linear map of persistence vector spaces $\omega' : \mathcal{W} \rightarrow \mathcal{V}$ such that $\omega' \circ \omega : \mathcal{V} \rightarrow \mathcal{V}$ and $\omega \circ \omega' : \mathcal{W} \rightarrow \mathcal{W}$ are identities. \mathcal{V} and \mathcal{W} are isomorphic if there is an isomorphism between them.

The existence of a linear map $\omega : \mathcal{V} \rightarrow \mathcal{W}$ so that each component $\omega_n : V_n \rightarrow W_n$ is an isomorphism for every $n \geq 0$ implies ω is an isomorphism. One takes $\omega'_n = (\omega_n)^{-1}$.

3) Let $\mathcal{U} = \{U_n, \psi_n : U_n \rightarrow U_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$, $\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$ and $\mathcal{W} = \{W_n, \phi_n : W_n \rightarrow W_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$ be persistence vector spaces. A short exact sequence of persistence vector spaces

$$0 \longrightarrow \mathcal{U} \xrightarrow{\mu} \mathcal{V} \xrightarrow{\nu} \mathcal{W} \longrightarrow 0$$

is a sequence of linear maps of persistence vector spaces such that

$$0 \longrightarrow U_n \xrightarrow{\mu_n} V_n \xrightarrow{\nu_n} W_n \longrightarrow 0$$

is short exact for all $n \geq 0$.

The short exact sequence $0 \longrightarrow \mathcal{U} \xrightarrow{\mu} \mathcal{V} \xrightarrow{\nu} \mathcal{W} \longrightarrow 0$ splits if there exists a linear map $\alpha : \mathcal{V} \rightarrow \mathcal{U}$ such that $\alpha \circ \mu = \text{identity}$ or if there exists a linear map $\beta : \mathcal{W} \rightarrow \mathcal{V}$ such that $\nu \circ \beta = \text{identity}$.

⁵Time here means the index of the component of the filtration.

Note. The alternative definitions are equivalent. Given $\alpha : \mathcal{V} \rightarrow \mathcal{U}$ such that $\alpha \circ \mu = \text{identity}$, we can define $\beta : \mathcal{W} \rightarrow \mathcal{V}$ such that $\nu \circ \beta = \text{identity}$ and vice versa.

i) Let $\alpha : \mathcal{V} \rightarrow \mathcal{U}$ such that $\alpha \circ \mu = \text{identity}$. For any $x \in W_n$, there exists $x' \in V_n$ such that $\nu_n(x') = x$. Define

$$\begin{aligned} \beta_n : W_n &\rightarrow V_n \\ x &\mapsto x' - \mu_n \circ \alpha_n(x'). \end{aligned}$$

To see that β_n is well defined, i.e. independent of the choice of x' , consider

$$\begin{aligned} \gamma_n : V_n &\rightarrow U_n \oplus W_n \\ x &\mapsto (\alpha_n x, \nu_n x). \end{aligned}$$

γ_n is injective. Indeed if $x \in \ker(\gamma_n)$ hence $\alpha_n x = 0$ and $\nu_n x = 0$. Since $\text{im } \mu_n = \ker \nu_n$, there exists $y \in U_n$ such that $x = \mu_n(y)$. Then $y = \alpha_n \circ \mu_n(y) = \alpha_n(x) = 0$, so $x = \mu_n(0) = 0$.

If $x'' \in V_n$ such that $\nu_n(x'') = x$, then γ_n will map both $x' - \mu_n \circ \alpha_n(x')$ and $x'' - \mu_n \circ \alpha_n(x'')$ to $(0, x)$. Since γ_n is injective, $x' - \mu_n \circ \alpha_n(x') = x'' - \mu_n \circ \alpha_n(x'')$ which shows that β_n is well defined.

It is straightforward that $\nu_n \circ \beta_n = \text{identity}$ and the following diagram is commutative

$$\begin{array}{ccc} W_n & \xrightarrow{\phi_n} & W_{n+1} \\ \downarrow \beta_n & & \downarrow \beta_{n+1} \\ V_n & \xrightarrow{\varphi_n} & V_{n+1} \end{array}$$

ii) Let $\beta : \mathcal{W} \rightarrow \mathcal{V}$ with $\nu \circ \beta = \text{identity}$. For any $x \in V_n$, let $y_x = x - \beta_n \circ \nu_n(x)$, then $\nu_n(y_x) = 0$ and there exist a unique $z_x \in \mu_n$ such that $\mu_n(z_x) = y_x$.

Define

$$\begin{aligned} \alpha_n : V_n &\rightarrow U_n \\ x &\mapsto z_x \end{aligned}$$

α_n is well defined since for each x , y_x and z_x is unique.

It is straightforward to check $\alpha_n \circ \mu_n = \text{identity}$ and the following diagram is commutative

$$\begin{array}{ccc} V_n & \xrightarrow{\varphi_n} & V_{n+1} \\ \downarrow \alpha_n & & \downarrow \alpha_{n+1} \\ U_n & \xrightarrow{\psi_n} & U_{n+1}. \end{array}$$

4) Direct sum of a finite collection of persistence vector spaces $\mathcal{V}^i = \{V_n^i, \varphi_n^i : V_n^i \rightarrow V_{n+1}^i \mid n \in \mathbb{Z}_{\geq 0}\}$, $i \in \Lambda$ (Λ finite), is defined by $\bigoplus_{i \in \Lambda} \mathcal{V}^i = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$, where $V_n = \bigoplus_{i \in \Lambda} V_n^i$, $\varphi_n = \bigoplus_{i \in \Lambda} \varphi_n^i$ for $n \geq 0$.

Note that a direct sum of a finite collection of tame persistence vector spaces is tame.

Observation 3.1. 1) Given a split short exact sequence

$$0 \longrightarrow \mathcal{U} \xrightarrow{\mu} \mathcal{V} \xrightarrow{\nu} \mathcal{W} \longrightarrow 0,$$

we have $\mathcal{V} \cong \mathcal{U} \oplus \mathcal{W}$, where $\mathcal{U}, \mathcal{V}, \mathcal{W}$ and μ, ν are the same as in Definition 3.2 3).

2) Not any short exact sequence is split.

Proof. 1) Since the two conditions for a short exact sequence to be split are equivalent, WLOG, we can assume that there exists a linear map $\beta : \mathcal{W} \rightarrow \mathcal{V}$ such that $\nu \circ \beta = \text{identity}$. Define the linear map

$$\begin{aligned} \delta_n : U_n \oplus W_n &\rightarrow V_n \\ (y, z) &\mapsto \mu_n y + \beta_n z \end{aligned}$$

for all $n \geq 0$.

It is easy to check the following diagram is commutative

$$\begin{array}{ccc} U_n \oplus W_n & \xrightarrow{(\psi_n, \phi_n)} & U_{n+1} \oplus W_{n+1} \\ \downarrow \delta_n & & \downarrow \delta_{n+1} \\ V_n & \xrightarrow{\varphi_n} & V_{n+1} \end{array}$$

Given any element $x \in V_n$. Let $z = \nu_n(x)$, then $\nu_n(x - \beta_n z) = z - z = 0$ and there exist a unique $y \in U_n$ such that $\mu_n(y) = x - \beta_n z$.

Then $\delta_n(y, z) = \mu_n y + \beta_n z = x - \beta_n z + \beta_n z = x$ and δ_n is surjective.

Let $(y, z) \in \ker \delta_n$, then $\mu_n y + \beta_n z = 0$.

Since $0 = \nu_n(\mu_n y + \beta_n z) = \nu_n \circ \mu_n(y) + \nu_n \circ \beta_n(z) = 0 + z = z$, we have $z = 0$.

Then $\mu_n y = 0$ implies $y = 0$, since μ_n is injective.

Therefore $(y, z) = (0, 0)$ and δ_n is injective.

Hence δ_n is an isomorphism for all $n \geq 0$.

Therefore $\delta : \mathcal{U} \oplus \mathcal{W} \rightarrow \mathcal{V}$ is an isomorphism and $\mathcal{V} \cong \mathcal{U} \oplus \mathcal{W}$.

2) Counterexample:

Consider the following short exact sequence:

$$\begin{array}{ccccccc}
 \mathcal{U} & & \mathbb{R} & \xrightarrow{0} & \mathbb{R} & \xrightarrow{0} & 0 \xrightarrow{0} 0 \xrightarrow{0} 0 \xrightarrow{0} \dots \\
 \downarrow \mu \nearrow \alpha & & \downarrow i_1 & & \downarrow i_1 & & \downarrow 0 \\
 \mathcal{V} & & \mathbb{R} \oplus \mathbb{R} & \xrightarrow{c} & \mathbb{R} \oplus \mathbb{R} & \xrightarrow{0} & 0 \xrightarrow{0} 0 \xrightarrow{0} 0 \xrightarrow{0} \dots \\
 \downarrow \nu & & \downarrow p_2 & & \downarrow 0 & & \downarrow 0 \\
 \mathcal{W} & & \mathbb{R} & \xrightarrow{0} & \mathbb{R} & \xrightarrow{0} & 0 \xrightarrow{0} 0 \xrightarrow{0} 0 \xrightarrow{0} \dots
 \end{array}$$

where

$$\begin{aligned}
 i_1 : \mathbb{R} &\rightarrow \mathbb{R} \oplus \mathbb{R}, & p_2 : \mathbb{R} \oplus \mathbb{R} &\rightarrow \mathbb{R}, & c : \mathbb{R} \oplus \mathbb{R} &\rightarrow \mathbb{R} \oplus \mathbb{R} \\
 x &\mapsto (x, 0) & (x, y) &\mapsto y & (x, y) &\mapsto (y, 0)
 \end{aligned}$$

If the sequence splits, in view of $\alpha \circ \mu = \text{identity}$, $\alpha_1(x, 0) = x$, which makes the commutativity of the diagram

$$\begin{array}{ccc}
 \mathbb{R} & \xrightarrow{0} & \mathbb{R} \\
 \alpha_0 \uparrow & & \uparrow \alpha_1 \\
 \mathbb{R} \oplus \mathbb{R} & \xrightarrow{c} & \mathbb{R} \oplus \mathbb{R}
 \end{array}$$

impossible.

□

Notation 3.3. Define the following tame persistence vector spaces as **basic tame persistence vector spaces**.

1) $\kappa[t]$ is the tame persistence vector space over a field κ

$$\{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$$

where $V_n = \kappa$ and $\varphi_n = \text{identity}$ for all $n \geq 0$.

It corresponds to the interval $[0, \infty)$.

2) $S^r \kappa[t]$ is the tame persistence vector space over a field κ

$$\{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$$

where $V_n = 0$ for $0 \leq n < r$, $V_n = \kappa$ for $n \geq r$, $\varphi_n = 0$ for $0 \leq n < r$ and $\varphi_n = \text{identity}$ for $n \geq r$.

It corresponds to the interval $[r, \infty)$.

The notation S^r indicates the right shift with r -units.

3) $T_{r+1} \kappa[t]$ is the tame persistence vector space over a field κ

$$\{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$$

where $V_n = \kappa$ for $0 \leq n \leq r$, $V_n = 0$ for $n > r$, $\varphi_n = \text{identity}$ for $0 \leq n < r$ and $\varphi_n = 0$ for $n \geq r$.

It corresponds to the interval $[0, r]$.

The notation T_{r+1} indicates the truncation at level $r + 1$.

4) $S^r T_{p+1} \kappa[t]$ is the tame persistence vector space over a field κ

$$\{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$$

where $V_n = \kappa$ for $r \leq n \leq r + p$ and $V_n = 0$ otherwise; $\varphi_n = \text{identity}$ for $r \leq n \leq (r + p - 1)$ and $\varphi_n = 0$ otherwise.

It corresponds to the interval $[r, r + p]$.

Notation 3.4. For tame persistence vector space $\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$

1) Denote $\varphi_{i,j} = \varphi_{j-1} \circ \cdots \circ \varphi_i : V_i \rightarrow V_j$ for $i < j$ and $\varphi_{i,i} = \text{identity} : V_i \rightarrow V_i$, with $i, j \in \mathbb{Z}_{\geq 0}$.

2) Denote $\beta(i, j) = \dim(\text{im}(\varphi_{i,j} : V_i \rightarrow V_j))$ with $i \leq j \in \mathbb{Z}_{\geq 0}$.

Note $\beta(i, i) = \dim V_i$ and $\beta(i, j) = \beta(i, j + 1)$ for j large enough. If φ_n is an isomorphism for $n \geq N$, denote $\beta(i, \infty) = \beta(i, m)$, where m is any integer larger than i and N .

Definition 3.5. Define an order \prec of all basic tame persistence vector spaces:

- 1) $S^r \kappa[t] \prec S^{r'} \kappa[t]$ if $r < r'$;
- 2) $S^r T_{p+1} \kappa[t] \prec S^{r'} T_{p'+1} \kappa[t]$ if $r < r'$ or ($r = r'$ and $p < p'$);
- 3) $S^r T_{p+1} \kappa[t] \prec S^{r'} \kappa[t]$.

Clearly, this order is a strict total order.

Lemma 3.2. If \mathcal{V} and \mathcal{W} are two basic tame persistence vector spaces such that $\mathcal{V} \prec \mathcal{W}$, then any map from \mathcal{V} to \mathcal{W} is trivial.

Proof. We check situation 2) of Definition 3.5 first.

Suppose $\mathcal{V} = S^r T_{p+1} \kappa[t]$ and $\mathcal{W} = S^{r'} T_{p'+1} \kappa[t]$ and there is a linear map $\omega : \mathcal{V} \rightarrow \mathcal{W}$. We want to show $\omega = 0$.

We have 2 cases:

Case 1: $r < r'$

In this case, $\omega_n = 0$ for all $n < r$ or $n > r + p$, since $V_n = 0$.

For $r \leq n \leq r + p$, consider the following commutative diagram:

$$\begin{array}{ccc} \kappa & \xrightarrow{id} & \kappa \\ \downarrow \omega_r & & \downarrow \omega_n \\ 0 & \xrightarrow{0} & W_n \end{array}$$

We have $\omega_n = \omega_n \circ id = 0 \circ \omega_r = 0$.

Hence, $\omega = 0$, if $r < r'$.

Case 2: $r = r'$ and $p < p'$

Consider the commutative diagram:

$$\begin{array}{ccc} \kappa & \xrightarrow{0} & 0 \\ \downarrow \omega_n & & \downarrow \omega_{r'+p'} \\ \kappa & \xrightarrow{id} & \kappa \end{array}$$

for $r \leq n \leq r + p$.

The proofs of situations 1) and 3) of Definition 3.5 are similar.

□

Proposition 3.3. *Any tame persistence vector space $\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} | n \in \mathbb{Z}_{\geq 0}\}$ over a field κ is isomorphic to*

$$\bigoplus_{1 \leq i \leq p} S^{r_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq q} S^{m_j} T_{n_j+1} \kappa[t]$$

where p, q, r_i, m_j and $n_j \in \mathbb{Z}_{\geq 0}$.

So, any tame persistence vector space can be decomposed in to a direct sum of a finite collection of basic tame persistence vector spaces.

Proof. Since \mathcal{V} is a tame persistence vector space $\dim(V_n)$ is finite for all $n \geq 0$ and there exists $N \geq 0$ such that φ_n is an isomorphisms for $n \geq N$.

Define $L(\mathcal{V}) = \sum_{0 \leq n \leq N} \dim V_n$.

We prove by induction on $L(\mathcal{V})$.

If $L(\mathcal{V}) = 0$, then $\mathcal{V} = 0$. We are done.

If $L(\mathcal{V}) \neq 0$, then $V_k \neq 0$ for some $0 \leq k \leq N$, with $V_i = 0$ if $i < k$.

WLOG, suppose $V_0 \neq 0$ and choose a nonzero element v_0 from V_0 .

Define $v_n = \varphi_{0,n}(v_0)$, $n \geq 0$.

Case 1: $v_n \neq 0$ for $0 \leq n < j$ and $v_n = 0$ for $n \geq j$, where $j \geq 1$.

Note. We must have $j \leq N$, otherwise $v_n \neq 0$ for all $n \geq 0$, a contradiction.

Consider the following “short exact sequence” of tame persistence vector

spaces

$$\begin{array}{ccccccccccc}
 \mathcal{W}_0 & \kappa & \xrightarrow{\alpha_0} & \kappa & \xrightarrow{\alpha_1} & \cdots & \xrightarrow{\alpha_{j-2}} & \kappa & \xrightarrow{\alpha_{j-1}} & 0 & \xrightarrow{\alpha_j} & 0 & \xrightarrow{\alpha_{j+1}} & \cdots \\
 & \downarrow f_0 & \lrcorner h_0 & \downarrow f_1 & \lrcorner h_1 & & & \downarrow f_{j-1} & \lrcorner h_{j-1} & \downarrow f_j & \lrcorner h_j & \downarrow f_{j+1} & \lrcorner h_{j+1} & \\
 \mathcal{V} & V_0 & \xrightarrow{\varphi_0} & V_1 & \xrightarrow{\varphi_1} & \cdots & \xrightarrow{\varphi_{j-2}} & V_{j-1} & \xrightarrow{\varphi_{j-1}} & V_j & \xrightarrow{\varphi_j} & V_{j+1} & \xrightarrow{\varphi_{j+1}} & \cdots \\
 & \downarrow g_0 & & \downarrow g_1 & & & & \downarrow g_{j-1} & & \downarrow g_j & & \downarrow g_{j+1} & & \\
 \mathcal{V}' & V_0/\kappa v_0 & \xrightarrow{\tilde{\varphi}_0} & V_1/\kappa v_1 & \xrightarrow{\tilde{\varphi}_1} & \cdots & \xrightarrow{\tilde{\varphi}_{j-2}} & V_{j-1}/\kappa v_{j-1} & \xrightarrow{\tilde{\varphi}_{j-1}} & V_j & \xrightarrow{\tilde{\varphi}_j} & V_{j+1} & \xrightarrow{\tilde{\varphi}_{j+1}} & \cdots
 \end{array}$$

where

$\alpha_n = \text{identity}$ for $0 \leq n \leq j-2$ and $\alpha_n = 0$ for $n \geq j-1$;

$f_n : \kappa \rightarrow V_n$ for $0 \leq n \leq j-1$

$\lambda \mapsto \lambda v_n$

and $f_n = 0$ for $n \geq j$;

$g_n : V_n \rightarrow V_n/\kappa v_n$ for $0 \leq n \leq j-1$

$x \mapsto x + \kappa v_n$

and $g_n = \text{identity}$ for $n \geq j$;

$\tilde{\varphi}_n : V_n/\kappa v_n \rightarrow V_{n+1}/\kappa v_{n+1}$ for $0 \leq n \leq j-1$

$x + \kappa v_n \mapsto \varphi_n(x) + \kappa v_{n+1}$

and $\tilde{\varphi}_n = \varphi_n$ for $n \geq j$.

The above diagram is commutative and is a short exact sequence of tame persistence vector spaces. Next we will show the above short exact sequence splits.

Extend v_{j-1} to a basis $\{x_1 = v_{j-1}, x_2, \dots, x_p\}$ of V_{j-1} .

Given any element $x \in V_{j-1}$, we have a unique expression

$$x = \sum_{i=1}^p a_i x_i$$

Define a linear map

$$h_{j-1} : V_{j-1} \rightarrow \kappa$$

$$\sum_{i=1}^p a_i x_i \mapsto a_1$$

For $0 \leq n \leq j-1$, define a linear map

$$h_n : V_n \rightarrow \kappa$$

$$x \mapsto h_{j-1} \circ \varphi_{n,j-1}(x)$$

For $n \geq j$, define $h_n = 0$.

Clearly $h : \mathcal{V} \rightarrow \mathcal{W}_0$ is well defined and $h \circ f = \text{identity}$.

Therefore the above short exact sequence splits, and

$$\begin{aligned}\mathcal{V} &\cong \mathcal{W}_0 \oplus \mathcal{V}' \\ &= T_j \kappa[t] \oplus \mathcal{V}'\end{aligned}$$

Hence $L(\mathcal{V}') = L(\mathcal{V}) - j < L(\mathcal{V})$.

Case 2: $v_n \neq 0$ for all $n \geq 0$.

Consider the following “short exact sequence” of tame persistence vector spaces

$$\begin{array}{ccccccc} \mathcal{W}_1 & \kappa & \xrightarrow{\alpha_0} \cdots \xrightarrow{\alpha_{N-2}} \kappa & \xrightarrow{\alpha_{N-1}} \kappa & \xrightarrow{\alpha_N} \kappa & \xrightarrow{\alpha_{N+1}} \cdots \\ & \downarrow f_0 \quad \uparrow h_0 & & \downarrow f_{N-1} \quad \uparrow h_{N-1} & \downarrow f_N \quad \uparrow h_N & \downarrow f_{N+1} \quad \uparrow h_{N+1} \\ \mathcal{V} & V_0 & \xrightarrow{\varphi_0} \cdots \xrightarrow{\varphi_{N-2}} V_{N-1} & \xrightarrow{\varphi_{N-1}} V_N & \xrightarrow{\varphi_N} V_{N+1} & \xrightarrow{\varphi_{N+1}} \cdots \\ & \downarrow g_0 & & \downarrow g_{N-1} & \downarrow g_N & \downarrow g_{N+1} \\ \mathcal{V}' & V_0/\kappa v_0 & \xrightarrow{\tilde{\varphi}_0} \cdots \xrightarrow{\tilde{\varphi}_{N-2}} V_{N-1}/\kappa v_{N-1} & \xrightarrow{\tilde{\varphi}_{N-1}} V_N/\kappa v_N & \xrightarrow{\tilde{\varphi}_N} V_{N+1}/\kappa v_{N+1} & \xrightarrow{\tilde{\varphi}_{N+1}} \cdots \end{array}$$

where $\alpha_n = \text{identity}$, $f_n : \kappa \rightarrow V_n$, $g_n : V_n \rightarrow V_n/\kappa v_n$
 $\lambda \mapsto \lambda v_n$ $x \mapsto x + \kappa v_n$

$$\begin{aligned}\tilde{\varphi}_n : V_n/\kappa v_n &\rightarrow V_{n+1}/\kappa v_{n+1} \\ x + \kappa v_n &\mapsto \varphi_n(x) + \kappa v_{n+1}\end{aligned}$$

for all $n \geq 0$.

We will show the above short exact sequence splits.

Extend v_N to a basis $\{y_1 = v_N, y_2, \dots, y_q\}$ of V_N .

Given any element $y \in V_N$, we have a unique expression

$$y = \sum_{j=1}^q b_j y_j$$

Define a linear map

$$\begin{aligned}h_N : V_N &\rightarrow \kappa \\ \sum_{j=1}^q b_j y_j &\mapsto b_1\end{aligned}$$

For $0 \leq n \leq N$, define a linear map

$$\begin{aligned}h_n : V_n &\rightarrow \kappa \\ x &\mapsto h_N \circ \varphi_{n,N}(x)\end{aligned}$$

For $n > N$, define a linear map

$$\begin{aligned}h_n : V_n &\rightarrow \kappa \\ x &\mapsto h_N \circ \varphi_{N,n}^{-1}(x)\end{aligned}$$

Clearly $h_n \circ f_n = \text{identity}$ and

$$\begin{array}{ccc} \kappa & \xrightarrow{\alpha_n} & \kappa \\ \uparrow h_n & & \uparrow h_{n+1} \\ V_n & \xrightarrow{\varphi_n} & V_{n+1} \end{array}$$

is commutative for all $n \geq 0$.

Therefore the short exact sequence in Case 2 splits, and

$$\begin{aligned} \mathcal{V} &\cong \mathcal{W}_1 \oplus \mathcal{V}' \\ &= \kappa[t] \oplus \mathcal{V}' \end{aligned}$$

Hence $L(\mathcal{V}') = L(\mathcal{V}) - (N + 1) < L(\mathcal{V})$.

In both cases $L(\mathcal{V}') < L(\mathcal{V})$.

By induction on $L(\mathcal{V})$,

$$\bigoplus_{1 \leq i \leq s} S^{r_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq t} S^{m_j} T_{n_j+1} \kappa[t]$$

□

Proposition 3.4. *Let*

$$\mathcal{V} = \bigoplus_{1 \leq i \leq p} S^{r_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq q} S^{m_j} T_{n_j+1} \kappa[t]$$

and

$$\mathcal{V}' = \bigoplus_{1 \leq i \leq p'} S^{r'_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq q'} S^{m'_j} T_{n'_j+1} \kappa[t]$$

.

If $\mathcal{V} \cong \mathcal{V}'$, then $p = p'$, $q = q'$ and $r_i = r'_i$, $m_j = m'_j$, $n_j = n'_j$ after a suitable permutation.

Proof. Reorder components in \mathcal{V} and \mathcal{V}' in increasing order (See Definition 3.5) and group all copies of the same basic tame persistence vector space together into isotypical components, so we have

$$\mathcal{W} = \bigoplus_{1 \leq n \leq a} W_n$$

and

$$\mathcal{W}' = \bigoplus_{1 \leq n \leq b} W'_n$$

Precisely each isotypical component of \mathcal{W} or \mathcal{W}' is a direct sum of isomorphic basic tame persistence vector spaces.

We only need to prove that $a = b$ and $W_n = W'_n$ for $1 \leq n \leq a = b$.

Define $S(\mathcal{W}) = a$, the cardinality of the isotypical components of \mathcal{W} . We will prove the statement by induction on $S(\mathcal{W})$.

If $S(\mathcal{W}) = 0$, clearly $\mathcal{W} = 0 = \mathcal{W}'$.

If $S(\mathcal{W}) > 0$, write $\mathcal{W} = W_1 \oplus R$ and $\mathcal{W}' = P_1 \oplus R'$,

where $R = \oplus_{2 \leq n \leq a} W_n$, $P_1 = \oplus_{1 \leq n \leq b_1} W'_n$, $R' = \oplus_{b_1+1 \leq n \leq b} W'_n$,

b_1 is an integer such that $W'_n \leq W_1$ for $1 \leq n \leq b_1$ and $W'_n > W_1$ for $n > b_1$.

Here the order of W_n and W'_n is determined by the order of their basic components.

Since $\mathcal{W} \cong \mathcal{W}'$, there is a pair of isomorphisms $\omega : \mathcal{W} \rightarrow \mathcal{W}'$ and $\omega' : \mathcal{W}' \rightarrow \mathcal{W}$ such that $\omega \circ \omega' = id$ and $\omega' \circ \omega = id$.

Write $\omega : \mathcal{W} \rightarrow \mathcal{W}'$ as a matrix

$$\begin{matrix} & W_1 & R \\ P_1 & \begin{pmatrix} A & C \end{pmatrix} \\ R' & \begin{pmatrix} B & D \end{pmatrix} \end{matrix}$$

Since any component of $W_1 <$ any component of R' , by Lemma 3.2, $B = 0$.

So matrix form of ω is

$$\begin{matrix} & W_1 & R \\ P_1 & \begin{pmatrix} A & C \end{pmatrix} \\ R' & \begin{pmatrix} 0 & D \end{pmatrix} \end{matrix}$$

Similarly, ω' has matrix form

$$\begin{matrix} & P_1 & R' \\ W_1 & \begin{pmatrix} A' & C' \end{pmatrix} \\ R & \begin{pmatrix} 0 & D' \end{pmatrix} \end{matrix}$$

Since $\omega \circ \omega' = id$ and $\omega' \circ \omega = id$, we have

$$\begin{pmatrix} A & C \\ 0 & D \end{pmatrix} \begin{pmatrix} A' & C' \\ 0 & D' \end{pmatrix} = I \text{ and } \begin{pmatrix} A' & C' \\ 0 & D' \end{pmatrix} \begin{pmatrix} A & C \\ 0 & D \end{pmatrix} = I.$$

So $AA' = I$, $A'A = I$, $DD' = I$, $D'D = I$.

Then $W_1 \cong P_1$ and $R \cong R'$.

Since W_1 and P_1 are isomorphic, each basic component of P_1 must be isomorphic to the basic component of W_1 . Their number in W_1 and P_1 should be the same.

So $W_1 = P_1$.

Since $R \cong R'$ and $S(R) = a - 1 < a$, by induction we finish the proof. \square

Definition 3.6. *Bar codes is a finite collection of intervals*

$$[i, j]$$

with $i \in \mathbb{Z}_{\geq 0}, j \in \mathbb{Z}_{\geq 0} \cup \{\infty\}$ and $i \leq j$.

Given a tame persistence vector space \mathcal{V} , there exist a decomposition

$$\mathcal{V} \cong \bigoplus_{1 \leq i \leq p} S^{r_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq q} S^{m_j} T_{n_j+1} \kappa[t]$$

by Proposition 3.3.

Then assign bar code

$$\mathcal{B}(\mathcal{V}) = \{[r_i, \infty], [m_j, m_j + n_j] | 1 \leq i \leq p, 1 \leq j \leq q\}$$

to \mathcal{V} . Call it the bar code of the tame persistence vector space \mathcal{V} .

This bar code $\mathcal{B}(\mathcal{V})$ is unique by Proposition 3.4.

Theorem 3.5. *Two tame persistence vector spaces are isomorphic iff their bar codes are the same.*

Proof. Theorem 3.5 is obtained directly from Proposition 3.3 and Proposition 3.4. \square

Observation 3.6. $\beta(i, j)$ = number of intervals in $\mathcal{B}(\mathcal{V})$ which contain $[i, j]$ for $i \in \mathbb{Z}_{\geq 0}, j \in \mathbb{Z}_{\geq 0} \cup \{\infty\}$ and $i \leq j$. In particular, $\dim(V_i)$ = number of intervals in $\mathcal{B}(\mathcal{V})$ which contain $\{i\}$ for $i \geq 0$.

Proof. Suppose

$$\mathcal{V} \cong \bigoplus_{1 \leq i \leq p} S^{r_i} \kappa[t] \oplus \bigoplus_{1 \leq j \leq q} S^{m_j} T_{n_j+1} \kappa[t] = \mathcal{W}.$$

Since $\mathcal{V} \cong \mathcal{W}$, there exist an isomorphism $f : \mathcal{V} \rightarrow \mathcal{W}$ and following commutative diagram:

$$\begin{array}{ccccccc} V_0 & \xrightarrow{\varphi_0} & V_1 & \xrightarrow{\varphi_1} & \dots & \xrightarrow{\varphi_{n-1}} & V_n & \xrightarrow{\varphi_n} & V_{n+1} & \xrightarrow{\varphi_{n+1}} & \dots \\ \downarrow f_0 & & \downarrow f_1 & & & & \downarrow f_n & & \downarrow f_{n+1} & & \\ W_0 & \xrightarrow{\phi_0} & W_1 & \xrightarrow{\phi_1} & \dots & \xrightarrow{\phi_{n-1}} & W_n & \xrightarrow{\phi_n} & W_{n+1} & \xrightarrow{\phi_{n+1}} & \dots \end{array}$$

Since $f_j : \text{im}(\varphi_{i,j} : V_i \rightarrow V_j) \rightarrow \text{im}(\phi_{i,j} : W_i \rightarrow W_j)$ is an isomorphism, $\beta(i, j) = \dim(\text{im}(\phi_{i,j} : W_i \rightarrow W_j))$.

Observe that $\phi_{i,j}$ is a direct sum of linear maps $\kappa \xrightarrow{id} \kappa$ or $0 \longrightarrow \kappa$ or $\kappa \longrightarrow 0$. $\beta(i, j) = \dim(\text{im} \phi_{i,j})$ is the number of linear maps $\kappa \xrightarrow{id} \kappa$. Each $\kappa \xrightarrow{id} \kappa$ corresponds to an interval in the bar code $\mathcal{B}(\mathcal{W})$ that contains $[i, j]$.

Therefore $\beta(i, j)$ is the number of intervals in bar code $\mathcal{B}(\mathcal{V})$ that contains $[i, j]$. □

Definition 3.7. Given a tame persistence vector space \mathcal{V} , define $\mu(i, j) =$ number of intervals in $\mathcal{B}(\mathcal{V})$ which equal to $[i, j]$.

Observation 3.7. 1) [19]

$$\begin{aligned} & \mu(i, j) \\ = & \begin{cases} \beta(i, j) - \beta(i-1, j) - \beta(i, j+1) + \beta(i-1, j+1), & 0 < i \leq j < \infty \\ \beta(0, j) - \beta(0, j+1), & i = 0, 0 \leq j < \infty \\ \beta(i, \infty) - \beta(i-1, \infty), & 0 < i < \infty, j = \infty \\ \beta(0, \infty), & i = 0, j = \infty \end{cases} \end{aligned}$$

$$2) \beta(i, j) = \sum_{l \leq i, m \geq j} \mu(l, m)$$

Proof. 1) In the case $0 < i \leq j < \infty$, we have

$$\begin{aligned} \mu(i, j) &= \#\{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau = [i, j]\} \\ &= \#\{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i, j]\} - \#\{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i-1, j]\} \\ &\quad - \#\{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i, j+1]\} + \#\{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i-1, j+1]\} \\ &= \beta(i, j) - \beta(i-1, j) - \beta(i, j+1) + \beta(i-1, j+1) \end{aligned}$$

The second identity in the above identities holds because

$$\begin{aligned} & \{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau = [i, j]\} \\ = & \{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i, j]\} - \{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i-1, j]\} \\ & - \{\tau \in \mathcal{B}(\mathcal{V}) \mid \tau \supseteq [i, j+1]\} \end{aligned}$$

The other three cases are easier to prove.

2) Follows directly from definition. □

Definition 3.8. Denote $k(i, j) = \dim(\ker(\varphi_{i,j} : V_i \rightarrow V_j))$, $k(i) = \dim(V_i)$ with $i \leq j \in \mathbb{Z}_{\geq 0}$.

Observation 3.8. 1)

$$k(i, j) = \beta(i, i) - \beta(i, j);$$

2)

$$\beta(i, j) = k(i) - k(i, j).$$

From the above observations, we see that the set of numbers $\{\mu(i, j)\}$, $\{\beta(i, j)\}$ and $\{k(i, j), k(i)\}$ are equivalent. Sometimes, it is more convenient to define and calculate bar codes using the numbers $k(i, j)$ and $k(i)$ instead of $\beta(i, j)$.

Edelsbrunner-Letscher-Zomorodian's Interpretation [20]

Let \mathcal{V} be a tame persistence vector space (see Definition 3.2).

One says that the nonzero element $x \in V_j$ is **born** in $V_i, i \leq j$ hence “is born at time i ” if it is in the image of $\varphi_{i,j}$ and not in the image of $\varphi_{i-1,j}$ and **dies** in $V_k, k \geq j+1$, hence “dies at time j ” if $\varphi_{j,k}(x) = 0$ but $\varphi_{j,k-1} \neq 0$.

A subset $S \subseteq V$ (V is a vector space) is called linearly independent if its elements form a collection of linearly independent vectors in V . If $S_i \subseteq V_i$ is linearly independent in V_i and $\varphi_{i,j}(S_i)$ is linearly independent in V_j , then we call S_i linearly independent on interval $[i, j]$.

With this in mind, $\beta(i, j)$ is the maximal cardinality of linearly independent sets on the interval $[i, j]$ and $\mu(i, j)$ = the maximal cardinality of linearly independent sets on the interval $[i, j]$, whose elements are born in V_i and die in V_{j+1} .

Each interval $[i, j]$ in $\mathcal{B}(\mathcal{V})$ guarantee the existence of an element in V_i , which is born in V_i , survives in each $V_r (i \leq r \leq j)$ and dies in V_{j+1} .

In [20] the authors propose to collect $\mu(i, j)$ as points in the extended half space $\overline{HR} := \{(x, y) \in \mathbb{R}^2 \sqcup R \times \infty \mid x \leq y\}$. With this convention Cohen-Steiner, Edelsbrunner, Harer [9] have established a strong stability result associated with a real valued map (cf. Stability Theorem, page 182 in [19]) which relates the distance between two real valued functions and the bottle neck distance of their persistence diagrams. This concept will not be used in the sequel.

Zomorodian - Carlsson's Interpretation [38]

Definition 3.9. 1) A $\kappa[t]$ -module V is the vector space over κ equipped with a linear map $A : V \rightarrow V$. The module action is defined as

$$\begin{aligned} \kappa[t] \times V &\rightarrow V \\ (\sum_{i=0}^n a_i t^i, v) &\mapsto \sum_{i=0}^n a_i A^i(v) \end{aligned}$$

2) A $\kappa[t]$ -module V with a linear map $A : V \rightarrow V$ is finitely generated if there exists $\{v_1, v_2, \dots, v_r\}$ such that any $v \in V$ is a linear combination of

$$\{v_1, v_2, \dots, v_r, A(v_1), A(v_2), \dots, A(v_r), A^2(v_1), A^2(v_2), \dots, A^2(v_r), \dots\}.$$

3) A graded $\kappa[t]$ -module V is a $\kappa[t]$ -module together with a decomposition of the vector space $V = \bigoplus_{n \geq 0} V_n$ and a linear map $A : V \rightarrow V$ such that $A(v_n) \in V_{n+1}, \forall v_n \in V_n$.

4) Let $V = \bigoplus_{n \geq 0} V_n$ with the linear map $A : V \rightarrow V$ and $W = \bigoplus_{n \geq 0} W_n$ with the linear map $B : W \rightarrow W$ be two graded $\kappa[t]$ -modules.

A morphism of graded $\kappa[t]$ -modules $f : V \rightarrow W$ is a linear map of vector spaces such that $f(V_n) \subseteq W_n$ and $f \circ A = B \circ f$.

A morphism of graded $\kappa[t]$ -module $f : V \rightarrow W$ is an isomorphism if there exists another morphism of graded $\kappa[t]$ -module $g : W \rightarrow V$ such that $g \circ f : V \rightarrow V$ and $f \circ g : W \rightarrow W$ are identities. V and W are isomorphic if there is an isomorphism between them.

One of the main results of [38] is the equivalence of the category of finitely generated $\kappa[t]$ -modules and the category of tame persistence vector spaces.

Proposition 3.9. *Finitely generated graded $\kappa[t]$ -modules identify to tame persistence vector spaces and so do their morphisms.*

Proof. 1) Let $V = \bigoplus_{n \geq 0} V_n$ be a finitely generated $\kappa[t]$ -module with a linear map $A : V \rightarrow V$ such that $A(v_n) \in V_{n+1}, \forall v_n \in V_n$.

There exist $\{v_1, v_2, \dots, v_r\}$ such that any $v \in V$ is a linear combination of

$$(*) \quad \{v_1, v_2, \dots, v_r, A(v_1), A(v_2), \dots, A(v_r), A^2(v_1), A^2(v_2), \dots, A^2(v_r), \dots\}.$$

Since each v_i is a sum of homogeneous components, WLOG, we can assume $\{v_1, v_2, \dots, v_r\}$ themselves are homogeneous, i.e., $v_i \in V_{n_i}, 1 \leq i \leq r$.

Let n_i be the degree of $v_i, 1 \leq i \leq r$. WLOG, we can assume $n_i \leq n_{i+1}$.

Let $\varphi_n = A|_{V_n} : V_n \rightarrow V_{n+1}$, then

$$\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} | n \in \mathbb{Z}_{\geq 0}\}$$

is a persistence vector space.

To see that \mathcal{V} is tame we denote by $N_0 = \max_{1 \leq i \leq r} n_i$. Clearly $\dim V_n (n \leq N_0)$ is finite and in view of $(*)$ the linear maps $\varphi_{n,n+k} = \varphi_{n+k-1} \circ \varphi_{n+k-2} \circ \dots \circ \varphi_n$

are surjective hence $\dim V_n$ is finite for any $n > N_0$. Since φ_n is surjective for any $n > N_0$ and $\dim V_{N_0}$ is finite, there exists N so that $\dim V_n$ is constant for $n > N$. Since any surjective map between vector spaces of the same finite dimension is an isomorphism, the linear map φ_n is an isomorphism for $n > N$.

Hence, \mathcal{V} is a tame persistence vector space.

If $\mathcal{V} = \{V_n, \varphi_n : V_n \rightarrow V_{n+1} \mid n \in \mathbb{Z}_{\geq 0}\}$ is a tame persistence vector space, define the graded $\kappa[t]$ -module $V = \bigoplus_{n \geq 0} V_n$ with linear map $A = \bigoplus_{n \geq 0} \varphi_n : V \rightarrow V$.

There exist $N \geq 0$ such that φ_n is an isomorphism for $n \geq 0$. Let $\{v_1, v_2, \dots, v_r\}$ be the set of all generators of V_1, V_2, \dots, V_N , then any $v \in V$ is a linear combination of

$$\{v_1, v_2, \dots, v_r, A(v_1), A(v_2), \dots, A(v_r), A^2(v_1), A^2(v_2), \dots, A^2(v_r), \dots\}.$$

2) From definitions of morphisms of finitely generated graded $\kappa[t]$ -modules and linear maps of tame persistence vector spaces, we can see that two finitely graded $\kappa[t]$ -modules are isomorphic iff the tame persistence vector spaces associated with them are isomorphic.

□

Recall that the ring $\kappa[t]$ is a principal ideal domain and a basic theorem in algebra [27] claims that any finitely generated modules over a principal ideal domain decompose uniquely. In particular, any finitely generated modules over $\kappa[t]$ decomposes uniquely as a finite direct sum of free modules $\kappa[t]$'s and torsion modules $T_{d_i} \kappa[t]$'s. As noticed by [38], the above result extends to finitely generated graded modules where the free module $\kappa[t]$ is to be replaced by $S^r \kappa[t]$ for some r and torsion module $T_d \kappa[t]$ by $S^r T_d \kappa[t]$ for some r and d . Notice that the module $S^r A$ has the component $(S^r A)_p = A_{p-r}$ for $p \geq r$ and equal to zero for $p < r$. Note that each component $S^r \kappa[t]$ corresponds to a bar code $[r, \infty)$ and each component $S^r T_{d+1} \kappa[t]$ corresponds to a bar code $[r, r+d]$ (cf. [38]).

Quiver Representation Perspective

As noticed by G. Carlsson and Vin de Silva [6], persistence vector spaces which stabilize for $n \geq N$ can be regarded as representation of the oriented graph

$$\bullet_1 \longrightarrow \bullet_2 \longrightarrow \bullet_3 \longrightarrow \cdots \longrightarrow \bullet_{N-1} \longrightarrow \bullet_N$$

Any such representation is a sum of indecomposable representations which are classified by the intervals $[i, j]$, $1 \leq i \leq j \leq N$. The interval $[i, j]$ with $j < N$ correspond to bar code $[i, j]$ while the interval $[i, N]$ to the bar code $[i, \infty)$. The interval $[i, j]$ corresponds to the representation

$$0_1 \longrightarrow \cdots \longrightarrow 0_{i-1} \longrightarrow \kappa_i \xrightarrow{id} \kappa_{i+1} \xrightarrow{id} \cdots \xrightarrow{id} \kappa_j \longrightarrow 0_{j+1} \longrightarrow \cdots$$

This is a result in the theory of quiver representation due to P. Gabriel[22].

3.4 Persistent Homology and Bar Codes of a Filtered Simplicial/Polytopal Complex

Definition 3.10. 1) A filtered space \mathcal{X} consists of a space X and a finite filtration

$$X_0 \subseteq X_1 \subseteq \cdots \subseteq X_N = X. \quad (*)$$

2) Sometimes it is convenient to suppose that X is embedded in X_∞ which is contractible. For a filtration without X_∞ , we can complete it with $X_\infty = C(X)$, the cone over X . Then $(*)$ becomes

$$X_0 \subseteq X_1 \subseteq \cdots \subseteq X_N = X \subseteq X_\infty. \quad (**)$$

Example 3.3. Filtrations of Vietoris-Rips complexes (3.1) - (3.3) associated to a PCD X with $N + 1$ points in \mathbb{R}^m provide examples of filtered simplicial complexes where $X_\infty = R_{\epsilon_N}(X) = \Delta^n$.

Below are two other relevant examples.

Definition 3.11. A continuous map $f : X \rightarrow \mathbb{R}$ is called **weakly tame** if X is compact and there exists finitely many values $\min(f(X)) = t_0 < t_1 < \cdots < t_N = \max(f(X))$ (so called **critical values**) so that:

(i) for any t the closed sub-level $X_{-\infty, t}$ is a deformation retract of an open neighborhood;

(ii) for any i and t , $t \in [t_i, t_{i+1})$, $X_{-\infty, t}$ retracts by deformation to $X_{-\infty, t_i}$.

Informally this means that each sub-level is homotopically well behaved (neighborhood retract) and the topology(homotopy type) of sub-levels change only for finitely many t 's.

Example 3.4. A weakly tame map $f : X \rightarrow \mathbb{R}$ with critical values $t_0 < t_1 < \cdots < t_N$ provides another example with $X_i = f^{-1}((-\infty, t_i])$, $0 \leq i \leq N$, $X_\infty = CX_N$, where CX_N is the cone with base X_N .

The properties (i) and (ii) are sufficient hypotheses to ensure that in each dimension the homology vector spaces of X_{t_i} provide *tame persistence vector spaces*.

Example 3.5. There is a natural filtration of an N -dimensional simplicial complex or polytopal complex, the skeleton filtration

$$X^0 \subseteq X^1 \subseteq \dots \subseteq X^N = X \subseteq X^\infty$$

where $X^i (0 \leq i \leq N)$ are i -skeletons of X and $X^\infty = CX$.

This subsection defines and studies persistent homology and bar codes of filtered spaces of finite simplicial/polytopal complexes.

Given a filtered space \mathcal{K} of a finite simplicial/polytopal complex K

$$K_0 \subseteq K_1 \subseteq \dots \subseteq K_N = K \quad (3.4)$$

we can associate a commutative diagram defined below.

Denote by C_r^s the κ -vector space $C_r(K_s)$ with basis the r -simplices/ r -cells of K_s , $\partial_r^s : C_r^s \rightarrow C_{r-1}^s$ the boundary map from $C_r(K_s)$ to $C_{r-1}(K_s)$, $i_r^s : C_r^s \rightarrow C_r^{s+1}$ the linear map induced by the inclusion from K_s to K_{s+1} (clearly i_r^s is one to one). Let $C_r^s = C_r^N$, $\partial_r^s = \partial_r^N$ and $i_r^s = \text{identity}$ when $s \geq N$, one obtains the commutative diagram

$$\begin{array}{ccccccc}
 \vdots & & \vdots & & \vdots & & \vdots \\
 \downarrow \partial_{M+1}^0 & & \downarrow \partial_{M+1}^s & & \downarrow \partial_{M+1}^N & & \downarrow \partial_{M+1}^{N+1} \\
 C_M^0 & \xrightarrow{i_M^0} \dots \xrightarrow{i_M^{s-1}} & C_M^s & \xrightarrow{i_M^s} \dots \xrightarrow{i_M^{N-1}} & C_M^N & \xrightarrow[i_M^N]{\cong} & C_M^{N+1} \xrightarrow[i_M^{N+1}]{\cong} \dots \\
 \downarrow \partial_M^0 & & \downarrow \partial_M^s & & \downarrow \partial_M^N & & \downarrow \partial_M^{N+1} \\
 \vdots & & \vdots & & \vdots & & \vdots \\
 \downarrow \partial_{r+1}^0 & & \downarrow \partial_{r+1}^s & & \downarrow \partial_{r+1}^N & & \downarrow \partial_{r+1}^{N+1} \\
 C_r^0 & \xrightarrow{i_r^0} \dots \xrightarrow{i_r^{s-1}} & C_r^s & \xrightarrow{i_r^s} \dots \xrightarrow{i_r^{N-1}} & C_r^N & \xrightarrow[i_r^N]{\cong} & C_r^{N+1} \xrightarrow[i_r^{N+1}]{\cong} \dots \\
 \downarrow \partial_r^0 & & \downarrow \partial_r^s & & \downarrow \partial_r^N & & \downarrow \partial_r^{N+1} \\
 C_{r-1}^0 & \xrightarrow{i_{r-1}^0} \dots \xrightarrow{i_{r-1}^{s-1}} & C_{r-1}^s & \xrightarrow{i_{r-1}^s} \dots \xrightarrow{i_{r-1}^{N-1}} & C_{r-1}^N & \xrightarrow[i_{r-1}^N]{\cong} & C_{r-1}^{N+1} \xrightarrow[i_{r-1}^{N+1}]{\cong} \dots \\
 \downarrow \partial_{r-1}^0 & & \downarrow \partial_{r-1}^s & & \downarrow \partial_{r-1}^N & & \downarrow \partial_{r-1}^{N+1} \\
 \vdots & & \vdots & & \vdots & & \vdots \\
 \downarrow \partial_1^0 & & \downarrow \partial_1^s & & \downarrow \partial_1^N & & \downarrow \partial_1^{N+1} \\
 C_0^0 & \xrightarrow{i_0^0} \dots \xrightarrow{i_0^{s-1}} & C_0^s & \xrightarrow{i_0^s} \dots \xrightarrow{i_0^{N-1}} & C_0^N & \xrightarrow[i_0^N]{\cong} & C_0^{N+1} \xrightarrow[i_0^{N+1}]{\cong} \dots \\
 \downarrow \partial_0^0 & & \downarrow \partial_0^s & & \downarrow \partial_0^N & & \downarrow \partial_0^{N+1} \\
 0 & & 0 & & 0 & & 0
 \end{array} \quad (3.5)$$

Note. Since K is finite, each row is a tame persistence vector space and $C_r^s = 0$ for $r > \dim(K)$.

Passing to homology with coefficient in a field κ , consider $\mathcal{H}_r^s = \ker(\partial_r^s) / \text{im}(\partial_{r+1}^s)$, and $\tilde{i}_r^s : \mathcal{H}_r^s \rightarrow \mathcal{H}_r^{s+1}$ the linear maps induced in homology by the linear inclusions i_r^s . For each $r(\geq 0)$

$$\mathcal{H}_r(\mathcal{K}) := \{\mathcal{H}_r^s, \tilde{i}_r^s : \mathcal{H}_r^s \rightarrow \mathcal{H}_r^{s+1} | s \in \mathbb{Z}_{\geq 0}\}$$

is a persistence vector space with bar code $\mathcal{B}(\mathcal{H}_r(\mathcal{K}))$.

Following [20], the collection of vector spaces $\mathcal{H}_r^{s,p} = \text{im}(\tilde{i}_r^{s,s+p} : \mathcal{H}_r^s \rightarrow \mathcal{H}_r^{s+p})$ is referred to as the persistent homology.

2) The collection of bar codes $\mathcal{B}(\mathcal{H}_r(\mathcal{K}))$ for all r is denoted by $\mathcal{B}(\mathcal{K})$ and referred to as the bar code of \mathcal{K} .

Suppose \mathcal{K} is a filtered simplicial / polytopal complex as in (3.4),

$$K_0 \subseteq K_1 \subseteq \cdots \subseteq K_N = K,$$

we calculate the bar codes of \mathcal{K} in two cases: $\kappa = \mathbb{Z}_2$ and $\kappa = \mathbb{R}$. We consider only filtered simplicial complex below, filtered polytopal complex can be dealt with similarly.

Case A: $\kappa = \mathbb{Z}_2$

In this case, we can calculate the bar codes of \mathcal{K} using the persistent algorithm given in [19]. There are two steps of this algorithm: 1) matrix reduction; 2) pairing.

Matrix reduction (ELZ Algorithm)

Suppose $S(K)$ is the set of all simplices of K and $\sharp S(K) = m$.

Consider the function $f_{ind} : S(K) \rightarrow \{0, 1, \dots, N\}$, such that $f_{ind}(\sigma) = i$ for $\sigma \in K_i \setminus K_{i-1}$.

Choose a *compatible ordering* of the simplices, that is, an ordering of all simplices of K , such that $\sigma \prec \tau$ if $f_{ind}(\sigma) < f_{ind}(\tau)$ or if σ is a face of τ . After we order all simplices of K according to this compatible ordering, we get a sequence of simplices $\sigma_1, \sigma_2, \dots, \sigma_m$.

Consider the m -by- m boundary matrix ∂ given by

$$\partial[i, j] = \begin{cases} 1 & \text{if } \sigma_i \text{ is a codimension-1 face of } \sigma_j; \\ 0 & \text{otherwise.} \end{cases}$$

Let $low(j)$ be the row index of the lowest 1 in column j . If the entire column is zero, then $low(j)$ is undefined. We call R *reduced* if $low(j) \neq low(j_0)$ whenever j and j_0 , with $j \neq j_0$, specify two non-zero columns.

Algorithm 2.1 The ELZ Algorithm reduces ∂ by adding columns from left to right. The running time is at most cubic in the number of simplices [26], [30]. In matrix notation, the algorithm generates the reduced matrix $R = \partial \cdot V$. Since each simplex is preceded by its proper faces, ∂ is upper triangular. Since we only add from left to right, V is also upper triangular and so is R (See page 152-153, [19] for details).

Pairing

Once there are two cases of columns of reduced matrix R .

Case 1: column j of R is zero. We call σ_j positive since it creates a new cycle and thus gives birth to a new homology class unless it dies in the same filtration.

Case 2: column j of R is non-zero. We call σ_j negative because it provides the death to a homology class.

See page 154-155, [19] for details.

We can read bar codes directly from the matrix R . Suppose column j is zero and $\dim(\sigma_j) = r$. If there exists a column k with $\text{low}(k) = j$ and $f_{\text{ind}}(\sigma_k) > f_{\text{ind}}(\sigma_j)$ then column j provides a closed interval $[f_{\text{ind}}(\sigma_j), f_{\text{ind}}(\sigma_k) - 1] \in \mathcal{B}(\mathcal{H}_r(\mathcal{K}))$. If there exists a column k with $\text{low}(k) = j$ and $f_{\text{ind}}(\sigma_k) = f_{\text{ind}}(\sigma_j)$ then column j doesn't provide any interval since it indicates a cycle which dies in the same filtration. Note that $f_{\text{ind}}(\sigma_k) < f_{\text{ind}}(\sigma_j)$ is impossible according to the definition of compatible ordering. If there is no column k with $\text{low}(k) = j$ then column j provides an infinite interval $[f_{\text{ind}}(\sigma_j), \infty] \in \mathcal{B}(\mathcal{H}_r(\mathcal{K}))$.

Case B: $\kappa = \mathbb{R}$

For each s consider the chain complex $\mathcal{C}(K_s)$ with coefficients in \mathbb{R}

$$\cdots \xrightarrow{\partial_{r+2}^s} C_{r+1}^s \xrightarrow{\partial_{r+1}^s} C_r^s \xrightarrow{\partial_r^s} C_{r-1}^s \xrightarrow{\partial_{r-1}^s} \cdots \xrightarrow{\partial_2^s} C_1^s \xrightarrow{\partial_1^s} C_0^s \xrightarrow{\partial_0^s} = 0$$

equipped with the scalar products defined by the standard basis (see Subsection 2.2).

Let

- (1) $\delta_{r-1}^s = (\partial_r^s)^*$ the adjoint operator of ∂_r^s ⁶ for $r \in \mathbb{Z}_{\geq 0}$,
- (2) $\Delta_r^s = \partial_{r+1}^s \circ \delta_r^s + \delta_{r-1}^s \circ \partial_r^s : C_r^s \rightarrow C_r^s$ for $r \in \mathbb{Z}_{\geq 0}$ and
- (3) $(C_r^s)_+ = \text{im}(\partial_{r+1}^s)$, $(C_r^s)_- = \text{im}(\delta_{r-1}^s)$ and $H_r^s = \ker(\Delta_r^s)$, with

$$H_r^s = \ker(\delta_r^s) \cap \ker(\partial_r^s)$$

and

$$C_r^s = (C_r^s)_+ \oplus H_r^s \oplus (C_r^s)_-$$

⁶When represented as a matrix w.r.t. the standard base, “adjoint” w.r.t. the scalar product defined by the standard base is actually “transpose”.

where $(C_r^s)_+$, H_r^s and $(C_r^s)_-$ are pairwise orthogonal.

Recall that $[A]$ represents the orthonormalization of matrix A (see Lemma 2.3). With respect to the standard basis provided by simplices / cells, ∂_r^s can be regarded as an $n_{r-1}^s \times n_r^s$ matrix. In view of the above considerations, the following linear maps are orthogonal projections onto $(C_r^s)_+$, $(C_r^s)_-$ and H_r^s , respectively.

$$\begin{aligned}
(p_r^s)_+ : C_r^s &\rightarrow C_r^s \\
y &\mapsto [\partial_{r+1}^s][\partial_{r+1}^s]^T y \\
(p_r^s)_- : C_r^s &\rightarrow C_r^s \\
y &\mapsto [(\partial_r^s)^T][(\partial_r^s)^T]^T y \\
(p_r^s)_H : C_r^s &\rightarrow C_r^s \\
y &\mapsto (I_{n_r^s} - [\partial_{r+1}^s][\partial_{r+1}^s]^T - [(\partial_r^s)^T][(\partial_r^s)^T]^T) y
\end{aligned} \tag{3.6}$$

$$\begin{aligned}
j_r^s : H_r^s = \ker(\delta_r^s) \cap \ker(\partial_r^s) &\rightarrow \mathcal{H}_r^s = \ker(\partial_r^s) / \text{im}(\partial_{r+1}^s) \\
x &\mapsto x + \text{im}(\partial_{r+1}^s)
\end{aligned}$$

and

$$\begin{aligned}
k_r^s : \mathcal{H}_r^s = \ker(\partial_r^s) / \text{im}(\partial_{r+1}^s) &\rightarrow H_r^s = \ker(\delta_r^s) \cap \ker(\partial_r^s) \\
y + \text{im}(\partial_{r+1}^s) &\mapsto (p_r^s)_H(y)
\end{aligned}$$

is a pair of isomorphisms between H_r^s and \mathcal{H}_r^s .

Since H_r^s identifies to \mathcal{H}_r^s by the pair of isomorphisms j_r^s and k_r^s , the persistence vector space

$$H_r(\mathcal{K}) = \{H_r^s, g_r^s : H_r^s \rightarrow H_r^{s+1} \mid s \in \mathbb{Z}_{\geq 0}\},$$

where $g_r^s = k_r^{s+1} \circ \widetilde{i}_r^s \circ j_r^s$, is isomorphic to

$$\mathcal{H}_r(\mathcal{K}) = \{\mathcal{H}_r^s, \widetilde{i}_r^s : \mathcal{H}_r^s \rightarrow \mathcal{H}_r^{s+1} \mid s \in \mathbb{Z}_{\geq 0}\}$$

By Theorem 3.5 one has $\mathcal{B}(H_r(\mathcal{K})) = \mathcal{B}(\mathcal{H}_r(\mathcal{K}))$.

Notation 3.12. 1) Denote

$$g_r^{s,t} = g_r^{t-1} \circ \cdots \circ g_r^s : H_r^s \rightarrow H_r^t$$

for $s < t$ and

$$g_r^{s,s} = \text{identity} : H_r^s \rightarrow H_r^s$$

with $s, t \in \mathbb{Z}_{\geq 0}$.

Define $i_r^{s,t} : C_r^s \rightarrow C_r^t$ and $\widetilde{i}_r^{s,t} : \mathcal{H}_r^s \rightarrow \mathcal{H}_r^t$ in similar way.

2) Denote $\beta_r(s, t) = \dim(\text{im}(g_r^{s,t} : H_r^s \rightarrow H_r^t))$ with $s \leq t \in \mathbb{Z}_{\geq 0}$.

Note. $\beta_r(s, s) = \dim H_r^s$ and

$$\beta_r(s, t) = \beta_r(s, t + 1) \quad (3.7)$$

for $t \geq N$.

Hence $\beta_r(s, \infty) = \beta_r(s, m)$ for m larger than s and N .

Notice that $\beta_r(s, t)$ = number of intervals in $\mathcal{B}(H_r(\mathcal{K}))$ which contain $[s, t]$ for $0 \leq s \leq t \leq \infty$. (See Observation 3.6)

3) Denote $\mu_r(s, t)$ = number of intervals in $\mathcal{B}(H_r(\mathcal{K}))$ which equal to $[s, t]$ for $0 \leq s \leq t \leq \infty$. We have

$$\begin{aligned} & \mu_r(s, t) \\ = & \begin{cases} \beta_r(s, t) - \beta_r(s-1, t) - \beta_r(s, t+1) + \beta_r(s-1, t+1), & 0 < s \leq t < \infty \\ \beta_r(0, t) - \beta_r(0, t+1), & s = 0, 0 \leq t < \infty \\ \beta_r(s, \infty) - \beta_r(s-1, \infty), & 0 < s < \infty, t = \infty \\ \beta_r(0, \infty), & s = 0, t = \infty \end{cases} \end{aligned} \quad (3.8)$$

Note. From the above formula (3.8) and formula (3.7) we only need to know $\beta_r(s, t)$ for all $0 \leq s \leq t \leq N$ in order to calculate $\mu_r(s, t)$ and the bar code $\mathcal{B}(H_r(\mathcal{K}))$.

Theorem 3.10. $\beta_r(s, s) = \text{rank}((p_r^s)_H)$ and $\beta_r(s, t) = \text{rank}((p_r^t)_H \circ i_r^{s,t} \circ (p_r^s)_H)$ for $s < t$, where

$$i_r^{s,t} = \begin{pmatrix} I_{n_r^s} \\ 0_{(n_r^t - n_r^s) \times n_r^s} \end{pmatrix}.$$

For the definition of $(p_r^s)_H$, refer to (2.1) and (3.6).

We will use MATLAB and in order to avoid the use of the function “rank” which is not reliable for large matrices we will need the following observations to complement Theorem 3.10.

Observation 3.11. The rank of a real matrix A equals to the number of positive eigenvalues of AA^T or $A^T A$.

Observation 3.12. $\beta_r(s, s) = \dim(H_r^s) = \dim(C_r^s) - \text{rank}(\partial_{r+1}^s) - \text{rank}(\partial_r^s)$.

Observation 3.13. If $\beta_r(s, t) \neq 0$, $s < t$, then

- i) H_r^i is nonzero for all $s \leq i \leq t$;
- ii) $\beta_r(s, i)$ is nonzero for all $s \leq i \leq t$.

Simultaneous Persistence

In the case $\kappa = \mathbb{Z}_2$ we describe how to calculate the bar codes of \mathbb{Z}_2 -homology groups of the filtered simplicial complex

$$K_0 \subseteq K_1 \subseteq \cdots \subseteq K_N = K.$$

Let's discuss a more general case where the filtration has two directions.

Suppose X is a polytopal complex and $X^\pm \subset X$ are two subcomplexes with $X_0 = X^+ \cap X^-$. Suppose X^+ and X^- are equipped with finite filtration $X_0^+ \subseteq X_1^+ \subseteq \cdots \subseteq X_{N^+}^+$ and $X_0^- \subseteq X_1^- \subseteq \cdots \subseteq X_{N^-}^-$ where $X_0^\pm = X_0$, $X_{N^-}^- = X^-$ and $X_{N^+}^+ = X^+$.

Define the simultaneous persistence number $\omega_r(s, t)$ of the above filtration to be the maximal number of linearly independent elements in $H_r(X_0)$ which dies in X^- at s and in X^+ at t (See [3] for details). These numbers are useful in the calculation of relevant level persistence numbers.

One can put the cells of X in three groups I, II, III. The group I contains the cells of X_0 , the group II the cells of $X^- \setminus X_0$ and the group III the cells of $X^+ \setminus X_0$.

Consider the following filtration of X

$$X_0 \subseteq X_1 \subseteq \cdots \subseteq X_{N^-+N^+} = X$$

where $X_i = X_i^-$ for $0 \leq i \leq N^-$, $X_i = X^- \cup X_{i-N^-}^+$ for $N^-+1 \leq i \leq N^-+N^+$.

Suppose $S(X)$ is the set of all cells of X and $\sharp S(X) = m$.

We can define a function $f_{ind} : S(X) \rightarrow \{0, 1, \dots, N^- + N^+\}$, such that $f(\sigma) = i$ for $\sigma \in X_i \setminus X_{i-1}$.

We use a *compatible ordering* of the cells, that is, an ordering of all cells of X , such that $\sigma \prec \tau$ if $f_{ind}(\sigma) < f_{ind}(\tau)$ or if σ is a face of τ . After we order all cells of X according to this compatible ordering, we get a sequence of cells $\sigma_1, \sigma_2, \dots, \sigma_m$.

With this compatible ordering, the incidence matrix M is given by

$$M = \begin{Bmatrix} A & B^- & C^+ \\ 0 & B & 0 \\ 0 & 0 & C \end{Bmatrix}$$

with A being the incidence matrix for the subcomplex X_0 , $M^- = \begin{Bmatrix} A & B^- \\ 0 & B \end{Bmatrix}$

being the incidence matrix for X^- and $M^+ = \begin{Bmatrix} A & C^+ \\ 0 & C \end{Bmatrix}$ being the incidence matrix for X^+ .

The matrix M is said to be in *relative reduced form* if both M^- and M^+ are in reduced form. This relative reduced form R_{rel} can be achieved by applying Algorithm 2.1 to matrix M .

We can read the number $\omega_r(i, j)$ directly from the relative reduced form R_{rel} of M . $\omega_r(i, j)$ = the number of all triples $(\sigma_k, \sigma_{k'}, \sigma_{k''})$ with $\sigma_k \in X_0$, column k of R_{rel} is zero, $\dim(\sigma_k) = r$, $low(k') = low(k'') = k$, $\sigma_{k'} \in X_i^- \setminus X_{i-1}^-$ and $\sigma_{k''} \in X_j^+ \setminus X_{j-1}^+$.

3.5 Calculate the Bar Codes of a PCD

Given a point cloud data $X = \{x_1, x_2, \dots, x_p\} \in \mathbb{R}^n$ (suppose all points in X are distinct), we will first consider its distance matrix

$$D = \begin{pmatrix} d(x_1, x_1) & d(x_1, x_2) & \cdots & d(x_1, x_p) \\ d(x_2, x_1) & d(x_2, x_2) & \cdots & d(x_2, x_p) \\ \vdots & \vdots & \ddots & \vdots \\ d(x_p, x_1) & d(x_p, x_2) & \cdots & d(x_p, x_p) \end{pmatrix}$$

Notice that the D is a symmetric matrix, the diagonal elements of D are zeros and all the other elements are positive real numbers.

Let S be the set of the entries in D and order these numbers increasingly (since they will be our epsilons), so we have

$$\mathcal{E} : 0 = \epsilon_0 < \epsilon_1 < \epsilon_2 < \cdots < \epsilon_N.$$

\mathcal{E} is exactly (3.0).

Given a positive integer m and P , let's compute the bar code of PCD X up to dimension $m - 1$ and step P (for definition see (3.3)).

To simplify the notation, from now on we will write K_i instead of $R_{\epsilon_i}(X, m)$. K_0 is just a collection of vertices $\{x_1, x_2, \dots, x_p\}$.

We will represent the vertex x_{i_0} simply by the index i_0 , the edge $[x_{i_0}, x_{i_1}]$ by $[i_0, i_1]$ and k -simplex $[x_{i_0}, \dots, x_{i_k}]$ is represented by $[i_0, i_1, \dots, i_k]$.

We want first to have the information about the filtered simplicial complex stored as

(a) An $(m + 1) \times (P + 1)$ matrix named **Dimension** where the entry **Dimension** $_{r+1, s+1}$ is the integer $n_r^s = \text{dimension of } C_r^s$.

(b) An $(m + 1)$ array of matrices named **Simplex**. The component **Simplex** $_{r+1}$ is a matrix describing r -simplices of K_P ; precisely is a matrix of $r + 1$ columns, each row representing an r -simplex $[i_0, i_1, \dots, i_k]$. These rows are ordered consistently with the filtration. The order for simplices of $K_s \setminus K_{s-1}$ is however arbitrary.

There is an obvious way to create these data, which is not very efficient but worth to mention. Given K_s , we can easily find its 1-simplices from the distance matrix D . All the nonzero numbers in D which are less than ϵ_s correspond to the edges in K_s . In order to find the r -simplices, we scan all possible combination of $r + 1$ numbers of $\{1, 2, \dots, p\}$. $[i_0, i_1, \dots, i_r]$ ($i_0 < i_1 < \dots < i_r$) belongs to the r -simplices of K_s iff every pair $[i_j, i_k]$ belongs to the 1-simplices of K_s . As s increases, we have more r -simplices in K_s . Store each r -simplex as a row in **Simplex** _{$r+1$} in filtration order. In another word, if $[i_0, i_1, \dots, i_r]$ belongs to K_s but not in K_{s-1} , then it is “new” and can be added as a new row to the rows provided by the r -simplices of K_{s-1} .

We use the package **JPlex** [1] [34] to compute **Dimension** and **Simplex**, which can be download from <http://comptop.stanford.edu/programs/jplex/>.

Once we get **Simplex**, we get the standard basis of C_r^P , which is stored in **Simplex** _{$r+1$} , for $0 \leq r \leq m$. Given **Simplex** _{r} and **Simplex** _{$r+1$} , we compute ∂_r^P with respect to the standard basis:

Write **Simplex** _{$r+1$} as

$$\begin{pmatrix} a_{11} & \cdots & a_{1,r+1} \\ \vdots & \ddots & \vdots \\ a_{n_r^P,1} & \cdots & a_{n_r^P,r+1} \end{pmatrix}$$

and **Simplex** _{r} as

$$\begin{pmatrix} b_{11} & \cdots & b_{1,r} \\ \vdots & \ddots & \vdots \\ b_{n_{r-1}^P,1} & \cdots & b_{n_{r-1}^P,r} \end{pmatrix},$$

where n_r^P and n_{r-1}^P can be obtained from **Dimension**.

∂_r^P is an $n_{r-1}^P \times n_r^P$ matrix with element $(\partial_r^P)_{i,j}$ equals to the incidence number of the simplices $[b_{i,1}, \dots, b_{i,r}]$ and $[a_{j,1}, \dots, a_{j,r+1}]$. If $[b_{i,1}, \dots, b_{i,r}]$ is not contained in $[a_{j,1}, \dots, a_{j,r+1}]$, $(\partial_r^P)_{i,j} = 0$. Otherwise, $[b_{i,1}, \dots, b_{i,r}]$ equals to $[a_{j,1}, \dots, \hat{a}_{j,k}, \dots, a_{j,r+1}]$ for some k , where “ $\hat{}$ ” means deletion. If k is odd, $(\partial_r^P)_{i,j} = 1$; if k is even, $(\partial_r^P)_{i,j} = -1$.

∂_r^s is the $n_{r-1}^s \times n_r^s$ upper-left block of ∂_r^P .

Once **Dimension** and **Simplex** are determined, we apply Theorem 3.10 and the Observations 3.11, 3.12 and 3.13 to get $\beta_r(s, t)$ and then $\mu_r(s, t)$ according to (3.8).

Here is a brief description of functions written in Matlab. Those Matlab programs contained in “Computing_Barcode_of_PCD.zip” can be downloaded from <http://people.math.osu.edu/du.50>

Initial Data:

A $p \times n$ matrix **X**, which stores p points x_1, \dots, x_p in \mathbb{R}^n as rows, two positive integers **m** and **S**. **m** is the given upper bound of dimension of restricted Vietoris-Rips Complex. **S** is the given upper bound of steps of filtration. $\mathbf{P} = \min(\mathbf{S}, \mathbf{N})$. See (3.3).

1. Function getDistance

Input: **X**.

Output: A $p \times p$ upper triangular matrix **D**. $\mathbf{D}_{ij} = d(x_i, x_j), 1 \leq i \leq p-1, i+1 \leq j \leq p$.

2. Function scaleX

Input: **X**, **D**.

Output: **X**.

Given a point cloud data **X**, we find all the distances between two points and sort them as $\epsilon_0 < \epsilon_1 < \dots < \epsilon_N$. **Note:** If $\epsilon_{i+1} - \epsilon_i$ is very small (less than $3 \cdot \text{diam}(\mathbf{X}) \cdot \text{eps}$, the upper bound of round off errors, where $\text{diam}(\mathbf{X})$ is the maximal distance between any two points of **X**, $\text{eps} = 2.2204 \times 10^{-16}$ is floating-point relative accuracy in Matlab), we delete ϵ_{i+1} . We scale **X** such that the minimum difference between two consecutive distances is larger than 3×10^{-4} . For details and more explanation, read the comments in my program.

3. Function getEpsilon

Input: **D**, **S**.

Output: Two positive integers $\mathbf{P} = \min(\mathbf{S}, \mathbf{N})$ and $\mathbf{e} = \epsilon_P$. Two row vectors **epsiorg**, which stores different distances $\epsilon_0 < \epsilon_1 < \dots < \epsilon_N$ of **D** in increasing order and **epsiavg**, which stores

$$\frac{\epsilon_0 + \epsilon_1}{2}, \dots, \frac{\epsilon_k + \epsilon_{k+1}}{2}, \dots, \frac{\epsilon_{N-1} + \epsilon_N}{2}, \epsilon_N + \frac{1}{2}.$$

4. Function getDimensionSimplex

Input: **X**, **P**, **epsiavg**, **m** and **e**.

Output:

- a) An integer $\mathbf{m}_a = \dim(R_{\epsilon_P}(X, m))$.
- b) An $(\mathbf{m}_a + 1) \times (\mathbf{P} + 1)$ matrix **Dimension**. $\text{Dimension}_{r+1, s+1} = n_r^s$, the dimension of C_r^s .
- c) An $\mathbf{m}_a + 1$ array **Simplex**.

The component Simplex_{r+1} is a matrix with $r+1$ columns, which stores r -simplex of $R_{\epsilon_P}(X, m)$ as rows in filtration order.

Note: Here we use the package **edu.stanford.math.plex.*** of **JPlex**[1][34].

5. Function getDeltaP

Input: **Dimension**, **Simplex** and \mathbf{m}_a .

Output: An \mathbf{m}_a array **DeltaP**. $\mathbf{DeltaP}_r =$ the matrix of $\partial_r^P : C_r^P \rightarrow C_{r-1}^P$ with respect to the the standard basis of C_r^P and C_{r-1}^P .

6. Function getDelta

Input: **DeltaP**, **Dimension**, r , s

Output: The matrix $\partial_r^s : C_r^s \rightarrow C_{r-1}^s$ with respect to standard basis of C_r^s and C_{r-1}^s .

7. Function getRank

Input: **P**, **m**, and \mathbf{m}_a .

Output: An $(\mathbf{m} + 1) \times (\mathbf{P} + 1)$ matrix **R**. $\mathbf{R}_{r+1,s+1} = \text{rank of } \partial_r^s$.

Algorithm: $\mathbf{R}_{r+1,s+1} =$ number of eigenvalues of $\partial_r^s \times \partial_r^{s'}$ which is greater than 10^{-8} . Here we use Observation 3.11.

8. Function getDimensionHrs

Input: **Dimension**, **R**, **P**, **m** and \mathbf{m}_a .

Output: An integer \mathbf{m}_c and an $(\mathbf{m}_c + 1) \times (\mathbf{P} + 1)$ matrix **dimHrs**. $\mathbf{dimHrs}_{r+1,s+1} = \dim(H_r^s)$.

Algorithm: $\mathbf{dimHrs}_{r+1,s+1} = \mathbf{Dimension}_{r+1,s+1} - \mathbf{R}_{r+2,s+1} - \mathbf{R}_{r+1,s+1}$. Here we use use Observation 3.12.

9. Function getHarmonicProjection

Input: **Dimension**, r and s .

Output: The harmonic projection matrix $(p_r^s)_H$ from C_r^s to itself with respect to the standard basis of C_r^s .

Algorithm: See (3.6).

10. Function getBeta

Input: **Dimension**, **P**, **dimHrs** and \mathbf{m}_c . \mathbf{m}_c is the maximal dimension of nonzero H_r^s 's.

Output: A $(\mathbf{P} + 1) \times (\mathbf{P} + 1) \times (\mathbf{m}_c + 1)$ matrix **Beta**. $\mathbf{Beta}(s + 1, t + 1, r + 1) = \beta_r(s, t)$.

Algorithm:

$$\beta_0(\cdot, \cdot) = \begin{pmatrix} \dim(H_0^0) & \dim(H_0^1) & \cdots & \dim(H_0^P) \\ & \dim(H_0^1) & \cdots & \dim(H_0^P) \\ & & \ddots & \vdots \\ & & & \dim(H_0^P) \end{pmatrix};$$

$\beta_r(s, s) = \text{rank}((p_r^s)_H) = \dim(H_r^s)$ (see Observation 3.12);

$\beta_r(s, t) = \text{rank}((p_r^t)_H \circ i_r^{s,t} \circ (p_r^s)_H)$ for $s < t$ (see Theorem 3.10).

Apply Observation 3.11 when computing rank as in function getRank.

Apply Observation 3.13 to avoid unnecessary computing.

11. Function getMu

Input: \mathbf{Beta} , \mathbf{P} and $\mathbf{m_c}$.

Output: A $(\mathbf{P} + 1) \times (\mathbf{P} + 1) \times (\mathbf{m_c} + 1)$ matrix \mathbf{Mu} . $\mathbf{Mu}(s + 1, t + 1, r + 1) = \mu_r(s, t)$

Algorithm: Apply (3.8).

12. Function `getBar codeMatrix`

Input: \mathbf{Mu} , \mathbf{P} , and $\mathbf{m_c}$.

Output: \mathbf{A} , a $q \times 4$ matrix which stores the bar code in increasing order according to dimension, left endpoint and right endpoint. M intervals $[a, b]$ in bar code $\mathcal{B}(\mathcal{H}_r(\mathcal{K}))$ will be represented by $[r \ a \ b \ M]$ in \mathbf{A} .

13. `drawBar code`

Input: \mathbf{A} and \mathbf{P} .

Output: A picture of bar code stored in `barcode.eps`.

14. Function `main`

Input: \mathbf{X} , \mathbf{S} and \mathbf{m} .

Output: A matrix \mathbf{A} stored in `A.mat` and a picture of bar code stored in `barcode.eps`.

Algorithm: Run above functions except 6. and 9. in order.

3.6 Numerical Experiment

There is a PCD in \mathbb{R}^2 with 53 points (Figure 3). The restricted filtered simplicial complex of this PCD ($m = 4$, $P = 35$) is illustrated in Figure 4. The bar codes of the filtered simplicial complex is illustrated in Figure 5.

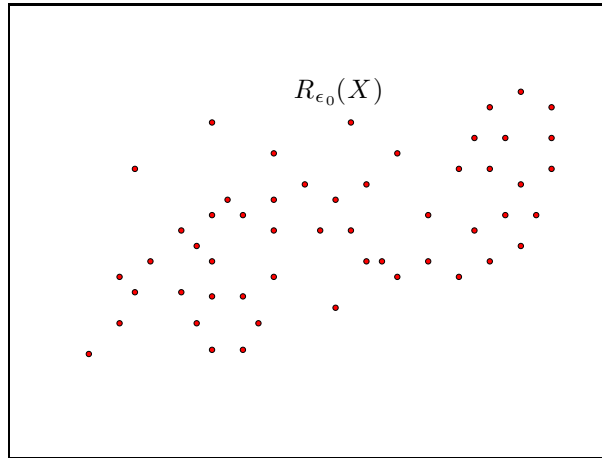


Figure 3: A PCD in \mathbb{R}^2 with 53 points

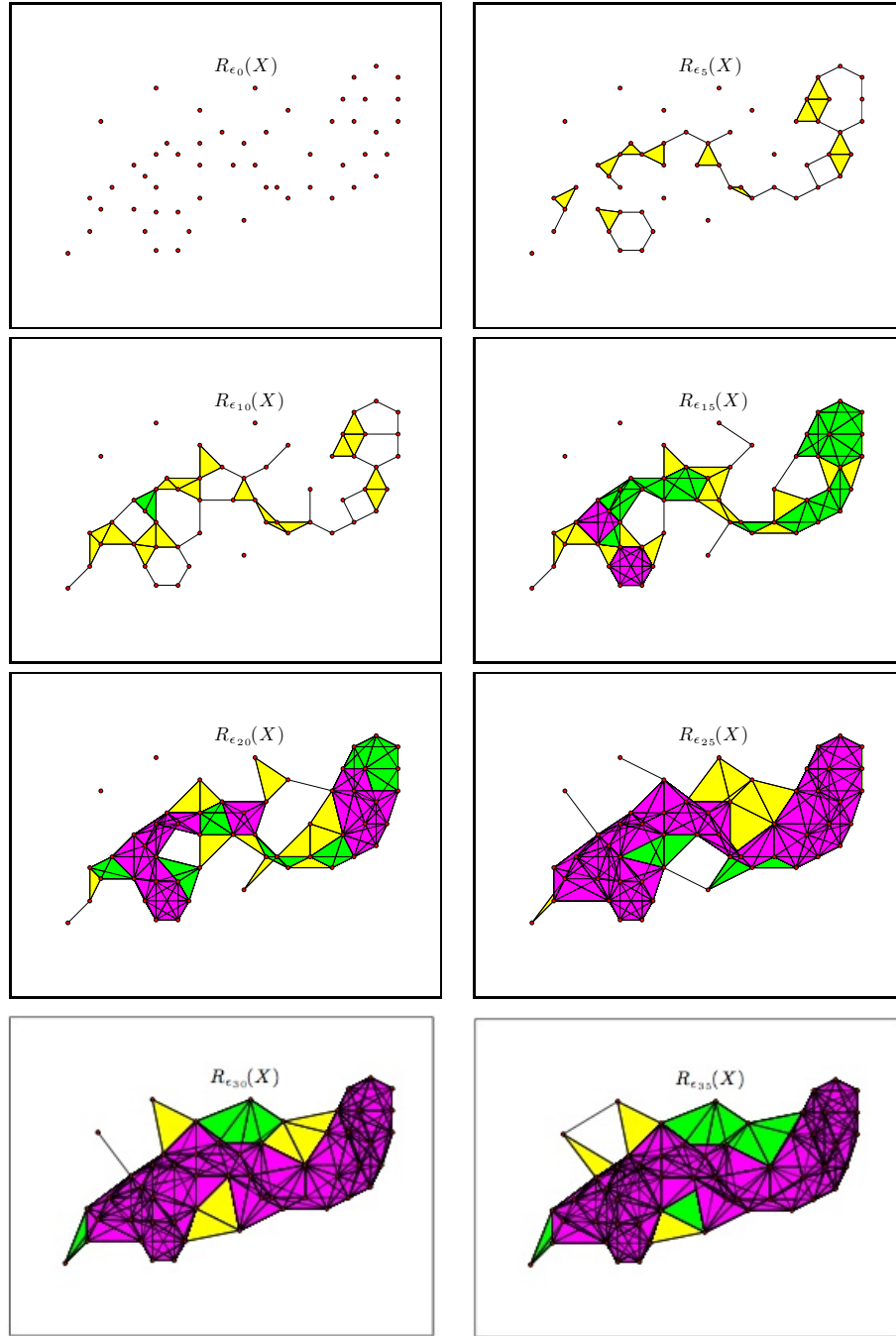


Figure 4: Restricted filtered simplicial complex of this PCD ($m = 4$, $P = 35$).

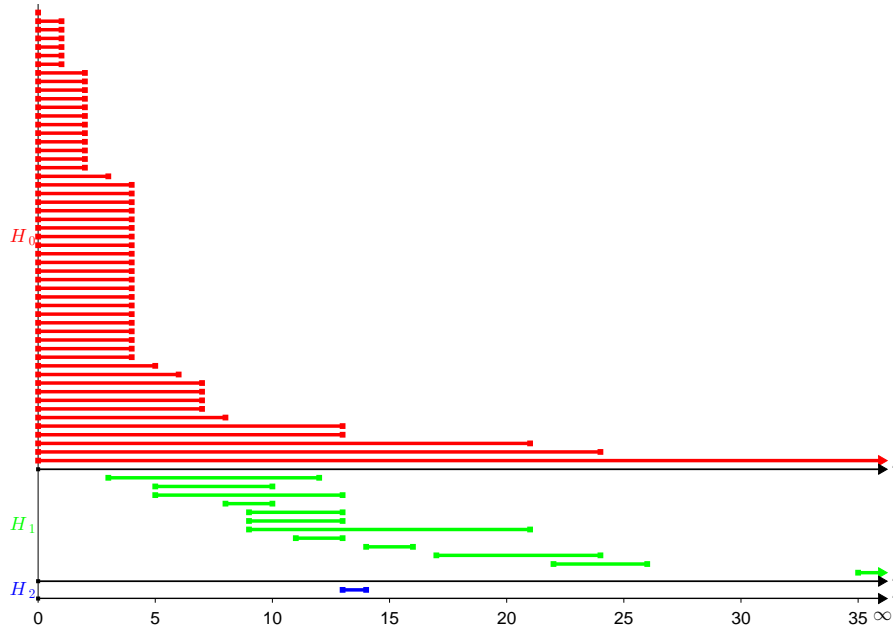


Figure 5: Bar codes of the above filtered simplicial complex.

4 Persistence Theory Refined

In this section we consider refinements of standard persistence for a real valued map and describe the calculation of its invariants in the case of a linear map $f : X \rightarrow \mathbb{R}$, X a simplicial complex. In order to understand the novelty of our considerations we will review in Subsection 4.1 the standard persistence named in this section sub-level persistence.

Let $f : X \rightarrow \mathbb{R}$ be a continuous map, denote by

$$X_{-\infty, t} := f^{-1}((-\infty, t])$$

$$X_{t, \infty} := f^{-1}([t, \infty))$$

$$X_t := f^{-1}(t) \text{ and}$$

$$X_{s, t} := f^{-1}([s, t]), s \leq t$$

4.1 Sub-level Persistence [3]

Given a continuous map $f : X \rightarrow \mathbb{R}$, the sub-level persistent homology introduced in [20] and further developed in [38] is concerned with the following

questions:

Q1. Does the class $x \in H_r(X_{-\infty,t})$ originate in $H_r(X_{-\infty,t''})$ for $t'' \leq t$?
Does the class $x \in H_r(X_{-\infty,t})$ vanish in $H_r(X_{-\infty,t'})$ for $t < t'$?

Q2. What are the smallest t' and t'' such that this happens?

The information that is contained in the linear maps $H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,t'})$ for any $t \leq t'$ is known as **sub-level persistence**.

Recall that (sub-level)**persistent homology** is the collection of vector spaces and linear maps $\{H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,t'}), t < t', t, t' \in \mathbb{R}\}$. Let

$0 \neq x \in H_r(X_{-\infty,t})$. One says that

(i) x is born at t'' , $t'' \leq t$, if x is contained in $\text{img}(H_r(X_{-\infty,t''}) \rightarrow H_r(X_{-\infty,t}))$ but is not contained in $\text{img}(H_r(X_{-\infty,t''-\epsilon}) \rightarrow H_r(X_{-\infty,t}))$ for any $\epsilon > 0$.

(ii) x dies at t' , $t' > t$, if its image is zero in $\text{img}(H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,t'}))$ but is nonzero in $\text{img}(H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,t'-\epsilon}))$ for any $0 < \epsilon < t' - t$.

(iii) x dies at ∞ , if its image is always nonzero in $\text{img}(H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,t'}))$ for any $t' > t$.

The standard construction “telescope” in homotopy theory permits to replace any finite filtered space $K_0 \subseteq K_1 \subseteq \cdots \subseteq K_N$ by a weakly tame map $f : X \rightarrow \mathbb{R}$, simply by taking $X = K_0 \times [t_0, t_1] \cup_{\phi_1} K_1 \times [t_1, t_2] \cup_{\phi_2} \cdots \cup_{\phi_{N-1}} K_{N-1} \times [t_{N-1}, t_N] \cup_{\phi_N} K_N$ where $\phi_i : K_i \times \{t_{i+1}\} \rightarrow K_{i+1} \times \{t_{i+1}\}$ is the inclusion and $f|_{K_i \times [t_i, t_{i+1}]}$ the projection of $K_i \times [t_i, t_{i+1}]$ on $[t_i, t_{i+1}]$.

The sub-level persistence for a filtered space is the sub-level persistence of the associated weakly tame map.

When f is weakly tame, the sub-level persistence for each $r = 0, 1, \dots, \dim X$ is determined by a finite collection of invariants referred to as **bar codes for sub-level persistence** [38]. The r -bar codes for sub-level persistence of f are intervals of the form $[t, t')$ or $[t, \infty)$ with $t < t'$.

The number $\mu_r(t, t')$ of r -bar codes which identify to the interval $[t, t')$ is the maximal number of linearly independent homology classes in $H_r(X_{-\infty,t})$, which are born at t , die at t' and remain independent in $\text{img}(H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,s}))$ for any $s, t \leq s < t'$.

The number $\mu_r(t, \infty)$ of r -bar codes which identify to the interval $[t, \infty)$ is the maximal number of linearly independent homology classes in $H_r(X_{-\infty,t})$ which are born at t , never die and remain independent in $\text{img}(H_r(X_{-\infty,t}) \rightarrow H_r(X_{-\infty,s}))$ for any $s > t$.

Lemma 4.1. *The set of r -bar codes for sub-level persistence is finite and any r -bar code is an intervals of the form $[t_i, t_j)$ or $[t_i, \infty)$ with $t_i < t_j$ t_i, t_j critical values of f .*

Proof. Indeed the ends of the “bar codes” have to be critical values. Suppose

there is an r -bar code of the form $[t, t']$, we have $t_0 \leq t < t' \leq t_N$. Suppose $t_i \leq t < t_{i+1}$ for some i and $t_j \leq t' < t_{j+1}$ for some j .

There exists a class $x \in H_r(X_{-\infty, t})$ that is born at t and dies at t' . Since the canonical inclusion $X_{-\infty, t_i} \hookrightarrow X_{-\infty, t}$ is a deformation retraction and induces an isomorphism on homology groups, we have $x \in \text{img}(H_r(X_{-\infty, t_i}) \rightarrow H_r(X_{-\infty, t}))$. If $t > t_i$, it will cause a contradiction with the fact that x is born at t . So we have $t = t_i$ and the bar code is of the form $[t_i, t']$.

Consider the commutative diagram below

$$\begin{array}{ccc} X_{-\infty, t_i} & \longrightarrow & X_{-\infty, t_j} \\ & \searrow & \downarrow \cong \\ & & X_{-\infty, t'} \end{array}$$

Since the canonical inclusion $X_{-\infty, t_j} \hookrightarrow X_{-\infty, t'}$ is a deformation retraction, the vertical map in the above diagram is an isomorphism. The image of x in $\text{img}(H_r(X_{-\infty, t_i}) \rightarrow H_r(X_{-\infty, t_j}))$ is zero. If $t' > t_j$, it will cause a contradiction with the fact that x dies at t' . So we have $t' = t_j$ and the bar code is of the form $[t_i, t_j]$.

For the same reason, an r -bar code $[t, \infty)$ is of the form $[t_i, \infty)$. The finiteness of the collection of r -bar codes follows from the finiteness of $\dim H_r(X_{-\infty, t})$. \square

From these bar codes one can derive the **Betti numbers** $\beta_r(t, t')$, the dimension of

$\text{img}(H_r(X_{-\infty, t}) \rightarrow H_r(X_{-\infty, t'}))$, for any $t \leq t'$ and get the answers to questions Q1 and Q2. For example,

$$\beta_r(t, t') = \text{the number of } r\text{-bar codes which contain the interval } [t, t']. \quad (4.1)$$

From the Betti numbers $\beta_r(t, t')$ one can also derive these r -bar codes.

Denote $\mu_r(t_i, t_j)$ = number of r -bar codes which equal to $[t_i, t_j]$ for $t_0 \leq t_i < t_j \leq \infty$, where t_0 is the smallest critical value. We have

$$\begin{aligned} & \mu_r(t_i, t_j) \\ = & \begin{cases} \beta_r(t_i, t_{j-1}) - \beta_r(t_{i-1}, t_{j-1}) - \beta_r(t_i, t_j) + \beta_r(t_{i-1}, t_j), & t_0 < t_i < t_j < \infty \\ \beta_r(t_0, t_{j-1}) - \beta_r(t_0, t_j), & t_i = t_0, t_0 < t_j < \infty \\ \beta_r(t_i, \infty) - \beta_r(t_{i-1}, \infty), & t_0 < t_i < \infty, t_j = \infty \\ \beta_r(t_0, \infty), & t_i = t_0, t_j = \infty \end{cases} \end{aligned} \quad (4.2)$$

The computation of the bar codes for a filtration of simplicial / polytopal complex is discussed in Subsection 3.4 with coefficient field of homology groups being \mathbb{Z}_2 or \mathbb{R} .

4.2 Level Persistence[3]

Level persistence for a map $f : X \rightarrow \mathbb{R}$ was first considered in [16] and was better understood when the Zigzag persistence was introduced and formulated in [7]. Given a continuous map $f : X \rightarrow \mathbb{R}$, level persistence is concerned with the homology of the fibers $H_r(X_t)$ and addresses questions of the following type.

Q1. Does the image of $x \in H_r(X_t)$ vanish in $H_r(X_{t,t'})$, where $t' > t$ or in $H_r(X_{t'',t})$, where $t'' < t$?

Q2. Can x be detected in $H_r(X_{t'})$ where $t' > t$ or in $H_r(X_{t''})$ where $t'' < t$? The precise meaning of detection is explained below.

Q3. What are the smallest t' and t'' for the answers to Q1 and Q2 to be affirmative?

To answer such questions one has to record information about the following maps:

$$H_r(X_t) \rightarrow H_r(X_{t,t'}) \leftarrow H_r(X_{t'}).$$

The **level persistence** is the information provided by this collection of vector spaces and linear maps considered above for all t, t' .

Let $0 \neq c \in H_r(X_t)$. One says that

(i) c dies downward at $t' < t$, if its image is zero in $\text{img}(H_r(X_t) \rightarrow H_r(X_{t',t}))$ but is nonzero in $\text{img}(H_r(X_t) \rightarrow H_r(X_{t'+\epsilon,t}))$ for any $0 < \epsilon < t - t'$.

(ii) c dies upward at $t'' > t$, if its image is zero in $\text{img}(H_r(X_t) \rightarrow H_r(X_{t,t''}))$ but is nonzero in $\text{img}(H_r(X_t) \rightarrow H_r(X_{t,t''-\epsilon}))$ for any $0 < \epsilon < t'' - t$.

We say that $x \in H_r(X_t)$ can be detected at $t' \geq t$, if its image in $H_r(X_{t,t'})$ is nonzero and is contained in the image of $H_r(X_{t'}) \rightarrow H_r(X_{t,t'})$. Similarly, the detection of x can be defined for $t'' < t$ also.

Definition 4.1. A continuous map $f : X \rightarrow \mathbb{R}$ is called **tame** (Definition 3.5, [3]) if X is compact and there exists finitely many values $\min(f(X)) = t_0 < t_1 < \dots < t_N = \max(f(X))$ (so called critical values) so that:

(i) for any $t \neq t_0, t_1, \dots, t_N$ there exists $\epsilon > 0$ so that $f : f^{-1}(t - \epsilon, t + \epsilon) \rightarrow (t - \epsilon, t + \epsilon)$ and the second factor projection $X_t \times (t - \epsilon, t + \epsilon) \rightarrow (t - \epsilon, t + \epsilon)$ are fiber-wise homotopy equivalent.

(ii) for any t_i there exists $\epsilon > 0$ so that canonical inclusions $X_{t_i} \hookrightarrow X_{t_i, t_i + \epsilon}$ and $X_{t_i} \hookrightarrow X_{t_i - \epsilon, t_i}$ are deformation retractions.

This means that the homology group of the level set changes at finitely many t 's. We have

Lemma 4.2. Given $t_{i-1} < s < t_i$, $X_{t_{i-1}, s}$ and X_{s, t_i} retracts by deformation onto $X_{t_{i-1}}$ and X_{t_i} respectively.

Proof. We will prove that X_{s,t_i} retracts by deformation onto $X_{t_{i-1}}$. The proof that X_{s,t_i} retracts by deformation onto X_{t_i} is essentially the same.

For any a , $t_{i-1} < a \leq s$, there exists an $\epsilon_a > 0$ such that $f : f^{-1}(a - \epsilon_a, a + \epsilon_a) \rightarrow (a - \epsilon_a, a + \epsilon_a)$ and the second factor projection $X_t \times (a - \epsilon_a, a + \epsilon_a) \rightarrow (a - \epsilon_a, a + \epsilon_a)$ are fiberwise homotopy equivalent. Hence for any $a - \epsilon_a < t'' < t' < a + \epsilon_a$, we have $X_{t'',t'}$ retract by deformation onto $X_{t''}$. For a , $t_{i-1} < a < s$ we can choose $0 < \epsilon_a < \min(a - t_{i-1}, s - a)$, so that $(a - \epsilon_a, a + \epsilon_a)$ is contained in (t_{i-1}, s) .

For t_{i-1} , there exists an $\epsilon_{t_{i-1}} > 0$ so that canonical inclusion $X_{t_i} \hookrightarrow X_{t_i, t_i + \epsilon_{t_{i-1}}}$ is a deformation retraction.

The union of open intervals $(a - \epsilon_a, a + \epsilon_a)$ for $t_{i-1} \leq a \leq s$ covers the closed interval $[t_{i-1}, s]$. Then finitely many of them (a_k, b_k) , $1 \leq k \leq m$ cover $[t_{i-1}, s]$. Suppose these intervals are not contained one in another, otherwise we can delete it. Observe that $(t_{i-1} - \epsilon_{t_{i-1}}, t_{i-1} + \epsilon_{t_{i-1}})$ and $(s - \epsilon_s, s + \epsilon_s)$ have to be contained in those finite collection of intervals.

Order (a_k, b_k) , $1 \leq k \leq m$ according to their right endpoints increasingly, we have $b_j < b_k$, when $j < k$. Observe that $b_{k-1} \in (a_k, b_k)$ and $a_k \in (a_{k-1}, b_{k-1})$. Therefore we also have $a_j < a_k$, when $j < k$ and the intersection of (a_{k-1}, b_{k-1}) and (a_k, b_k) is always nonempty. Since both left endpoints and right endpoints are ordered increasingly $(a_1, b_1) = (t_{i-1} - \epsilon_{t_{i-1}}, t_{i-1} + \epsilon_{t_{i-1}})$ and $(a_m, b_m) = (s - \epsilon_s, s + \epsilon_s)$.

Choose any $c_k \in (a_k, b_k) \cap (a_{k+1}, b_{k+1})$, $2 \leq k \leq m - 1$. Let $c_0 = t_{i-1}$, $c_1 = t_{i-1} + \epsilon_{t_{i-1}}$ and $c_m = s$.

We have X_{c_{k-1}, c_k} retracts by deformation onto $X_{c_{k-1}}$ for any $1 \leq k \leq m$.

Hence $X_{t_{i-1}, s} = X_{c_0, c_m}$ retracts by deformation onto $X_{t_{i-1}} = X_{c_0}$.

□

Corollary 4.3. *Tame maps are weakly tame.*

Note that the tame maps form a generic set of maps in the set of all continuous maps on a simplicial complex or on a smooth manifold (and more general on a compact ANR).

Lemma 4.4. *Given $t_{i-1} < s < t < t_i$, $X_{s,t}$ retracts by deformation onto X_s or X_t .*

In case of a tame map the collection of the vector spaces and linear maps is determined up to coherent isomorphisms by a collection of invariants called **bar codes for level persistence** which are intervals of the form $[t, t']$ with $t \leq t'$ and (t, t') , $(t, t']$, $[t, t')$ with $t < t'$.

These bar codes are called invariants because two tame maps $f : X \rightarrow \mathbb{R}$ and $g : Y \rightarrow \mathbb{R}$ which are fiber-wise homotopy equivalent have the same associated bar codes. The above result was established for the Zigzag persistence

from which can be derived, but can also be proven directly. However the proof is not contained in this paper.

An open end of an interval signifies the death of a homology class at that end (left or right) whereas a closed end signifies that a homology class cannot be detected beyond this level (left or right).

There exists an r -bar code (t'', t') if there exists a class $x \in H_r(X_t)$ for some $t'' < t < t'$ which is detectable for $t'' < s < t'$ and dies at t'' and t' . The multiplicity of (t'', t') is the maximal number of linearly independent classes in $H_r(X_t)$ such that:

- (i) all remain linearly independent in $\text{img}(H_r(X_t) \rightarrow H_r(X_{t,s}))$ for $t \leq s < t'$ and $\text{img}(H_r(X_t) \rightarrow H_r(X_{s,t}))$ for $t'' < s \leq t$;
- (ii) all die at t'' and t' .

Notice that the change of t above will not affect the multiplicity of (t'', t') .

There exists an r -bar code $(t'', t']$ if there exists an element $x \in H_r(X_{t'})$ which is not detectable for $s > t'$ and detectable for $t'' < s \leq t'$ and dies at t'' . The multiplicity of $(t'', t']$ is the maximal number of linearly independent elements in $H_r(X_{t'})$ such that

- (i) neither one is detectable for $s > t'$;
- (ii) all remain linearly independent in $\text{img}(H_r(X_{t'}) \rightarrow H_r(X_{s,t'}))$ for $t'' < s \leq t'$;
- (iii) all dies at t'' .

There exists an r -bar code $[t'', t')$ if there exists an element $x \in H_r(X_{t''})$ which is not detectable for $s < t''$ and detectable for $t'' \leq s < t'$ and dies at t' . The multiplicity of $[t'', t')$ is the maximal number of linearly independent elements in $H_r(X_{t''})$ such that

- (i) neither one is detectable for $s < t''$;
- (ii) all remain linearly independent in $\text{img}(H_r(X_{t''}) \rightarrow H_r(X_{t'',s}))$ for $t'' \leq s < t'$;
- (iii) all dies at t' .

There exists an r -bar code $[t'', t']$ if there exists an element $x \in H_r(X_{t''})$ which is not detectable for $s < t''$ or $s > t'$ and detectable for $t'' \leq s \leq t'$. The multiplicity of $[t'', t']$ is the maximal number of linearly independent elements in $H_r(X_{t''})$ such that

- (i) neither one is detectable for $s < t''$ or $s > t'$;
- (ii) all remain linearly independent in $\text{img}(H_r(X_{t''}) \rightarrow H_r(X_{t'',s}))$ for $t'' \leq s \leq t'$.

Lemma 4.5. *For a tame map, the set of r -bar codes for level persistence is finite. Any r -bar code is an interval of the form $[t_i, t_j]$ with $t_i \leq t_j$ critical values or (t_i, t_j) , $(t_i, t_j]$, $[t_i, t_j)$ with $t_i < t_j$, t_i, t_j critical values.*

Proof. Using Lemma 4.2 and Lemma 4.4, the proof is similar to Lemma 4.1. \square

Notation 4.2. Given a tame map $f : X \rightarrow \mathbb{R}$ with critical values $t_0 < \dots < t_N$, denote by

$BL_r(f) :=$ the number of all r -bar codes for level persistence
(with respect to r -th homology groups).

$N_r(t_i, t_j) :=$ the number of intervals (t_i, t_j) in $BL_r(f)$.

$N_r(t_i, t_j] :=$ the number of intervals $(t_i, t_j]$ in $BL_r(f)$.

$N_r[t_i, t_j) :=$ the number of intervals $[t_i, t_j)$ in $BL_r(f)$.

$N_r[t_i, t_j] :=$ the number of intervals $[t_i, t_j]$ in $BL_r(f)$.

Hence $\sharp BL_r(f) = N_r(t_i, t_j) + N_r(t_i, t_j] + N_r[t_i, t_j) + N_r[t_i, t_j]$.

In Figure 6, we indicate the bar codes both for sub-level and level persistence for some simple map in order to illustrate their differences and connections. The class consisting of the sum of two circles at level t is not detected on the right, but is detected at all levels on the left up to (but not including) the level t' .

Level persistence provides considerably more information than the sub-level persistence [3] and the bar codes for the sub-level persistence can be recovered from the bar codes for the level persistence. An r -bar code $[s, t)$ for level persistence contributes an r -bar code $[s, t)$ for sub-level persistence. An r -bar code $[s, t]$ for level persistence contributes an r -bar code $[s, \infty)$ for sub-level persistence. r -bar codes $(s, t]$ and (s, t) for level persistence contribute nothing to r -bar codes for sub-level persistence. An r -bar codes (s, t) for level persistence contributes an $r + 1$ -bar code $[t, \infty)$ for sub-level persistence. See Figure 6 and Lemma 4.6 below.

Lemma 4.6. Given a tame map $f : X \rightarrow \mathbb{R}$ with critical values $t_0 < t_1 < \dots < t_N$. We have

$$\begin{aligned} \mu_r(t_i, t_j) &= N_r[t_i, t_j) \\ \mu_r(t_i, \infty) &= \sum_{l=i}^N N_r[t_i, t_l] + \sum_{l=0}^{i-1} N_{r-1}(t_l, t_i) \end{aligned}$$

for any critical values $t_i < t_j$.

Proof. Item 1 follows from formulas (4.1) and (4.2).

Item 2 is more elaborate. One uses formula (4.2) which calculates $\mu_r(t_i, \infty)$ as $\mu_r(t_i, \infty) = \beta_r(t_i, \infty) - \beta_r(t_{i-1}, \infty)$. A calculation of $\beta_r(t_i, \infty)$ can be recovered from Corollary 3.4 in [3] which implies that this number is exactly

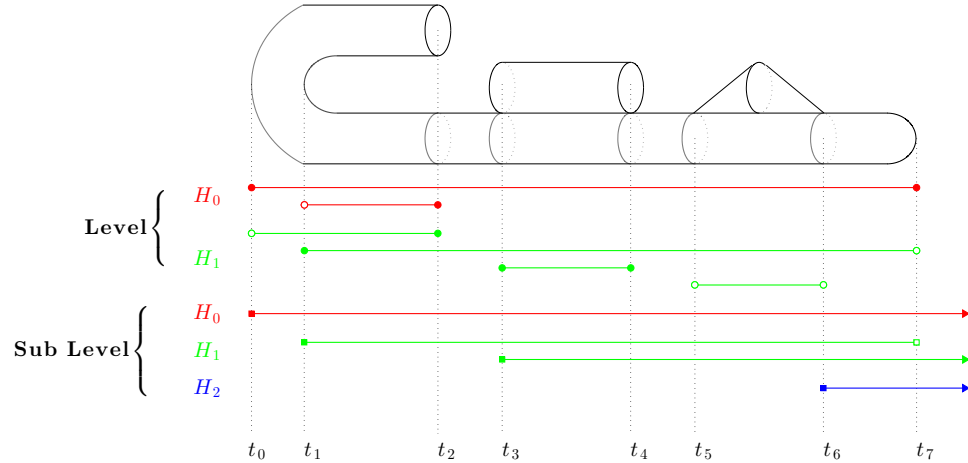


Figure 6: Bar codes for level and sub-level persistence.

the number of $(r-1)$ -bar codes of the form (t_l, t_i) , $l = 0, 1, \dots, i-1$ plus the number of r -bar codes of the form $[a, b]$ with $a \leq t_i$. Clearly a, b should be critical values. A different derivation can be achieved independently of [3]. \square

The bar codes for the level persistence can be also recovered from the bar codes for the sub-level persistence but from the bar codes of a collection of tame maps canonically associated to f . This will be described in the next subsection.

For this purpose one uses an alternative but equivalent way to describe the level persistence based on a different collection of numbers, referred below as relevant persistence numbers, $l_r, l_r^+, l_r^-, e_r, i_r$.

Definition 4.3. For a continuous map $f : X \rightarrow \mathbb{R}$ and $t'' \leq t \leq t'$, let $L_r(t) := H_r(X_t)$, $L_r^+(t; t') := \ker(H_r(X_t) \rightarrow H_r(X_{t,t'}))$, $L_r^-(t; t'') := \ker(H_r(X_t) \rightarrow H_r(X_{t'',t}))$ and

$$I_r(t, t') := \text{img}(H_r(X_t) \rightarrow H_r(X_{t,t'})) \cap \text{img}(H_r(X_{t'}) \rightarrow H_r(X_{t,t'})).$$

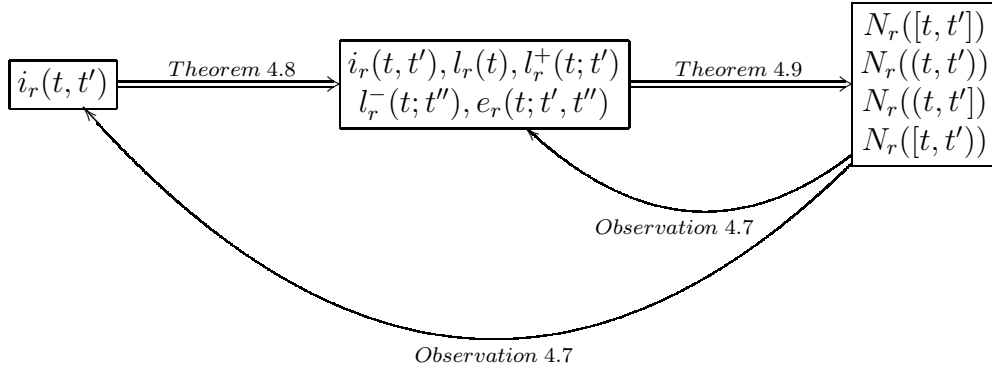
Define the **relevant level persistent numbers**

1. $l_r(t) := \dim L_r(t)$
2. $l_r^+(t; t') := \dim L_r^+(t; t')$
3. $l_r^-(t; t'') := \dim L_r^-(t; t'')$
4. $e_r(t; t', t'') := \dim(L_r^+(t; t') \cap L_r^-(t; t''))$

and

$$5. i_r(t, t') := \dim(I_r(t, t'))$$

The relation between these collections of numbers is illustrated in the diagram below.



The first four have geometric meaning the last ones (the fifth) are more technical. However the first four l_r, l_r^+, l_r^-, e_r can be derived from the last ones i_r .

One can derive all the numbers $l_r, l_r^+, l_r^-, e_r, i_r$ from the number of bar codes $N_r(t_i, t_j), N_r(t_i, t_j], N_r[t_i, t_j), N_r[t_i, t_j]$.

Observation 4.7. *For a tame map we can derive relevant level persistent numbers from the numbers N 's of bar codes for level persistence.*

Proof. For $t'' \leq t \leq t'$

1. $l_r(t)$ = number of intervals in $BL_r(f)$ which contain t ;
2. $i_r(t, t')$ = number of intervals in $BL_r(f)$ which contain $[t, t']$;

$$3. \quad l_r^+(t; t') = \sum_{t_i \leq t < t_j \leq t'} N_r[t_i, t_j) + \sum_{t_i < t < t_j \leq t'} N_r(t_i, t_j); \quad (4.3)$$

$$4. \quad l_r^-(t; t'') = \sum_{t'' \leq t_i < t \leq t_j} N_r(t_i, t_j] + \sum_{t'' \leq t_i < t < t_j} N_r(t_i, t_j); \quad (4.4)$$

$$5. \quad e_r(t; t', t'') = \sum_{t'' \leq t_i < t < t_j \leq t'} N_r(t_i, t_j). \quad (4.5)$$

□

Theorem 4.8. *For a tame map the numbers $i_r(t, t')$ determine the numbers $l_r(t)$, $l_r^+(t; t')$, $l_r^-(t; t'')$ and $e_r(t; t', t'')$.*

Proof.

$$l_r(t) = i_r(t, t).$$

$$N_r(t_k, t_j) = i_r(t'', t') - i_r(t_k, t') - i_r(t'', t_j) + i_r(t_k, t_j)$$

for any $t'' \leq t'$ such that $t_k < t'' < t_{k+1}$, $t_{j-1} < t' < t_j$.

$$N_r[t_k, t_j] = i_r(t_k, t_j) - i_r(t'', t_j) - i_r(t_k, t') + i_r(t'', t')$$

for any $t_{k-1} < t'' < t_k$ and $t_j < t' < t_{j+1}$.

$$N_r(t_k, t_j] = i_r(t'', t_j) - i_r(t_k, t_j) - i_r(t'', t') + i_r(t_k, t')$$

for any $t_k < t'' < t_{k+1}$ and $t_j < t' < t_{j+1}$.

$$N_r[t_k, t_j) = i_r(t_k, t') - i_r(t_k, t_j) - i_r(t'', t') + i_r(t'', t_j)$$

for any $t_{k-1} < t'' < t_k$ and $t_{j-1} < t' < t_j$.

Plug in equation (4.3), (4.4) and (4.5) we get $l_r^+(t; t')$, $l_r^-(t; t'')$ and $e_r(t; t', t'')$. \square

Theorem 4.9. *For a tame map we can get the r -bar codes for level persistence from relevant persistent numbers.*

Proof. First the numbers $N_r(t_k, t_j)$ can be calculated by the formula.

$$N_r(t_k, t_j) = e_r(t; t_j, t_k) - e_r(t; t_j, t_{k+1}) - e_r(t; t_{j-1}, t_k) + e_r(t; t_{j-1}, t_{k+1}), \quad (4.6)$$

for any $t_k < t < t_j$.

To determine the numbers $N_r(t_i, t_j]$, $N_r[t_i, t_j)$ and $N_r[t_i, t_j]$, we introduce the following auxiliary numbers

$n_r\{t_i, t_j\} :=$ the number of intervals in $BL_r(f)$ which intersect the levels X_{t_i} and X_{t_j} ;

$n_r\{t_i, t_j\} :=$ the number of intervals in $BL_r(f)$ which intersect the level X_{t_i} with open end at t_j ;

$n_r\{t_i, t_j] :=$ the number of intervals in $BL_r(f)$ which intersect the level X_{t_i} with closed end at t_j ;

$n_r(t_i, t_j\} :=$ the number of intervals in $BL_r(f)$ which intersect the level X_{t_j} with open end at t_i ;

$n_r[t_i, t_j\} :=$ the number of intervals in $BL_r(f)$ which intersect the level X_{t_j} with closed end at t_i .

The numbers $n_r\{t_i, t_j\}$, $n_r(t_i, t_j)$, $n_r\{t_i, t_j\}$ and $n_r[t_i, t_j]$ can be derived from the relevant persistent numbers as indicated below:

$$\begin{aligned} n_r\{t_i, t_j\} &= l_r^+(t_i; t_j) - l_r^+(t_i; t_{j-1}) \\ n_r(t_i, t_j) &= l_r^-(t_j; t_i) - l_r^-(t_j; t_{i+1}) \\ n_r\{t_i, t_j\} &= i_r(t_i, t_j) \\ n_r[t_i, t_j] &= n_r\{t_i, t_j\} - n_r\{t_{i-1}, t_j\} - n_r(t_{i-1}, t_j) \end{aligned}$$

With their help one derives

$$N_r(t_i, t_j] = n_r(t_i, t_j) - n_r(t_i, t_{j+1}) - N_r(t_i, t_{j+1}) \quad (4.7)$$

$$N_r[t_i, t_j) = n_r\{t_i, t_j\} - n_r\{t_{i-1}, t_j\} - N_r(t_{i-1}, t_j) \quad (4.8)$$

$$N_r[t_i, t_j] = n_r[t_i, t_j] - n_r[t_i, t_{j+1}] - N_r[t_i, t_{j+1}] \quad (4.9)$$

□

In order to calculate the relevant persistence numbers $l_r(t)$, $l_r^+(t; t')$, $l_r^-(t; t'')$ and $e_r(t; t', t'')$ we induce so called **positive and negative bar codes** in Subsection 4.4. Positive and negative bar codes are defined by sub-level persistence, so they can be calculated with the methods as described in Subsection 3.4.

We can also calculate $i_r(t, t')$ with similar method. Details will be given in Subsection 4.3.

Once we get all relevant numbers, we can get the bar codes for the level persistence providing an alternative to the Carson deSilva algorithm cf. [6] to calculate the level persistence bar codes as bar codes for Zigzag persistence.

4.3 The General Framework for the Calculation of Bar Codes for Level Persistence of Generic \mathbb{R} -valued Linear Maps on Simplicial Complexes

Definition 4.4. (*Generic \mathbb{R} -valued Linear Maps on Simplicial Complexes*)

Let X be a simplicial complex and $|X|$ its spatial realization. An \mathbb{R} -valued linear map on X is a continuous map $f : |X| \rightarrow \mathbb{R}$ whose restriction to each simplex is linear. Let X_0 be the set of vertices of X . f is called generic if $f : X_0 \rightarrow \mathbb{R}$ is injective. (See page 2, [3])

Theorem 4.10. *Generic \mathbb{R} -valued linear maps on simplicial complexes are tame.*⁷

Proof. Let f be a generic \mathbb{R} -valued linear map on a simplicial complex X with $N + 1$ vertices. Let $t_0 < \dots < t_N$ be the values of f on the vertices of X . Label vertices of X such that $f(x_i) = t_i$, $0 \leq i \leq N$.

(i) For $s < t$ with $t_{i-1} < s, t < t_i$ one produces a canonical homeomorphism $\theta_{s,t} : X_{s,t} \rightarrow X_t \times [s, t]$ which intertwines the map f with the second factor projection of $X_t \times [s, t]$ on $[s, t]$.

Given σ with $\sigma \cap X_{s,t} \neq \emptyset$. Each point $x \in X_{s,t} \cap \sigma$ lies on a unique straight line segment $l(x)$ which is the trajectory of $\nabla f|_\sigma$ w.r.t. the flat metric on σ . Denote by $L(l)$ (resp. $R(l)$) the intersection of this line with X_s (resp. X_t) and by $p_{s,t}(x) := L(l(x))$ (resp. $q_{s,t}(x) := R(l(x))$) the obviously continuous map $X_{s,t} \cap \sigma \rightarrow X_s \cap \sigma$ (resp. $X_{s,t} \cap \sigma \rightarrow X_t \cap \sigma$). Consider $\tilde{p}_{s,t} : X_{s,t} \cap \sigma \rightarrow X_s \cap \sigma \times [s, t]$ (resp. $\tilde{q}_{s,t} : X_{s,t} \cap \sigma \rightarrow X_t \cap \sigma \times [s, t]$) the map defined by $\tilde{p}_{s,t}(x) = (p_{s,t}(x), f(x))$ (resp. $\tilde{q}_{s,t}(x) = (q_{s,t}(x), f(x))$). Clearly $\tilde{p}_{s,t}$ and $\tilde{q}_{s,t}$ are homeomorphisms which intertwines the map f with the second factor projection of $X_t \times [s, t]$ on $[s, t]$.

For any σ_1 and σ_2 , $\tilde{p}_{s,t}$ and $\tilde{q}_{s,t}$ agree on the intersection $\sigma_1 \cap \sigma_2$. So we can define $\theta_{s,t}|_\sigma$, the restriction of $\theta_{s,t}$ on each σ , as $\tilde{q}_{s,t}$. In this way, one produces a canonical homeomorphism $\theta_{s,t} : X_{s,t} \rightarrow X_t \times [s, t]$ which intertwines the map f with the second factor projection of $X_t \times [s, t]$ on $[s, t]$.

Therefore, for any $t \neq t_0, t_1, \dots, t_N$, there exists $\epsilon > 0$ so that $f : f^{-1}(t - \epsilon, t + \epsilon) \rightarrow (t - \epsilon, t + \epsilon)$ and the second factor projection $X_t \times (t - \epsilon, t + \epsilon) \rightarrow (t - \epsilon, t + \epsilon)$ are fiber-wise homotopy equivalent.

(ii) For $s < t$ with $s = t_{i-1}, t < t_i$ and σ with $\sigma \cap X_{t_{i-1},t} \neq \emptyset$, it is easy to see that $X_{t_{i-1},t} \cap \sigma$ is the mapping cylinder of a canonical map $X_t \cap \sigma \rightarrow X_{t_{i-1}} \cap \sigma$ which prove that $X_{t_{i-1},t} \cap \sigma$ retracts by deformation on $X_{t_{i-1}} \cap \sigma$. For any σ_1 and σ_2 , the above deformation retractions agree on the intersection $\sigma_1 \cap \sigma_2$. So $X_{t_{i-1},t}$ retracts by deformation on $X_{t_{i-1}}$.

For $s < t$ with $t_{i-1} < s, t = t_i$, we can prove X_{s,t_i} retracts by deformation on X_{t_i} in the same way.

□

Let f be a generic \mathbb{R} -valued linear map on a simplicial complex X with $N + 1$ vertices. Let $t_0 < \dots < t_N$ be the values of f on the vertices of X . Since f is tame, we can define bar codes for level persistence of $f : |X| \rightarrow \mathbb{R}$. By Lemma 4.5, the bar codes for level persistence of f are intervals of the form $[t_i, t_j]$ with $t_i \leq t_j$ and (t_i, t_j) , $(t_i, t_j]$, $[t_i, t_j)$ with $t_i < t_j$.

⁷Any simplicial maps are tame, the arguments although similar are more elaborate to write down in the general case.

Lemma 4.11. *For the map $f : |X| \rightarrow \mathbb{R}$ defined above, the bar codes of the form (t_i, t_{i+1}) do not exist.*

Let f and X be defined as above. Let M_X be the boundary matrix of X which records the boundary map $\partial : X \rightarrow X$. Start with M_X and the values on vertices $t_0 < t_1 < \dots < t_N$, we can calculate the bar codes for level persistence with two methods:

Method 1, which involves only $i_r(t, t')$, and

Method 2, which involves $l_r^+(t; t')$, $l_r^-(t; t'')$, $e_r(t; t'', t')$ and $i_r(t, t')$.

Finally one use the pervious results to determine the bar codes. It appears that Method 1 is simpler, however we will see later that Method 2 is about 8 times faster than Method 1.

Definition 4.5. *Given any t , we get a polytopal complex*

$$X_t := \{ |\sigma| \cap |X|_t \mid \sigma \in X \}.$$

Given any $s < t$, we get aother polytopal complex

$$X_{s,t} := X_s \cup X_t \cup \{ |\sigma| \cap |X|_{s,t} \mid \sigma \in X \}.$$

For $s = t$, observe that $X_{s,t}$ is X_t .

Method 1 (with coefficients in \mathbb{Z}_2)

Given $s < t$, order cells in $X_{s,t}$ in the following way:

- (i) if $\tau \in X_s$, $\tau' \in X_t$ and $\tau'' \in X_{s,t}$, then $\tau \prec \tau' \prec \tau''$;
- (ii) if $\dim(\tau) < \dim(\tau')$, then $\tau \prec \tau'$.

We can write a matrix $M_{X_{s,t}}$ for the boundary map $\partial_{s,t} : X_{s,t} \rightarrow X_{s,t}$ with respect to the order of cells above. The ordering of the vertices provides orientation on simplices and then on the cells obtained as “cuts”.

From the reduced form $R(M_{X_{s,t}})$ calculated by Algorithm 4.1., we can read the $\dim(H_r(|X|_s) \rightarrow H_r(|X|_{s,t}))$ and $\dim(H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t}))$.

Let $n_1 = \sharp X_s$, $n_2 = \sharp X_t$ and $n = \sharp X_{s,t}$. A zero column j , $1 \leq j \leq n_1$, contributes a generator in the image of $H_r(|X|_s) \rightarrow H_r(|X|_{s,t})$, if there is no column $k > j$ such that $\text{low}(k) = j$. Therefore $\dim(H_r(|X|_s) \rightarrow H_r(|X|_{s,t})) = \text{number of such columns } j$.

A zero column j' , $1 \leq j' \leq n_1 + n_2$, contributes a generator in the image of $H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t})$, if there is no column $k > j'$ such that $\text{low}(k) = j'$. Therefore $\dim(H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t})) = \text{number of such columns } j'$.

Let's change the order of cells in $X_{s,t}$.

Given $s < t$, order cells in $X_{s,t}$ in the following way:

- (i) if $\tau \in X_t$, $\tau' \in X_s$ and $\tau'' \in X_{s,t}$, then $\tau \prec \tau' \prec \tau''$;
- (ii) if $\dim(\tau) < \dim(\tau')$, then $\tau \prec \tau'$.

We can write a matrix $M'_{X_{s,t}}$ for the boundary map $\partial_{s,t} : X_{s,t} \rightarrow X_{s,t}$ with respect to the order above.

From the reduced form $R(M'_{X_{s,t}})$ calculated by Algorithm 4.1., we can read the $\dim(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))$ and $\dim(H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t}))$ as describe above.

Since

$$I_r(s, t) = \text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t})) \cap \text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))$$

we have

$$\begin{aligned} & i_r(s, t) \\ = & \dim(I_r(s, t)) \\ = & \dim(\text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t})) \cap \text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))) \\ = & \dim(\text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t}))) + \dim(\text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))) \\ & - \dim(\text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t})) + \text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))). \end{aligned}$$

Since

$$\begin{aligned} & \text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t})) + \text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t})) \\ = & \text{img}((H_r(|X|_s) \oplus H_r(|X|_t)) \rightarrow H_r(|X|_{s,t})) \\ = & \text{img}(H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t})) \end{aligned}$$

we have

$$\begin{aligned} & i_r(s, t) \\ = & \dim(\text{img}(H_r(|X|_s) \rightarrow H_r(|X|_{s,t}))) + \dim(\text{img}(H_r(|X|_t) \rightarrow H_r(|X|_{s,t}))) \\ & - \dim(\text{img}(H_r(|X|_s \sqcup |X|_t) \rightarrow H_r(|X|_{s,t}))) \end{aligned}$$

for $s < t$.

For $s = t$, we have

$$i_r(s, s) = \dim(H_r(|X|_s))$$

and this can be read from the reduced form $R(M_{X_s})$ directly.

Once we get $i_r(s, t)$ for any $s \leq t$, we can get the numbers of all bar codes of level persistence according to Theorem 4.8 and Theorem 4.9.

Observe that in view of the tameness of f we only need finitely many $i_r(s, t)$. Let $t_{-1} = t_0 - 1$, $t_{N+1} = t_N + 1$ and $t_{i+1/2} = (t_i + t_{i+1})/2$ for $-1 \leq i \leq N$. It is enough to know $i_r(s, t)$ for $s, t \in \{t_{i/2} | -1 \leq i \leq 2N + 1\}$.

Therefore we need to calculate reduced forms of about $4N^2$ matrices in order to get r -bar codes of level persistence.

Method 2 (with coefficients in \mathbb{Z}_2)

Method 2 is based on the formulas of Theorem 4.9.

So called positive and negative bar codes are introduced in the next subsection in order to calculate $l_r(t_i)$, $l_r^+(t_i; t_j)$, $l_r^-(t_i; t_j)$ and $e_r(t_i; t_k, t_j)$ for $t_k \leq t_i \leq t_j$ efficiently. We need to calculate reduced forms of about $2N$ matrices in order to get $l_r(t_i)$, $l_r^+(t_i; t_j)$, $l_r^-(t_i; t_j)$ and $e_r(t_i; t_k, t_j)$ for $t_k \leq t_i \leq t_j$.

We still need to calculate $i_r(t_i; t_j)$ for $t_j < t_i$. We have

$$\begin{aligned} & i_r(t_i, t_j) \\ &= \dim(\text{img}(H_r(|X|_{t_i}) \rightarrow H_r(|X|_{t_i, t_j}))) + \dim(\text{img}(H_r(|X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j}))) \\ & \quad - \dim(\text{img}(H_r(|X|_{t_i} \sqcup |X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j}))). \end{aligned}$$

for $t_i < t_j$.

$\dim(\text{img}(H_r(|X|_{t_i} \sqcup |X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j})))$ can be calculated in the same way as described in Method 1.

$\dim(\text{img}(H_r(|X|_{t_i}) \rightarrow H_r(|X|_{t_i, t_j})))$ and $\dim(\text{img}(H_r(|X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j})))$ can be calculated at very little cost with the knowledge of $l_r(t_i)$, $l_r^+(t_i; t_j)$ and $l_r^-(t_i; t_j)$.

$$\begin{aligned} & \dim(\text{img}(H_r(|X|_{t_i}) \rightarrow H_r(|X|_{t_i, t_j}))) \\ &= \dim(H_r(|X|_{t_i})) - \dim(\ker(H_r(|X|_{t_i}) \rightarrow H_r(|X|_{t_i, t_j}))) \\ &= l_r(t_i) - l_r^+(t_i; t_j) \\ & \quad \dim(\text{img}(H_r(|X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j}))) \\ &= \dim(H_r(|X|_{t_j})) - \dim(\ker(H_r(|X|_{t_j}) \rightarrow H_r(|X|_{t_i, t_j}))) \\ &= l_r(t_j) - l_r^-(t_i; t_j) \end{aligned}$$

We need to calculate reduced forms of about $N^2/2$ matrices in order to get $i_r(t_i; t_j)$ for $t_j < t_i$.

Therefore we need to calculate reduced forms of about $N^2/2 + 2N$ matrices in order to get r -bar codes of level persistence.

4.4 Definition of Positive and Negative Bar Codes of a Tame Map

Suppose $f : X \rightarrow \mathbb{R}$ is a tame map with $t_1 < t_2 < \dots < t_N$ the critical values. Define $t_{m+1/2} = (t_m + t_{m+1})/2$, $1 \leq m \leq N-1$. Given $t_1 \leq u < s < t \leq t_N$, there exist k, i, j such that $t_{k-1} < u \leq t_k$, $t_i \leq s < t_{i+1}$, $t_j \leq t < t_{j+1}$, we have

Proposition 4.12.

$$H_r(X_s) \cong \begin{cases} H_r(X_{t_i}), & s = t_i \\ H_r(X_{t_{i+1/2}}), & t_i < s < t_{i+1} \end{cases}$$

$$\cong \begin{cases} \ker(H_r(X_s) \rightarrow H_r(X_{s,t})) \\ \begin{cases} \ker(H_r(X_{t_i}) \rightarrow H_r(X_{t_i,t_j})), & s = t_i \\ \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_{i+1/2},t_j})), & t_i < s < t_{i+1} \leq t \\ 0, & t_i < s < t < t_{i+1} \end{cases} \end{cases}$$

$$\cong \begin{cases} \ker(H_r(X_s) \rightarrow H_r(X_{u,s})) \\ \begin{cases} \ker(H_r(X_{t_i}) \rightarrow H_r(X_{t_k,t_i})), & s = t_i \\ \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_k,t_{i+1/2}})), & u \leq t_i < s < t_{i+1} \\ 0, & t_i < u < s < t_{i+1} \end{cases} \end{cases}$$

$$\cong \begin{cases} \ker(H_r(X_s) \rightarrow H_r(X_{s,t})) \cap \ker(H_r(X_s) \rightarrow H_r(X_{u,s})) \\ \begin{cases} \ker(H_r(X_{t_i}) \rightarrow H_r(X_{t_i,t_j})) \cap \ker(H_r(X_{t_i}) \rightarrow H_r(X_{t_k,t_i})), & s = t_i \\ \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_{i+1/2},t_j})) \cap \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_k,t_{i+1/2}})), & u \leq t_i < s < t_{i+1} \leq t \\ 0, & \text{otherwise} \end{cases} \end{cases}$$

Proof. We will verify the case $u \leq t_i < s < t_{i+1} \leq t$, other cases can be proved similarly.

The commutative diagram (all maps are inclusions)

$$\begin{array}{ccccc} X_{t_k,s} & \longleftarrow & X_s & \longrightarrow & X_{s,t_j} \\ & \searrow & & \searrow & \downarrow \\ & & X_{u,s} & & X_{s,t} \end{array}$$

induces

$$\begin{array}{ccccc} H_r(X_{t_k,s}) & \longleftarrow & H_r(X_s) & \longrightarrow & H_r(X_{s,t_j}) \\ & \searrow & & \searrow & \downarrow \\ & & H_r(X_{u,s}) & & H_r(X_{s,t}) \end{array}$$

We have two isomorphisms in the diagram above, since $X_{u,s}$ and $X_{s,t}$ retracts by deformation onto $X_{t_k,s}$ and X_{s,t_j} by Lemma 4.2. So we have

$$\begin{aligned} \ker(H_r(X_s) \rightarrow H_r(X_{s,t})) &= \ker(H_r(X_s) \rightarrow H_r(X_{s,t_j})) \\ \ker(H_r(X_s) \rightarrow H_r(X_{u,s})) &= \ker(H_r(X_s) \rightarrow H_r(X_{t_k,s})) \\ \text{and} \quad & \\ \ker(H_r(X_s) \rightarrow H_r(X_{s,t})) \cap \ker(H_r(X_s) \rightarrow H_r(X_{u,s})) & \\ = \ker(H_r(X_s) \rightarrow H_r(X_{s,t_j})) \cap \ker(H_r(X_s) \rightarrow H_r(X_{t_k,s})) & \quad (*) \end{aligned}$$

When $s = t_i$, this provides the verification of the isomorphism.

When $t_i < s < t_{i+1}$, wlog, suppose $t_{i+1/2} \leq s < t_{i+1}$. By Lemma 4.4, $X_{t_{i+1/2},s}$ retracts by deformation onto $X_{t_{i+1/2}}$ and X_s , the commutative diagram below (all maps are inclusions)

$$\begin{array}{ccccc}
 & & X_s & \longrightarrow & X_{s,t_j} \\
 & \swarrow & \downarrow & & \downarrow \\
 X_{t_k,s} & \longleftarrow & X_{t_{i+1/2},s} & \longrightarrow & X_{t_{i+1/2},t_j} \\
 \uparrow & & \uparrow & \nearrow & \\
 X_{t_k,t_{i+1/2}} & \longleftarrow & X_{t_{i+1/2}} & &
 \end{array}$$

induces

$$\begin{array}{ccccc}
 & & H_r(X_s) & \longrightarrow & H_r(X_{s,t_j}) \\
 & \swarrow & \downarrow \cong & & \downarrow \cong \\
 H_r(X_{t_k,s}) & \longleftarrow & H_r(X_{t_{i+1/2},s}) & \longrightarrow & H_r(X_{t_{i+1/2},t_j}) \\
 \cong \uparrow & & \cong \uparrow & \nearrow & \\
 H_r(X_{t_k,t_{i+1/2}}) & \longleftarrow & H_r(X_{t_{i+1/2}}) & &
 \end{array}$$

From the diagram above, we have

$$\begin{aligned}
 & H_r(X_s) \cong H_r(X_{t_{i+1/2}}) \\
 & \ker(H_r(X_s) \rightarrow H_r(X_{s,t_j})) \cong \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_{i+1/2},t_j})) \\
 & \ker(H_r(X_s) \rightarrow H_r(X_{t_k,s})) \cong \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_k,t_{i+1/2}})) \\
 & \ker(H_r(X_s) \rightarrow H_r(X_{s,t_j})) \cap \ker(H_r(X_s) \rightarrow H_r(X_{t_k,s})) \\
 \cong & \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_{i+1/2},t_j})) \cap \ker(H_r(X_{t_{i+1/2}}) \rightarrow H_r(X_{t_k,t_{i+1/2}}))
 \end{aligned}$$

Combined with (*), the isomorphisms follow as stated. \square

From Proposition 4.12, we can see that relevant level persistent numbers $l_r(s)$, $l_r^+(s; t)$, $l_r^-(s; u)$ and $e_r(s; t, u)$ for $t_1 \leq u \leq s \leq t \leq t_N$ are determined by $l_r(t_{(i+1)/2})$, $l_r^+(t_{(i+1)/2}; t_j)$, $l_r^-(t_{(i+1)/2}; t_k)$ and $e_r(t_{(i+1)/2}; t_j, t_k)$ for $1 \leq i \leq 2N - 1$, $1 \leq k \leq (i + 1)/2 \leq j \leq N$.

Proposition 4.13. *Let*

$$V_1 \xleftarrow{f_1} V_2 \xleftarrow{f_2} \cdots \xleftarrow{f_{i-1}} V_i \xrightarrow{f_i} V_{i+1} \xrightarrow{f_{i+1}} \cdots \xrightarrow{f_{N-1}} V_N$$

be a diagram of vector spaces V_1, \dots, V_N over a field κ .

Denote

$$g_k = \begin{cases} f_k \circ \cdots \circ f_{i-2} \circ f_{i-1} & k < i \\ id & k = i \\ f_k \circ \cdots \circ f_{i+1} \circ f_i & k > i \end{cases}$$

$$W_k = \text{img}(g_k : V_i \rightarrow V_k)$$

One can always choose a basis $\{v_1, \dots, v_n\}$ of V_i such that condition (P_1) below is satisfied for $1 \leq k \leq N$.

Condition (P_1) : Nonzero elements of $\{g_k(v_1), \dots, g_k(v_n)\}$ is a basis of W_k .

Proof. First we obtain a base $\{y_1, \dots, y_n\}$ of V_i such that P_1 is satisfied for all $k, i \leq k \leq N$ and a base $\{z_1, \dots, z_n\}$ of V_i such that P_1 is satisfied for all $k, 1 \leq k \leq i$.

(a) Let $\{x_1, \dots, x_n\}$ be a basis of V_i , then (P_1) is true for $k = i$. Suppose (P_1) is true for $i \leq k \leq j$, we can modify the basis so that (P_1) is also true for $i \leq k \leq j + 1$.

If $g_{j+1}(x_i) = 0$ for $1 \leq i \leq n$, then (P_1) is satisfied for $i \leq k \leq j + 1$; otherwise, there exist $1 \leq n_1 < n_2 < \cdots < n_r \leq n$ such that $g_{j+1}(x_{n_1}), \dots, g_{j+1}(x_{n_r})$ are nonzero and $g_{j+1}(x_l) = 0$ for $l \notin \{n_1, \dots, n_r\}$.

If $r = \dim(W_{j+1})$, then $\{g_{j+1}(x_{n_1}), \dots, g_{j+1}(x_{n_r})\}$ forms a basis of W_{j+1} ; otherwise, there exist $2 \leq r_1 \leq r$ such that $g_{j+1}(x_{n_1}), \dots, g_{j+1}(x_{n_{r_1-1}})$ are

linearly independent, but $g_{j+1}(x_{n_{r_1}}) = \sum_{l=1}^{r_1-1} a_l g_{j+1}(x_{n_l})$ for some $a_l \in \kappa$.

Let $x'_{n_{r_1}} = x_{n_{r_1}} - \sum_{l=1}^{r_1-1} a_l x_{n_l}$. Replace x_{n_r} by $x'_{n_{r_1}}$, $\{x_1, \dots, x'_{n_{r_1}}, \dots, x_n\}$

is still a basis of V_i . Since $g_{j+1}(x_{n_1}), \dots, g_{j+1}(x_{n_{r_1}})$ are nonzero, $g_k(x_{n_1}), \dots, g_k(x_{n_{r_1}})$ are nonzero for $i \leq k \leq j$. Since (P_1) is true for $\{x_1, \dots, x_n\}$ and $i \leq k \leq j$, $g_k(x_{n_1}), \dots, g_k(x_{n_{r_1}})$ are linearly independent for $i \leq k \leq j$. Hence $g_k(x'_{n_{r_1}}) \neq 0$ for $i \leq k \leq j$. Replace $g_k(x_{n_{r_1}})$ by $g_k(x'_{n_{r_1}})$, $\{g_k(x_1), \dots, g_k(x_{n_{r_1}}), \dots, g_k(x_n)\}$ and $\{g_k(x_1), \dots, g_k(x'_{n_{r_1}}), \dots, g_k(x_n)\}$ have the same number of nonzero elements for $i \leq k \leq j$, since $g_k(x_{n_{r_1}})$ and $g_k(x'_{n_{r_1}})$ are nonzero for $i \leq k \leq j$. In addition, those nonzero elements can be linearly represented by each other. So nonzero elements of $\{g_k(x_1), \dots, g_k(x'_{n_{r_1}}), \dots, g_k(x_n)\}$ is still a basis of W_k for $i \leq k \leq j$. Hence (P_1) is still true for $\{x_1, \dots, x'_{n_{r_1}}, \dots, x_n\}$ and $i \leq k \leq j$, but the number of nonzero elements of $\{g_{j+1}(x_1), \dots, g_{j+1}(x'_{n_{r_1}}), \dots, g_{j+1}(x_n)\}$ is one less than the number of nonzero elements of $\{g_{j+1}(x_1), \dots, g_{j+1}(x'_{n_{r_1}}), \dots, g_{j+1}(x_n)\}$ since $g_{j+1}(x'_{n_{r_1}}) = 0$ by definition and $g_{j+1}(x_{n_{r_1}}) \neq 0$.

By induction on the number of nonzero elements of $\{g_{j+1}(x_1), \dots, g_{j+1}(x_{n_{r_1}}), \dots, g_{j+1}(x_n)\}$, there exist a basis $\{x'_1, \dots, x'_n\}$ of V_i such that (P_1) is true for $i \leq k \leq j+1$.

Then by induction on j , we get a basis $\{y_1, \dots, y_n\}$ of V_i such that (P_1) is true for $i \leq k \leq N$.

Similarly, we can get a basis $\{z_1, \dots, z_n\}$ of V_i such that (P_1) is true for $1 \leq k \leq i$. To finalize the proof we need a definition.

(b) **Definition: Negative Life** Given $z_l \in V_i$ as in the basis above, if $g_k(z_l) \neq 0$ and $g_{k-1}(z_l) = 0$ for some $2 \leq k \leq i$ or $g_k(z_l) \neq 0$ for $k = 1$, define the *negative life* of z_l to be $i - k + 1$.

Order z_1, \dots, z_n according to their negative life increasingly.

Claim: For any given $1 \leq l \leq n-1$, if we replace z_{l+1} by $z'_{l+1} = z_{l+1} - \sum_{m=1}^l b_m z_m$ for any given $b_m \in \kappa$, $\{z_1, \dots, z'_{l+1}, \dots, z_n\}$ still satisfies (P_1) for $1 \leq k \leq i$.

Proof of the claim:

Since z_1, \dots, z_n are ordered according to their negative life increasingly, the negative life of z_m is less than the negative life of z_{l+1} for $1 \leq m \leq l$. Suppose $g_{k_1}(z_{l+1}) \neq 0$ and $g_{k_1-1}(z_{l+1}) = 0$ for some $2 \leq k_1 \leq i$, then we must have $g_k(z_m) = 0$ for all $1 \leq k \leq k_1-1$ and $1 \leq m \leq l+1$. Therefore $g_k(z_{l+1}) = g_k(z'_{l+1}) = 0$ for $1 \leq k \leq k_1-1$. Hence $\{g_k(z_1), \dots, g_k(z'_{l+1}), \dots, g_k(z_n)\}$ is the same as $\{g_k(z_1), \dots, g_k(z_n)\}$ for $1 \leq k \leq k_1-1$. So $\{z_1, \dots, z'_{l+1}, \dots, z_n\}$ satisfies (P_1) for $1 \leq k \leq k_1-1$.

Given $k_1 \leq k' \leq i$, we have $g_{k'}(z_{l+1}) \neq 0$. There exist $1 \leq l_1 \leq l$, such that $g_{k'}(z_m) = 0$ for $1 \leq m \leq l_1$ and $g_{k'}(z_m) \neq 0$ for $l_1+1 \leq m \leq l+1$.

Hence $g_{k'}(z'_{l+1}) = g_{k'}(z_{l+1}) - \sum_{m=l_1+1}^l b_m g_{k'}(z_m)$. Therefore nonzero elements of $\{g_{k'}(z_1), \dots, g_{k'}(z_n)\}$ and $\{g_{k'}(z_1), \dots, g_{k'}(z'_{l+1}), \dots, g_{k'}(z_n)\}$ can be linearly represented by each other. Since z_1, \dots, z_n satisfies (P_1) for $1 \leq k \leq i$, nonzero elements of $\{g_{k'}(z_1), \dots, g_{k'}(z_n)\}$ form a basis of $W_{k'}$. Then nonzero elements of $\{g_{k'}(z_1), \dots, g_{k'}(z'_{l+1}), \dots, g_{k'}(z_n)\}$ also form a basis of $W_{k'}$. So $\{z_1, \dots, z'_{l+1}, \dots, z_n\}$ satisfies (P_1) for $k_1 \leq k \leq i$.

The claim is proved.

Notice that in part(a), we replace $x_{n_{r_1}}$ by $x'_{n_{r_1}} = x_{n_{r_1}} - \sum_{l=1}^{r_1-1} a_l x_{n_l}$. We only need this type of replacement to modify a basis $\{x_1, \dots, x_n\}$ into a basis $\{y_1, \dots, y_n\}$ that satisfies (P_1) for $i \leq k \leq N$. And the replacement $x_{n_{r_1}}$ by

$x'_{n_{r_1}} = x_{n_{r_1}} - \sum_{l=1}^{r_1-1} a_l x_{n_l}$ is only a special case of the replacement in the above claim. Therefore, start from a basis $\{z_1, \dots, z_n\}$ of V_i such that (P_1) is true for $1 \leq k \leq i$ which is ordered according to their negative life increasingly, we can use the same process in part(a) to modify this basis until (P_1) is satisfied for $1 \leq k \leq N$.

□

We use the concept “positive bar code” because it is consistent with the sub-level persistence associated to the space $X_{[t_{(i+1)/2}, \infty)}$ and the map $f|_{X_{[t_{(i+1)/2}, \infty)}}$. Similarly, we use the concept “negative bar code” because it is consistent with the sub-level persistence associated to the space $X_{(-\infty, t_{(i+1)/2}]}$ and the map $f|_{X_{(-\infty, t_{(i+1)/2}]}}$.

Let $d_{(i+1)/2}$ denote $\dim(H_r(X_{t_{(i+1)/2}}))$. Choose a basis $\{c_1, \dots, c_{d_{(i+1)/2}}\}$ of $H_r(X_{t_{(i+1)/2}})$ such that for any $k < (i+1)/2$ and $j > (i+1)/2$, nonzero elements of $\{\sigma_k(c_1), \dots, \sigma_k(c_{d_{(i+1)/2}})\}$ is always a basis of $\text{img}(\sigma_k : H_r(X_{t_{(i+1)/2}}) \rightarrow H_r(X_{t_k, t_{(i+1)/2}}))$ and nonzero elements of $\{\tau_j(c_1), \dots, \tau_j(c_{d_{(i+1)/2}})\}$ is always a basis of $\text{img}(\tau_j : H_r(X_{t_{(i+1)/2}}) \rightarrow H_r(X_{t_{(i+1)/2}, t_j}))$, where σ_k and τ_j are induced by inclusion. This is possible by Proposition 4.13.

Definition 4.6. *Positive and Negative Bar Codes*

(i) The set of r -positive bar codes $B_r^+(f; t_{(i+1)/2})$ of $X_{t_{(i+1)/2}, \infty}$ is a collection of $d_{(i+1)/2}$ intervals described as follows. For any generator $c_a \in \{c_1, \dots, c_{d_{(i+1)/2}}\}$, if c_a dies upward at $t_{j'}$ ($t_{j'} > t_{(i+1)/2}$), then we have an interval $\langle t_{(i+1)/2}, t_{j'} \rangle$ in $B_r^+(f; t_{(i+1)/2})$; otherwise, if c_a never dies upward, we have an interval $\langle t_{(i+1)/2}, \infty \rangle$ in $B_r^+(f; t_{(i+1)/2})$.

(ii) The set of r -negative bar codes $B_r^-(f; t_{(i+1)/2})$ of $X_{-\infty, t_{(i+1)/2}}$ is a collection of $d_{(i+1)/2}$ intervals described as follows. For any generator $c_b \in \{c_1, \dots, c_{d_{(i+1)/2}}\}$, if c_b dies downward at $t_{k'}$ ($t_{k'} < t_{(i+1)/2}$), then we have an interval $\langle t_{k'}, t_{(i+1)/2} \rangle$ in $B_r^-(f; t_{(i+1)/2})$; otherwise, if c_b never dies downward, we have an interval $\langle -\infty, t_{(i+1)/2} \rangle$ in $B_r^-(f; t_{(i+1)/2})$.

(iii) Each generator c_k correspond to an interval $\langle a_k, t_{(i+1)/2} \rangle$ in the negative bar code $B_r^-(f, t_{(i+1)/2})$ and an interval $\langle t_{(i+1)/2}, b_k \rangle$ in the positive bar code $B_r^+(f, t_{(i+1)/2})$. Therefore each generator c_k defines an interval (a_k, b_k) . Define $B_r(f, t_{(i+1)/2})$ as the collection of the above intervals $\{(a_k, b_k) | 1 \leq k \leq d_{(i+1)/2}\}$.

We have introduced these additional numbers given the fact they can be efficiently calculated as indicated in Subsection 4.5 and they are equivalent to the relevant persistent numbers $l_r(t_{(i+1)/2})$, $l_r^+(t_{(i+1)/2}; t_j)$, $l_r^-(t_{(i+1)/2}; t_k)$ and $e_r(t_{(i+1)/2}; t_j, t_k)$ as indicated below.

We can get the relevant persistent numbers $l_r(t_{(i+1)/2})$, $l_r^+(t_{(i+1)/2}; t_j)$, $l_r^-(t_{(i+1)/2}; t_k)$ and $e_r(t_{(i+1)/2}; t_j, t_k)$ from the positive and negative bar codes:

$$l_r(t_{(i+1)/2}) = \dim H_r(X_{t_{(i+1)/2}}) = |B_r^+(f, t_{(i+1)/2})| = |B_r^-(f, t_{(i+1)/2})| \quad (4.10)$$

$$l_r^+(t_{(i+1)/2}; t_j) = l_r(t_{(i+1)/2}) - \sharp\{\langle t_{(i+1)/2}, b \rangle \in B_r^+(f, t_{(i+1)/2}) \mid b > t_j\} \quad (4.11)$$

$$l_r^-(t_{(i+1)/2}; t_k) = l_r(t_{(i+1)/2}) - \sharp\{\langle a, t_{(i+1)/2} \rangle \in B_r^-(f, t_{(i+1)/2}) \mid a < t_k\} \quad (4.12)$$

$$\begin{aligned} & e_r(t_{(i+1)/2}; t_j, t_k) \\ = & l_r(t_{(i+1)/2}) - (l_r(t_{(i+1)/2}) - l_r^+(t_{(i+1)/2}; t_j)) - (l_r(t_{(i+1)/2}) - l_r^-(t_{(i+1)/2}; t_k)) \\ & + \sharp\{\langle a, b \rangle \in B_r(f, t_{(i+1)/2}) \mid a < t_k, b > t_j\} \\ = & l_r^+(t_{(i+1)/2}; t_j) + l_r^-(t_{(i+1)/2}; t_k) - l_r(t_{(i+1)/2}) \\ & + \sharp\{\langle a, b \rangle \in B_r(f, t_{(i+1)/2}) \mid a < t_k, b > t_j\} \end{aligned} \quad (4.13)$$

We can also get the positive and negative bar codes from the relevant persistent numbers $l_r(t_{(i+1)/2})$, $l_r^+(t_{(i+1)/2}; t_j)$, $l_r^-(t_{(i+1)/2}; t_k)$ and $e_r(t_{(i+1)/2}; t_j, t_k)$:

$$\begin{aligned} \sharp\{\langle t_{(i+1)/2}, t_j \rangle \in B_r^+(f, t_{(i+1)/2})\} &= l_r^+(t_{(i+1)/2}; t_j) - l_r^+(t_{(i+1)/2}; t_{j-1}) \\ \sharp\{\langle t_k, t_{(i+1)/2} \rangle \in B_r^-(f, t_{(i+1)/2})\} &= l_r^-(t_{(i+1)/2}; t_k) - l_r^-(t_{(i+1)/2}; t_{k+1}) \\ \sharp\{\langle t_k, t_j \rangle \in B_r(f, t_{(i+1)/2})\} &= e(t_{(i+1)/2}; t_j, t_k) - e(t_{(i+1)/2}; t_j, t_{k+1}) \\ &\quad - e(t_{(i+1)/2}; t_{j-1}, t_k) + e(t_{(i+1)/2}; t_{j-1}, t_{k+1}) \end{aligned}$$

Therefore, positive and negative bar codes and relevant persistent numbers $l_r(t_{(i+1)/2})$, $l_r^+(t_{(i+1)/2}; t_j)$, $l_r^-(t_{(i+1)/2}; t_k)$ and $e_r(t_{(i+1)/2}; t_j, t_k)$ are equivalent.

4.5 Computing Positive and Negative Bar codes of Generic \mathbb{R} -valued Linear Maps on Simplicial Complexes

Let X be a simplicial complex with N vertices. Let $f : |X| \rightarrow \mathbb{R}$ be a generic linear map on X , with $t_1 < \dots < t_N$ being values on vertices x_1, \dots, x_N of X . According to Method 2 in Subsection 4.3, we need to calculate $l_r(t_i)$, $l_r^+(t_i; t_j)$, $l_r^-(t_i; t_j)$ and $e_r(t_i; t_k, t_j)$ for $1 \leq k \leq i \leq j \leq N$. So we only need to calculate $B_r^+(|X|; t_i)$ and $B_r^-(|X|; t_i)$ for $1 \leq i \leq N$, then plug in (4.10)-(4.13).

We'll describe how to compute positive bar codes in details and negative bar codes briefly at the end.

In order to get $B_r^+(|X|; t_i)$, we have to consider the filtration below

$$|X|_{t_i} \subseteq |X|_{t_i, t_{i+1}} \subseteq \dots \subseteq |X|_{t_i, t_N}$$

(see Definition 4.1 and 4.2, [3]).

It is not economical to compute the positive bar codes from this filtration directly. Instead, we'll consider an equivalent filtration.

Suppose that the vertices are ordered from 1 to N .

Definition 4.7. Given a j -simplex $\sigma = [x_{n_0}, \dots, x_{n_j}]$ of X , $1 \leq n_0 < \dots < n_j \leq N$, define $t_{\min}(|\sigma|) = t_{n_0}$, $t_{\max}(|\sigma|) = t_{n_j}$. For $t_{n_0} < s < t_{n_j}$, let $|\sigma|_s$ denote $|\sigma| \cap |X|_s$, $|\sigma|_{s,\infty}$ denote $|\sigma| \cap |X|_{s,\infty}$, $|\sigma|_{-\infty,s}$ denote $|\sigma| \cap |X|_{-\infty,s}$.

Extend t_{\min} and t_{\max} to $|\sigma|_s$, $|\sigma|_{s,\infty}$ and $|\sigma|_{-\infty,s}$, we have

$$\begin{aligned} t_{\min}(|\sigma|_s) &= t_{\max}(|\sigma|_s) = s, \\ t_{\min}(|\sigma|_{s,\infty}) &= s, & t_{\max}(|\sigma|_{s,\infty}) &= t_{n_j}, \\ t_{\min}(|\sigma|_{-\infty,s}) &= t_{n_0}, & t_{\max}(|\sigma|_{-\infty,s}) &= s. \end{aligned}$$

Note: It is easy to see that all cells in $X_{t_i,\infty}$ are $\{|\sigma| \mid t_{\min}(|\sigma|) \geq t_i\}$, $\{|\sigma|_{t_i} \mid t_{\min}(|\sigma|) < t_i < t_{\max}(|\sigma|)\}$ or $\{|\sigma|_{t_i,\infty} \mid t_{\min}(|\sigma|) < t_i < t_{\max}(|\sigma|)\}$.

Definition 4.8. Given s, t such that $t_1 \leq s < t \leq t_N$, define $Y_{s,t}$ to be a cell complex consist of cells in $X_{s,\infty}$ that are contained in the space $X_{-\infty,t}$. In another words, the cells in $Y_{s,t}$ are the cells \mathbf{c} in $X_{s,\infty}$ which satisfy $t_{\max}(\mathbf{c}) \leq t$.

Lemma 4.14. Let $\sigma = [x_0, \dots, x_n]$ be an n -simplex and $f : |\sigma| \rightarrow \mathbb{R}$ be a generic linear map on it. Suppose $f(x_0) < f(x_1) < \dots < f(x_n)$, let t_j denote $f(x_j)$, $0 \leq j \leq n$. Given $t_0 < s < t_n$, suppose $t_i \leq s < t_{i+1}$ for some $0 \leq i \leq n-1$. Let $|\sigma|_{-\infty,s}$ denote $f^{-1}(-\infty, s]$, then $|\sigma|_{-\infty,s}$ retracts by deformation onto $[x_0, \dots, x_i]$.

Proof. Given $x \in |\sigma|_{-\infty,s}$, write x as $\sum_{j=0}^n a_j x_j$, where $0 \leq a_j \leq 1, \sum_{j=0}^n a_j = 1$.

We have $\sum_{j=0}^i a_j > 0$, otherwise $a_j = 0, 0 \leq j \leq i$, $x = \sum_{j=i+1}^n a_j x_j$, $f(x) =$

$\sum_{j=i+1}^n a_j t_j > s$, since $t_j > s$ for $j \geq i+1$ and $\sum_{j=i+1}^n a_j = 1$. This is a

contradiction with $x \in |\sigma|_{-\infty,s}$.

Define $g_0 : |\sigma|_{-\infty,s} \rightarrow |\sigma|_{-\infty,s}$ to be the identity map.

Define

$$\begin{aligned} g_1 : |\sigma|_{-\infty,s} &\rightarrow |\sigma|_{-\infty,s} \\ x = \sum_{j=0}^n a_j x_j &\mapsto \frac{1}{\sum_{j=0}^i a_j} \sum_{j=0}^i a_j x_j \end{aligned}$$

g_1 is a well-defined continuous map and $g_1|_{|[x_0, \dots, x_i]|}$ is the identity.

Define

$$\begin{aligned} G : |\sigma|_{-\infty, s} \times I &\rightarrow |\sigma|_{-\infty, s} \\ (x, \tau) &\mapsto (1 - \tau)g_0(x) + \tau g_1(x) \end{aligned}$$

G is a deformation retraction from $|\sigma|_{-\infty, s}$ onto $|[x_0, \dots, x_i]|$ and it is canonical.

□

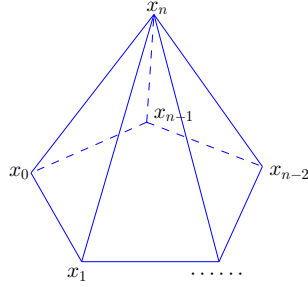


Figure 7: A cone with apex x_n and base a convex cell with vertices x_0, \dots, x_{n-1} .

Lemma 4.15. *Let σ be a cone with apex x_n and base a convex cell with vertices x_0, \dots, x_{n-1} . See Figure 7. Let f be a linear map on $|\sigma|$ such that $f(x_0) \leq f(x_1) \leq \dots \leq f(x_{n-1}) < f(x_n)$. For $f(x_{n-1}) \leq s < f(x_n)$, let $|\sigma|_{-\infty, s}$ denote $f^{-1}(-\infty, s]$, $|\sigma|_{-\infty, s}$ retracts by deformation onto the base, again by a canonical retraction.*

Proof. Since σ is a cone, for any point $x \in |\sigma| (x \neq x_n)$, there exists a unique x' in the base such that $x = (1 - a)x' + ax_n$, for some $0 \leq a \leq 1$.

Define $g_0 : |\sigma|_{-\infty, s} \rightarrow |\sigma|_{-\infty, s}$ to be the identity map.

Define

$$\begin{aligned} g_1 : |\sigma|_{-\infty, s} &\rightarrow |\sigma|_{-\infty, s} \\ x &\mapsto x' \end{aligned}$$

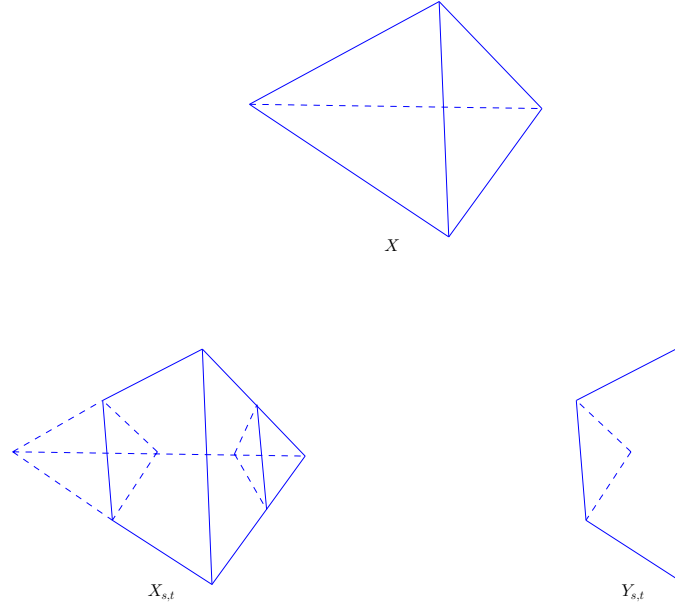
g_1 is a well-defined continuous map and $g_1|_{\text{the base}}$ is the identity.

Define

$$\begin{aligned} G : |\sigma|_{-\infty, s} \times I &\rightarrow |\sigma|_{-\infty, s} \\ (x, \tau) &\mapsto (1 - \tau)g_0(x) + \tau g_1(x) \end{aligned}$$

G is a deformation retraction from $|\sigma|_{-\infty, s}$ onto the base.

□

Figure 8: X , $X_{s,t}$ and $Y_{s,t}$.

Lemma 4.16. *Let $X = [x_0, x_1, \dots, x_n]$ be an n -simplex and f be a generic linear map on $|X|$. Suppose $f(x_0) < f(x_1) < \dots < f(x_n)$, let t_j denote $f(x_j)$, $0 \leq j \leq n$. Given $t_0 < s < t < t_n$, let $Y_{s,t}$ be as in Definition 4.8, then $X_{s,t}$ retracts by deformation onto $Y_{s,t}$. See Figure 8.*

Proof. Suppose $t_i < s \leq t_{i+1}$, $t_j \leq t < t_{j+1}$, $0 \leq i \leq j \leq n-1$. When $i = j$, $Y_{s,t} = X_s$ and $X_{s,t}$ retracts by deformation onto X_s by Lemma 4.4. When $i < j$, $X_{s,t}$ retracts by deformation onto X_{s,t_j} . Since $Y_{s,t} = Y_{s,t_j}$, we only need to prove $X_{s,t_j} \searrow Y_{s,t_j}$.

Note: “ \searrow ” means “retracts by deformation onto” here and after.

$X_{s,\infty}$ is an n -dim convex cell, which can be viewed as union of n -dim cones with apex x_n and base $(n-1)$ -dim faces of $X_{s,\infty}$ which do not contain x_n .

By Lemma 4.15, the intersection of each of the cones and $X_{s,t_{n-1}}$ retracts by deformation to the base of that cone.

Since the deformation retractions are compatible on the intersection of the cones, $X_{s,t_{n-1}} = \text{union of intersections of the cones and } X_{s,t_{n-1}} \searrow \text{union of bases of the cones} = (n-1)\text{-dim faces of } X_{s,\infty} \text{ which do not contain } x_n = Y_{s,t_{n-1}}$.

Since $Y_{s,t_{n-1}}$ are union of $(n-1)$ -dim faces of $X_{s,\infty}$ which do not contain

x_n , $Y_{s,t_{n-1}} = Y_{s,t_{n-2}} \cup$ cells of $Y_{s,t_{n-1}}$ which contain x_{n-1} .

Let σ_1, σ_2 be $(n-1)$ -dim cells of $Y_{s,t_{n-1}}$ which contain x_{n-1} , $\sigma_1 \cap \sigma_2 \neq \emptyset$. σ_1 can be viewed as union of $(n-1)$ -dim cones with apex x_{n-1} .

By Lemma 4.15, $\sigma_1 \cap X_{s,t_{n-2}} \searrow$ a cell sub complex of $Y_{s,t_{n-2}}$. Similarly, $\sigma_2 \cap X_{s,t_{n-2}} \searrow$ a cell sub complex of $Y_{s,t_{n-2}}$.

The deformation retractions on $\sigma_1 \cap X_{s,t_{n-2}}$ and $\sigma_2 \cap X_{s,t_{n-2}}$ induce the same deformation retractions on $\sigma_1 \cap \sigma_2 \cap X_{s,t_{n-2}}$. Therefore $Y_{s,t_{n-1}} \cap X_{s,t_{n-2}} \searrow Y_{s,t_{n-2}}$.

$X_{s,t_{n-1}} \searrow Y_{s,t_{n-1}}$ implies $X_{s,t_{n-2}} \searrow Y_{s,t_{n-1}} \cap X_{s,t_{n-2}}$. So $X_{s,t_{n-2}} \searrow Y_{s,t_{n-2}}$. Suppose $X_{s,t_k} \searrow Y_{s,t_k}$, we have $X_{s,t_{k-1}} \searrow Y_{s,t_k} \cap X_{s,t_{k-1}}$.

Let σ be a cell of Y_{s,t_k} which contains x_k . σ can be viewed as union of cones with apex x_k and base convex cells in $Y_{s,t_{k-1}}$.

By Lemma 4.15, $\sigma \cap X_{s,t_{k-1}}$ retracts by deformation onto a cell sub complex of $Y_{s,t_{k-1}}$.

The above deformation retractions are coherent on the boundary of the cones, so they induce a deformation retraction from $Y_{s,t_k} \cap X_{s,t_{k-1}}$ onto $Y_{s,t_{k-1}}$.

By induction $X_{s,t_j} \searrow Y_{s,t_j}$.

□

Proposition 4.17. $X_{s,t}$ retracts by deformation onto $Y_{s,t}$.

Proof. Four types of simplices of X have nonempty intersection with space $|X|_{s,t}$:

- (1) $\{\sigma \in X \mid t_{\min}(\sigma) < s, s < t_{\max}(\sigma) \leq t\}$
- (2) $\{\sigma \in X \mid t_{\min}(\sigma) < s, t_{\max}(\sigma) > t\}$
- (3) $\{\sigma \in X \mid t_{\min}(\sigma) \geq s, t_{\max}(\sigma) \leq t\}$
- (4) $\{\sigma \in X \mid s \leq t_{\min}(\sigma) \leq t, t_{\max}(\sigma) > t\}$

Therefore, there are four types of cells of $X_{s,t}$:

- (1) $\{\sigma_{s,\infty} \mid \sigma \in X, t_{\min}(\sigma) < s, s < t_{\max}(\sigma) \leq t\}$
- (2) $\{\sigma_{s,t} \mid \sigma \in X, t_{\min}(\sigma) < s, t_{\max}(\sigma) > t\}$
- (3) $\{\sigma \in X \mid t_{\min}(\sigma) \geq s, t_{\max}(\sigma) \leq t\}$
- (4) $\{\sigma_{-\infty,t} \mid \sigma \in X, s \leq t_{\min}(\sigma) \leq t, t_{\max}(\sigma) > t\}$

There are five types of cells of $X_{s,\infty}$ with respect to t ($s < t$):

- (1) $\{\sigma_{s,\infty} \mid \sigma \in X, t_{\min}(\sigma) < s, s < t_{\max}(\sigma) \leq t\}$
- (2) $\{\sigma_{s,\infty} \mid \sigma \in X, t_{\min}(\sigma) < s, t_{\max}(\sigma) > t\}$
- (3) $\{\sigma \in X \mid t_{\min}(\sigma) \geq s, t_{\max}(\sigma) \leq t\}$
- (4) $\{\sigma \in X \mid s \leq t_{\min}(\sigma) \leq t, t_{\max}(\sigma) > t\}$
- (5) $\{\sigma \in X \mid t_{\min}(\sigma) > t\}$

There are four types of cell complexes of $Y_{s,t}$:

- (1) $\{\sigma_{s,\infty} \mid \sigma \in X, t_{\min}(\sigma) < s, s < t_{\max}(\sigma) \leq t\}$

(2) $\{\text{subcomplexes of } \sigma_{s,\infty} \text{ which do not contain } x_{m+1}, \dots, x_n \mid \sigma = [x_0, \dots, x_n] \in X, t_{\min}(\sigma) < s, t_{\max}(\sigma) > t, f([x_0, \dots, x_i]) \leq t, 0 \leq i \leq m, f([x_0, \dots, x_i]) > t, m+1 \leq i \leq n\}$

(3) $\{\sigma \in X \mid t_{\min}(\sigma) \geq s, t_{\max}(\sigma) \leq t\}$

(4) $\{[x_0, \dots, x_m] \mid \sigma = [x_0, \dots, x_n] \in X, s \leq t_{\min}(\sigma) \leq t, t_{\max}(\sigma) > t, f([x_0, \dots, x_i]) \leq t, 0 \leq i \leq m, f([x_0, \dots, x_i]) > t, m+1 \leq i \leq n\}$

Notice that type (1) and (3) cells of $X_{s,t}$ and $Y_{s,t}$ are the same. By Lemma 4.15 type (2) cells of $X_{s,t}$ retract by deformation onto type (2) cell complexes of $Y_{s,t}$. By Lemma 4.14 type (4) cells of $X_{s,t}$ retract by deformation onto type (4) simplices of $Y_{s,t}$. □

By Proposition 4.17, the filtration

$$Y_{t_i} \subseteq Y_{t_i, t_{i+1}} \subseteq \dots \subseteq Y_{t_i, t_N}$$

provides the same bar codes as the filtration

$$X_{t_i} \subseteq X_{t_i, t_{i+1}} \subseteq \dots \subseteq X_{t_i, t_N}.$$

The advantage to consider Y_{t_i, t_j} instead of X_{t_i, t_j} is that Y_{t_i, t_j} keeps more simplices from X and generates less “new” cells, which makes representation of cells and construction of boundary matrices more economical.

Note: From now on, we'll only consider homology groups with \mathbb{Z}_2 coefficients. $H_r(X; \mathbb{Z}_2)$ will be simply written as $H_r(X)$.

Let X be a simplicial complex containing N vertices x_1, x_2, \dots, x_N . Let n denote $\dim X$, m_j denote number of j -simplices of X , $0 \leq j \leq n$. Let f be a generic linear map on X . For simplicity, let $f(x_i) = i$, $1 \leq i \leq N$. Represent j -simplex of X by (a_0, a_1, \dots, a_j) , $1 \leq a_0 < a_1 < \dots < a_j \leq N$.

All information on the above simplicial complex X can be stored in $n+1$ matrices named $\text{Simplex}_0, \text{Simplex}_1, \dots, \text{Simplex}_n$. The matrix Simplex_j is an $m_j \times (j+1)$ matrix, which stores j -simplices of X as rows in lexicographic order.

Example 4.1.

$$\text{Simplex}_0 = \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{pmatrix} \quad \text{Simplex}_1 = \begin{pmatrix} 12 \\ 13 \\ 23 \\ 24 \\ 25 \\ 35 \\ 36 \\ 45 \\ 56 \end{pmatrix} \quad \text{Simplex}_2 = (2 \ 3 \ 5)$$

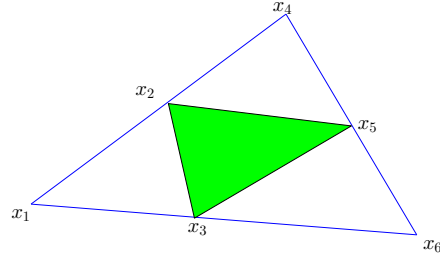


Figure 9: Example 4.1.

Definition 4.9. Initial Order

Order all simplices of X first according to dimension increasingly, then according to lexicographic order, we call this order as initial order of simplices of X .

Example 4.2. The initial order for Example 4.1 is $1 < 2 < 3 < 4 < 5 < 6 < 12 < 13 < 23 < 24 < 25 < 35 < 36 < 45 < 56 < 235$.

Let $\sigma_1, \sigma_2, \dots, \sigma_m$ be simplices of X in initial order, where $m = \sum_{j=0}^n m_j$.

We have a boundary matrix of \mathbb{Z}_2 coefficients

$$\partial := \begin{matrix} & \sigma_1 & \sigma_2 & \cdots & \sigma_j & \cdots & \sigma_m \\ \begin{matrix} \sigma_1 \\ \sigma_2 \\ \vdots \\ \sigma_i \\ \vdots \\ \sigma_m \end{matrix} & \left(\begin{array}{cccccc} & & & & & \\ & & & & & \\ & & & & & \\ & & & & \partial_{ij} & \\ & & & & & \\ & & & & & \end{array} \right) \end{matrix}$$

$$\partial_{ij} = \begin{cases} 1 & \text{if } \sigma_i \text{ is a codimension-1 face of } \sigma_j \\ 0 & \text{otherwise} \end{cases}$$

We'll compute positive and negative bar codes from this boundary matrix ∂ .

4.6 Computing of Positive Bar codes $B_r^+(f; i)$

We'll describe how to compute positive bar codes $B_r^+(f; i)$ of $X_{i,\infty}$ first.

There are five classes of cells in $X_{i,\infty}$:

$$P_1 = \{i\}$$

$$P_2 = \{\sigma_i \mid \sigma \in X, t_{\min}(\sigma) < i \text{ and } t_{\max}(\sigma) > i\}$$

$$P_3 = \{\sigma_{i,\infty} \mid \sigma \in X, t_{\min}(\sigma) < i \text{ and } t_{\max}(\sigma) > i\}$$

$$P_4 = \{\sigma \in X \mid \dim(\sigma) > 0 \text{ and } t_{\min}(\sigma) = i\}$$

$$P_5 = \{\sigma \in X \mid t_{\min}(\sigma) > i\}$$

Order cells in $X_{i,\infty}$ first according to these five classes, then according to initial order within each class. We can easily construct a boundary matrix ∂_i^+ from ∂ according to this order:

$$\partial_i^+ := \begin{matrix} & \begin{matrix} P_1 & P_2 & P_3 & P_4 & P_5 \end{matrix} \\ \begin{matrix} P_1 \\ P_2 \\ P_3 \\ P_4 \\ P_5 \end{matrix} & \left(\begin{array}{ccccc} & \boxed{Q_1} & & \boxed{Q_3} & \\ \boxed{M_1} & & \boxed{Q_2} & & \\ & & & & \\ & & \boxed{M_2} & & \\ & & & \boxed{M_3} & \boxed{M_4} \end{array} \right) \end{matrix}$$

where $M_i (1 \leq i \leq 4)$ are correspondent submatrices of ∂ , elements in $Q_1 = 1$ iff $\dim \sigma = 2$ and correspondent column in M_1 contains only one “1”, Q_2 is identity matrix, elements in $Q_3 = 1$ iff $\dim \sigma = 1$.

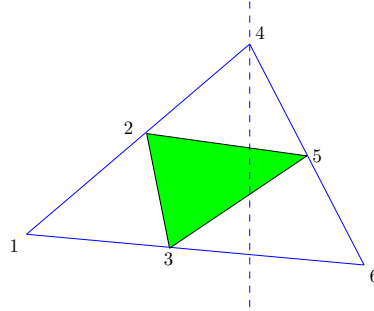


Figure 10: Example 4.3.

Example 4.3.

Let X be the same as in Example 4.1 and $i = 4$, we have

$$P_1 = \{4\}$$

$$P_2 = \{25|_4, 35|_4, 36|_4, 235|_4\}$$

$$P_3 = \{25|_{4,\infty}, 35|_{4,\infty}, 36|_{4,\infty}, 235|_{4,\infty}\}$$

$$P_4 = \{45\}$$

$$P_5 = \{5, 6, 56\}$$

$$\partial_4^+ := \begin{array}{c} P_1 \quad 4 \quad \overbrace{25|_4 \quad 35|_4 \quad 36|_4 \quad 235|_4}^{P_2} \quad \overbrace{25|_{4,\infty} \quad 35|_{4,\infty} \quad 36|_{4,\infty} \quad 235|_{4,\infty}}^{P_3} \quad P_4 \quad 45 \quad \overbrace{5 \quad 6 \quad 56}^{P_5} \\ \left(\begin{array}{c} P_1 \quad 4 \\ P_2 \left\{ \begin{array}{l} 25|_4 \\ 35|_4 \\ 36|_4 \\ 235|_4 \end{array} \right. \\ P_3 \left\{ \begin{array}{l} 25|_{4,\infty} \\ 35|_{4,\infty} \\ 36|_{4,\infty} \\ 235|_{4,\infty} \end{array} \right. \\ P_4 \quad 45 \\ P_5 \left\{ \begin{array}{l} 5 \\ 6 \\ 56 \end{array} \right. \end{array} \right. \end{array}$$

Q_1

1

M_1

Q_2

M_2

M_3

M_4

It's convenient to construct ∂_4^+ using the above order. However, since this order is neither topologically consistent nor filtration compatible with $Y_i \subseteq Y_{i+1} \subseteq \cdots \subseteq Y_{i,N}$ (see Definition 4.2,[3]), we have to reorder those cells in $X_{i,\infty}$.

Definition 4.10. Positive Order

Order cells in $X_{i,\infty}$ first according to t_{max} increasingly, then according to the initial order, we call this order positive order.

For example, τ_1 and τ_2 are two cells in $X_{i,\infty}$. If $t_{max}(\tau_1) < t_{max}(\tau_2)$, then $\tau_1 < \tau_2$ in positive order. If $t_{max}(\tau_1) = t_{max}(\tau_2)$, according to the note after Definition 4.7, $\tau_j (j = 1, 2)$ can be uniquely represented by $\sigma_j, \sigma_j|_i$ or $\sigma_j|_{i,\infty}$ for some $\sigma_j \in X$. If $\sigma_1 < \sigma_2$ in initial order, we require $\tau_1 < \tau_2$ in positive order.

Notice that P_1 and P_2 do not change in positive order. The cells in P_3 , P_4 and P_5 permute. Denote the cells in P_3 , P_4 and P_5 in positive order by $\tilde{P}_{3,4,5}$. Denote the boundary matrix of $X_{i,\infty}$ by $\tilde{\partial}_i^+$ when the cells are in positive order.

Example 4.4. Consider Example 4.3, reorder all cells in $X_{4,\infty}$ in positive order, we have

$$\tilde{\partial}_4^+ := \begin{matrix} & \begin{matrix} P_1 & \overbrace{P_2} & \overbrace{\tilde{P}_{3,4,5}} \end{matrix} \\ & \begin{matrix} 4 & 25|_4 & 35|_4 & 36|_4 & 235|_4 & 5 & 25|_{4,\infty} & 35|_{4,\infty} & 45 & 235|_{4,\infty} & 6 & 36|_{4,\infty} & 56 \end{matrix} \\ \begin{matrix} P_1 \\ P_2 \\ \tilde{P}_{3,4,5} \end{matrix} & \begin{matrix} 4 \\ \left\{ \begin{matrix} 25|_4 \\ 35|_4 \\ 36|_4 \\ 235|_4 \end{matrix} \right\} \\ \left\{ \begin{matrix} 5 \\ 25|_{4,\infty} \\ 35|_{4,\infty} \\ 45 \\ 235|_{4,\infty} \\ 6 \\ 36|_{4,\infty} \\ 56 \end{matrix} \right\} \end{matrix} \end{matrix} \left(\begin{array}{cccccccccccc} & & & & & & & & 1 & & & & \\ & & & & & & & & & & & & \\ & & & 1 & & 1 & & & & & & & \\ & & & 1 & & & 1 & & & & & & \\ & & & & & & & & & & 1 & & \\ & & & & & & & & & 1 & & & \\ & & & & & & 1 & 1 & 1 & & & & 1 \\ & & & & & & & & & 1 & & & \\ & & & & & & & & & 1 & & & \\ & & & & & & & & & & & 1 & 1 \end{array} \right)$$

From $\tilde{\partial}_i^+$, we will get its reduced form $R(\partial_i^+)$ and read the bar code directly from it.

Given an $m \times m$ matrix ∂ with \mathbb{Z}_2 coefficients. Let $low(j)$ be the row index of lowest 1 in column j . If the entire column is zero, then $low(j)$ is undefined. We call ∂ *reduced* if for any two non-zero columns j and j_0 , we have $low(j) \neq low(j_0)$. See page 153,[19]. The following algorithm reduces ∂ by adding columns from left to right.

Algorithm 4.1.

```

for  $j = 1$  to  $m - 1$  do
  if column  $j$  is nonzero do
    while there exists  $j_0 > j$  with  $low(j_0) = low(j)$  do
      add column  $j$  to column  $j_0$ 
    endwhile
    let  $\{i_1, \dots, i_k\} = \{i | 1 \leq i \leq j - 1, low(i) < low(j)\}$ 
      such that  $low(i_1) > low(i_2) > \dots > low(i_k)$ 
    for  $l=1$  to  $k$  do
      while there exists  $j_0 > j$  with  $low(j_0) = low(l)$  do
        add column  $l$  to column  $j_0$ 
      endwhile
    endfor
  endif
endfor

```


$$N_4 = \{\sigma \in X \mid \dim(\sigma) > 0 \text{ and } t_{\max}(\sigma) = i\}$$

$$N_5 = \{\sigma \in X \mid t_{\max}(\sigma) < i\}$$

If we order all cells in $X_{-\infty,i}$ first according to these five classes then according to the initial order, we can easily construct ∂_i^- from the submatrices of ∂ . Third, order cells in $X_{-\infty,i}$ first according to t_{\min} decreasingly then according to the initial order (so called *negative order*), the boundary matrix of $X_{-\infty,i}$ becomes $\tilde{\partial}_i^-$. Fourth, apply Algorithm 4.1 to $\tilde{\partial}_i^-$ and get reduced form $R(\tilde{\partial}_i^-)$. Finally, read the negative bar code $B_r^-(f; i)$ from the reduced form $R(\tilde{\partial}_i^-)$.

Notice: The first $|P_1| + |P_2|$ columns of ∂_i^+ and ∂_i^- are essentially the same, since $P_1 \cup P_2$ and $N_1 \cup N_2$ contain the same cells in the same order. It's the same situation for $\tilde{\partial}_i^+$ and $\tilde{\partial}_i^-$, since reordering cells in positive and negative order will not affect the first $|P_1| + |P_2|$ cells. The same thing happened for the reduced form $R(\tilde{\partial}_i^+)$ and $R(\tilde{\partial}_i^-)$, since we apply the same algorithm to them and we only add columns from left to right in that algorithm. We have the same amount of zero columns in $R(\tilde{\partial}_i^+)$ and $R(\tilde{\partial}_i^-)$, which are in the same position and correspond to the same cycles of X_i . Therefore, the zero columns in $R(\tilde{\partial}_i^+)$ and $R(\tilde{\partial}_i^-)$ which represent generators of $H_*(X_i)$ are also in one to one correspondence. We can pair up intervals in positive and negative bar codes according to generators.

4.8 Numerical Experiments

The bar codes in this section are generated by Matlab. Positive and negative bar codes are paired up in the given order.

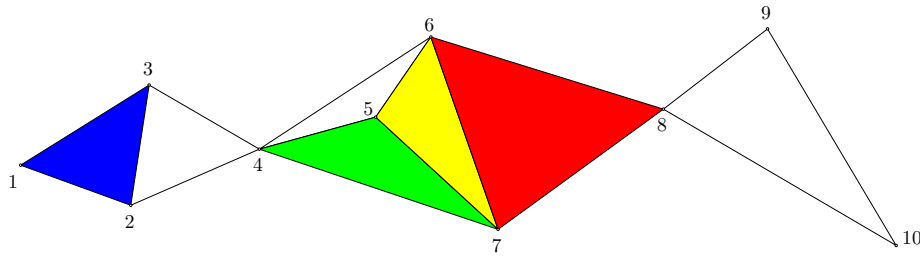


Figure 11: f^1 , a linear map on a 2-simplicial complex.

Example 4.7. Bar codes of f^1 in Figure 11.

$$\begin{array}{ll}
B_0^+(f^1; 1) = \{ [1, \infty) \} & B_0^-(f^1; 1) = \{ (-\infty, 1] \} \\
B_0^+(f^1; 2) = \{ [2, \infty) \} & B_0^-(f^1; 2) = \{ (-\infty, 2] \} \\
B_0^+(f^1; 3) = \{ [3, \infty), [3, 4) \} & B_0^-(f^1; 3) = \{ (-\infty, 3], (2, 3] \} \\
B_0^+(f^1; 4) = \{ [4, \infty) \} & B_0^-(f^1; 4) = \{ (-\infty, 4] \} \\
B_0^+(f^1; 5) = \{ [5, \infty), [5, 6) \} & B_0^-(f^1; 5) = \{ (-\infty, 5], (4, 5] \} \\
B_0^+(f^1; 6) = \{ [6, \infty) \} & B_0^-(f^1; 6) = \{ (-\infty, 6] \} \\
B_0^+(f^1; 7) = \{ [7, \infty) \} & B_0^-(f^1; 7) = \{ (-\infty, 7] \} \\
B_0^+(f^1; 8) = \{ [8, \infty) \} & B_0^-(f^1; 8) = \{ (-\infty, 8] \} \\
B_0^+(f^1; 9) = \{ [9, \infty), [9, 10) \} & B_0^-(f^1; 9) = \{ (-\infty, 9], (8, 9] \} \\
B_0^+(f^1; 10) = \{ [10, \infty) \} & B_0^-(f^1; 10) = \{ (-\infty, 10] \}
\end{array}$$

The set of 0-bar codes for level persistence is

$$BL_0(f^1) = \{[1, 10], (2, 4), (4, 6), (8, 10)\}$$

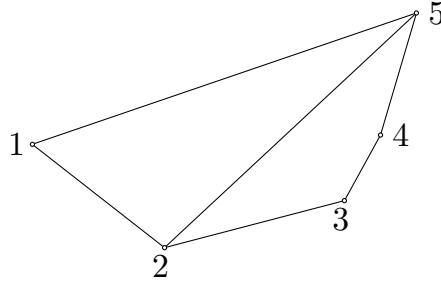


Figure 12: f^2 , a linear map on a planar graph.

Example 4.8. Bar codes of f^2 in Figure 12.

$$\begin{array}{ll}
B_0^+(f^2; 1) = \{ [1, \infty) \} & B_0^-(f^2; 1) = \{ (-\infty, 1] \} \\
B_0^+(f^2; 2) = \{ [2, \infty), [2, 5) \} & B_0^-(f^2; 2) = \{ (-\infty, 2], (1, 2] \} \\
B_0^+(f^2; 3) = \{ [3, \infty), [3, 5), [3, 5) \} & B_0^-(f^2; 3) = \{ (-\infty, 3], (1, 3], (2, 3] \} \\
B_0^+(f^2; 4) = \{ [4, \infty), [4, 5), [4, 5) \} & B_0^-(f^2; 4) = \{ (-\infty, 4], (1, 4], (2, 4] \} \\
B_0^+(f^2; 5) = \{ [5, \infty) \} & B_0^-(f^2; 5) = \{ (-\infty, 5] \}
\end{array}$$

The set of 0-bar codes for level persistence is

$$BL_0(f^2) = \{[1, 5], (1, 5), (2, 5)\}$$

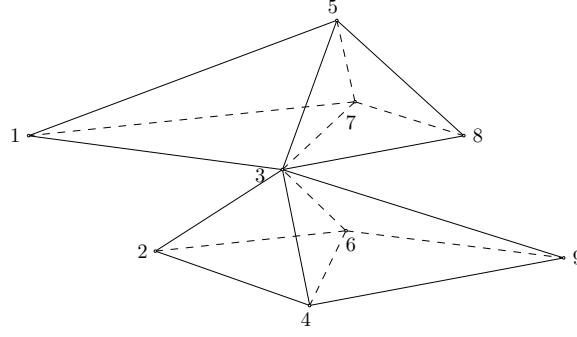


Figure 13: f^3 , a linear map on the union of the surfaces of two tetrahedra.

Example 4.9. Bar codes of f^3 in Figure 13.

$$\begin{array}{ll}
B_0^+(f^3; 1) = \{ [1, \infty) \} & B_0^-(f^3; 1) = \{ (-\infty, 1] \} \\
B_0^+(f^3; 2) = \{ [2, \infty), [2, 3) \} & B_0^-(f^3; 2) = \{ (-\infty, 2], (-\infty, 2] \} \\
B_0^+(f^3; 3) = \{ [3, \infty) \} & B_0^-(f^3; 3) = \{ (-\infty, 3] \} \\
B_0^+(f^3; 4) = \{ [4, \infty), [4, \infty) \} & B_0^-(f^3; 4) = \{ (-\infty, 4], (3, 4] \} \\
B_0^+(f^3; 5) = \{ [5, \infty), [5, \infty) \} & B_0^-(f^3; 5) = \{ (-\infty, 5], (3, 5] \} \\
B_0^+(f^3; 6) = \{ [6, \infty), [6, \infty) \} & B_0^-(f^3; 6) = \{ (-\infty, 6], (3, 6] \} \\
B_0^+(f^3; 7) = \{ [7, \infty), [7, \infty) \} & B_0^-(f^3; 7) = \{ (-\infty, 7], (3, 7] \} \\
B_0^+(f^3; 8) = \{ [8, \infty), [8, \infty) \} & B_0^-(f^3; 8) = \{ (-\infty, 8], (3, 8] \} \\
B_0^+(f^3; 9) = \{ [9, \infty) \} & B_0^-(f^3; 9) = \{ (-\infty, 9] \} \\
B_1^+(f^3; 2) = \{ [2, 8) \} & B_1^-(f^3; 2) = \{ (1, 2] \} \\
B_1^+(f^3; 3) = \{ [3, 8), [3, 9) \} & B_1^-(f^3; 3) = \{ (1, 3], (2, 3] \} \\
B_1^+(f^3; 4) = \{ [4, 9), [4, 8) \} & B_1^-(f^3; 4) = \{ (2, 4], (1, 4] \} \\
B_1^+(f^3; 5) = \{ [5, 8), [5, 9) \} & B_1^-(f^3; 5) = \{ (1, 5], (2, 5] \} \\
B_1^+(f^3; 6) = \{ [6, 9), [6, 8) \} & B_1^-(f^3; 6) = \{ (2, 6], (1, 6] \} \\
B_1^+(f^3; 7) = \{ [7, 9), [7, 8) \} & B_1^-(f^3; 7) = \{ (2, 7], (1, 7] \} \\
B_1^+(f^3; 8) = \{ [8, 9) \} & B_1^-(f^3; 8) = \{ (2, 8] \}
\end{array}$$

The set of 0-bar codes for level persistence is

$$BL_0(f^3) = \{[1, 9], [2, 3), (3, 8]\}$$

The set of 1-bar codes for level persistence is

$$BL_1(f^3) = \{(1, 8), (2, 9)\}$$

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